# Table of Contents

1 About This Guide .............................................. 2
  1.1 What This Reference Manual Contains ...................... 2
  1.2 Conventions ............................................ 3
  1.3 Related Information ..................................... 3

2 Implementation Defined Pragmas .......................... 5
  2.1 Pragma Abort_Defer ...................................... 5
  2.2 Pragma Abstract_State ................................... 5
  2.3 Pragma Ada_83 .............................................. 6
  2.4 Pragma Ada_95 ............................................. 7
  2.5 Pragma Ada_05 ............................................. 7
  2.6 Pragma Ada_2005 .......................................... 7
  2.7 Pragma Ada_12 ............................................ 7
  2.8 Pragma Ada_2012 .......................................... 8
  2.9 Pragma Ada_2022 .......................................... 8
  2.10 Pragma Aggregate_Individually_Assign ..................... 8
  2.11 Pragma Allow_Integer_Address ............................ 8
  2.12 Pragma Annotate ......................................... 9
  2.13 Pragma Assert ........................................... 10
  2.14 Pragma Assert_And_Cut .................................. 11
  2.15 Pragma Assertion_Policy ................................ 11
  2.16 Pragma Assume .......................................... 12
  2.17 Pragma Assume_No_Invalid_Values ......................... 13
  2.18 Pragma Async_Readers .................................... 13
  2.19 Pragma Async_Writers .................................... 13
  2.20 Pragma Attribute_Definition ............................. 14
  2.21 Pragma C_Pass_By_Copy .................................. 14
  2.22 Pragma Check ............................................ 14
  2.23 Pragma Check_Float_Overflow ............................. 15
  2.24 Pragma Check_Name ...................................... 16
  2.25 Pragma Check_Policy .................................... 16
  2.26 Pragma Comment ......................................... 17
  2.27 Pragma Common_Object .................................. 17
  2.28 Pragma Compile_Time_Error .............................. 18
  2.29 Pragma Compile_Time_Warning ........................... 18
  2.30 Pragma Complete_Representation ........................ 18
  2.31 Pragma Complex_Representation ........................ 19
  2.32 Pragma Component_Alignment ............................ 19
  2.33 Pragma Constant_After_Elaboration ....................... 20
  2.34 Pragma Contract_Cases .................................. 20
  2.35 Pragma Convention_Identifier ........................... 21
  2.36 Pragma CPP_Class ....................................... 22
  2.37 Pragma CPP_Constructor ................................ 22
2.38 Pragma CPP_Virtual  ........................................ 23
2.39 Pragma CPP_Vtable .......................................... 23
2.40 Pragma CPU .................................................. 23
2.41 Pragma Deadline_Floor ....................................... 23
2.42 Pragma Debug .............................................. 23
2.43 Pragma Debug_Policy ......................................... 24
2.44 Pragma Default_Initial_Condition ......................... 24
2.45 Pragma Default_Scalar_Storage_Order .................... 24
2.46 Pragma Default_Storage_Pool ............................... 25
2.47 Pragma Depends ............................................ 25
2.48 Pragma Detect_Blocking ..................................... 26
2.49 Pragma Disable_Atomic_Synchronization .................. 26
2.50 Pragma Dispatching_Domain ................................ 26
2.51 Pragma Effective_Reads .................................... 27
2.52 Pragma Effective_Writes ................................... 27
2.53 Pragma Elaboration_Checks .................................. 27
2.54 PragmaEliminate ............................................ 27
2.55 Pragma Enable_Atomic_Synchronization ................... 29
2.56 Pragma Export_Function ..................................... 30
2.57 Pragma Export_Object ....................................... 31
2.58 Pragma Export_Procedure ................................... 31
2.59 Pragma Export_Valued_Procedure .......................... 32
2.60 Pragma Extend_System ...................................... 33
2.61 Pragma Extensions_Allowed .................................. 34
2.62 Pragma Extensions_Visible ................................... 37
2.63 Pragma External ............................................ 38
2.64 Pragma External_Name_Casing ............................... 38
2.65 Pragma Fast_Math ........................................... 39
2.66 Pragma Favor_Top_Level ..................................... 39
2.67 Pragma Finalize_Storage_Only .............................. 40
2.68 Pragma Float_Representation ............................... 40
2.69 Pragma Ghost ............................................... 40
2.70 Pragma Global ............................................... 40
2.71 Pragma Ident ................................................ 41
2.72 Pragma Ignore_Pragma ....................................... 41
2.73 Pragma Implementation_Defined .............................. 41
2.74 Pragma Implemented ......................................... 41
2.75 Pragma Implicit_Packing ..................................... 42
2.76 Pragma Import_Function ...................................... 43
2.77 Pragma Import_Object ........................................ 44
2.78 Pragma Import_Procedure .................................... 44
2.79 Pragma Import_Valued_Procedure .......................... 45
2.80 Pragma Independent ........................................... 46
2.81 Pragma Independent_Components ............................ 46
2.83 Pragma Initialize_Scalars ................................... 46
2.84 Pragma Initializes ............................................ 48
2.85 Pragma Inline_Always ......................................... 48
2.86 Pragma Inline_Generic ...................................... 49
2.87 Pragma Interface ............................................. 49
2.88 Pragma Interface_Name ....................................... 49
2.89 Pragma Interrupt_Handler ................................. 49
2.90 Pragma Interrupt_State ..................................... 50
2.91 Pragma Invariant ............................................. 51
2.92 Pragma Keep_Names ......................................... 51
2.93 Pragma License ................................................. 51
2.94 Pragma Link_With ............................................. 53
2.95 Pragma Linker_Alias ........................................ 53
2.96 Pragma Linker_Destructor .................................... 53
2.97 Pragma Linker_Destructor .................................... 54
2.98 Pragma Linker_Section ...................................... 54
2.99 Pragma Lock_Free ............................................. 55
2.100 Pragma Loop_Invariant .................................... 56
2.101 Pragma Loop_Optimize ..................................... 56
2.102 Pragma Loop_Variant ....................................... 57
2.103 Pragma Machine_Attribute ................................. 58
2.104 Pragma Main ................................................ 58
2.105 Pragma Main_Storage ....................................... 58
2.106 Pragma Max_Queue_Length ................................. 58
2.107 Pragma No_Body .............................................. 59
2.108 Pragma No_Caching .......................................... 59
2.109 Pragma No_Component_Reordering ......................... 59
2.110 Pragma No Elaboration_Code_All ......................... 59
2.111 Pragma No_Heap_Finalization .............................. 60
2.112 Pragma No_Inline ........................................... 60
2.113 Pragma No_Return .......................................... 60
2.114 Pragma No_Strict_Aliasing ................................. 61
2.115 Pragma No_Tagged_Streams ................................. 61
2.116 Pragma Normalize_Scalars ................................. 61
2.117 Pragma Obsolescent ....................................... 62
2.118 Pragma Optimize_Alignment ............................... 64
2.119 Pragma Ordered ............................................. 65
2.120 Pragma Overflow_Mode ..................................... 66
2.121 Pragma Overriding_Renamings ............................ 67
2.122 Pragma Part_Of ............................................. 67
2.123 Pragma Partition_Elaboration_Policy ...................... 68
2.124 Pragma Passive ............................................. 68
2.125 Pragma Persistent_BSS ..................................... 68
2.126 Pragma Post ................................................ 68
2.127 Pragma Postcondition ..................................... 69
2.128 Pragma Post_Class .......................................... 71
2.129 Pragma Pre ................................................ 71
2.130 Pragma Precondition ...................................... 72
2.131 Pragma Predicate .......................................... 72
2.132 Pragma Predicate_Failure ................................. 73
2.133 Pragma Preelaborable_Initialization ...................... 73
2.134 Pragma Prefix_Exception_Messages                      73
2.135 Pragma Pre_Class                                      74
2.136 Pragma Priority_Specific_Dispatching                 74
2.137 Pragma Profile                                        74
2.138 Pragma Profile_Warnings                              77
2.139 Pragma Propagate_Exceptions                          77
2.140 Pragma Provide_Shift_Operators                        78
2.141 Pragma Psect_Object                                  78
2.142 Pragma Pure_Function                                  78
2.143 Pragma Rational                                      79
2.144 Pragma Ravenscar                                      79
2.145 Pragma Refined_Deps                                  79
2.146 Pragma Refined_Global                                 80
2.147 Pragma Refined_Post                                   80
2.148 Pragma Refined_State                                 80
2.149 Pragma Relative_Deadline                             81
2.150 Pragma Remote_Access_Type                            81
2.151 Pragma Rename_Pragma                                  81
2.152 Pragma Restricted_Run_Time                            82
2.153 Pragma Restriction_Warnings                           82
2.154 Pragma Reviewable                                    82
2.155 Pragma Secondary_Stack_Size                          83
2.156 Pragma Share_Generic                                 84
2.157 Pragma Shared                                         84
2.158 Pragma Short_Circuit_And_Or                          84
2.159 Pragma Short_Descriptors                             84
2.160 Pragma Simple_Storage_Pool_Type                      85
2.161 Pragma Source_File_Name                              85
2.162 Pragma Source_File_Name_Project                      87
2.163 Pragma Source_Reference                              87
2.164 Pragma SPARK_Mode                                    87
2.165 Pragma Static_Elaboration_Desired                    88
2.166 Pragma Stream_Convert                                89
2.167 Pragma Style_Checks                                   90
2.168 Pragma Subtitle                                      91
2.169 Pragma Suppress                                       91
2.170 Pragma Suppress_All                                  92
2.171 Pragma Suppress_Debug_Info                           92
2.172 Pragma Suppress_Exception_Locations                   92
2.173 Pragma Suppress_Initialization                        92
2.174 Pragma Task_Name                                      93
2.175 Pragma Task_Storage                                   94
2.176 Pragma Test_Case                                      94
2.177 Pragma Thread_Local_Storage                          95
2.178 Pragma Time_Slice                                     95
2.179 Pragma Title                                         95
2.180 Pragma Type.Invariant                                 96
2.181 Pragma Type.Invariant.Class                         96
3 Implementation Defined Aspects ............. 108

3.1 Aspect Abstract_State ........................... 108
3.2 Aspect Annotate ................................ 108
3.3 Aspect Async_Readers ......................... 109
3.4 Aspect Async_Writers ......................... 109
3.5 Aspect Constant_After_Elab .................... 109
3.6 Aspect Contract_Cases ......................... 109
3.7 Aspect Depends ................................ 109
3.8 Aspect Default_Initial_Condition ............. 109
3.9 Aspect Dimension ............................. 109
3.10 Aspect Dimension_System ...................... 110
3.11 Aspect Disable_Controlled ..................... 110
3.12 Aspect Effective_Reads ....................... 111
3.13 Aspect Effective_Writes ...................... 111
3.14 Aspect Extensions_Visible .................... 111
3.15 Aspect Favor_Top_Level ....................... 111
3.16 Aspect Ghost .................................. 111
3.17 Aspect Global ................................ 111
3.18 Aspect Initial_Condition ..................... 111
3.19 Aspect Initializes ............................ 111
3.20 Aspect Inline_Aways ........................... 111
3.21 Aspect Invariant ............................. 111
3.22 Aspect Invariant_Class ....................... 111
3.23 Aspect Iterable ............................... 112
3.24 Aspect Linker_Section ......................... 112
3.25 Aspect Lock_Free .............................. 112
3.26 Aspect Max_Queue_Length ..................... 113
3.27 Aspect No_Caching ............................ 113
4 Implementation Defined Attributes ...... 116

4.1 Attribute Abort_Signal ......................... 116
4.2 Attribute Address_Size .......................... 116
4.3 Attribute Asm_Input .............................. 116
4.4 Attribute Asm_Output .............................. 116
4.5 Attribute Atomic_Always_Lock_Free .............. 117
4.6 Attribute Bit .......................... 117
4.7 Attribute Bit_Position .......................... 117
4.8 Attribute Code_Address ......................... 117
4.9 Attribute Compiler_Version ..................... 118
4.10 Attribute Constrained ......................... 118
4.11 Attribute Default_Bit_Order .................. 118
4.12 Attribute Default_Scalar_Storage_Order ........ 118
4.13  Attribute Deref ............................................. 118
4.14  Attribute Descriptor_Size .................................. 118
4.15  Attribute Elaborated ........................................ 119
4.16  Attribute Elab_Body ......................................... 119
4.17  Attribute Elab_Spec .......................................... 119
4.18  Attribute Elab_Subp_Body ................................... 119
4.19  Attribute Emax ................................................ 119
4.20  Attribute Enabled ............................................ 119
4.21  Attribute Enum_Rep .......................................... 120
4.22  Attribute Enum_Val ........................................... 120
4.23  Attribute Epsilon ............................................. 120
4.24  Attribute Fast_math ......................................... 121
4.25  Attribute Finalization_Size ................................. 121
4.26  Attribute Fixed_Value ........................................ 121
4.27  Attribute From_Any ........................................... 121
4.28  Attribute Has_Access_Values ............................... 121
4.29  Attribute Has_Discriminants ............................... 121
4.30  Attribute Has_Tagged_Values ............................... 122
4.31  Attribute Img ................................................ 122
4.32  Attribute Initialized ........................................ 122
4.33  Attribute Integer_Value ..................................... 122
4.34  Attribute Invalid_Value ..................................... 122
4.35  Attribute Iterable ........................................... 123
4.36  Attribute Large .............................................. 123
4.37  Attribute Library_Level ..................................... 123
4.38  Attribute Loop_Entry ......................................... 123
4.39  Attribute Machine_Size ..................................... 123
4.40  Attribute Mantissa ........................................... 124
4.41  Attribute Maximum_Alignment ............................. 124
4.42  Attribute Max_Integer_Size ................................. 124
4.43  Attribute Mechanism_Code .................................. 124
4.44  Attribute Null_Parameter .................................... 124
4.45  Attribute Object_Size ....................................... 124
4.46  Attribute Old ................................................ 125
4.47  Attribute Passed_By_Reference ............................. 125
4.48  Attribute Pool_Address ....................................... 126
4.49  Attribute Range_Length ..................................... 126
4.50  Attribute Restriction_Set ................................... 126
4.51  Attribute Result .............................................. 127
4.52  Attribute Safe_Emax ......................................... 127
4.53  Attribute Safe_Large ......................................... 127
4.54  Attribute Safe_Small .......................................... 127
4.55  Attribute Scalar_Storage_Order ............................ 127
4.56  Attribute Simple_Storage_Pool ............................ 130
4.57  Attribute Small ............................................... 131
4.58  Attribute Small_Denominator .............................. 131
4.59  Attribute Small_Numerator .................................. 131
4.60  Attribute Storage_Unit ..................................... 131
5 Standard and Implementation Defined Restrictions ............................. 139

5.1 Partition-Wide Restrictions ............................................. 139
  5.1.1 Immediate_Reclamation ............................................. 139
  5.1.2 Max_Async_Select_Nesting ........................................ 139
  5.1.3 Max_Entry_Queue_Length ......................................... 139
  5.1.4 Max_Protected_Entries ............................................ 139
  5.1.5 Max_Select_Alt .................................................... 139
  5.1.6 Max_Storage_AT_Blocking ....................................... 140
  5.1.7 Max_Task_Entries ................................................ 140
  5.1.8 Max_Tasks ......................................................... 140
  5.1.9 No_Abort_Statements ............................................. 140
  5.1.10 No_Access_Parameter_Allocators ................................. 140
  5.1.11 No_Access_Subs ............................................... 140
  5.1.12 No_Allocators .................................................. 140
  5.1.13 No_Anon_Allocators ............................................. 140
  5.1.14 No_Async_Control ............................................... 140
  5.1.15 No_Calendar ................................................... 140
  5.1.16 No_Coextensions ................................................. 140
  5.1.17 No_Default_Initialization ...................................... 141
  5.1.18 No_Delay ....................................................... 141
  5.1.19 No_Dependence .................................................. 141
  5.1.20 No_Direct_Boolean_Ops ........................................ 141
  5.1.21 No_Dispatch ................................................... 141
  5.1.22 No_Dispatching_Calls .......................................... 141
  5.1.23 No_Dynamic_Attach ............................................. 142
  5.1.24 No_Dynamic_Priors ............................................. 142
  5.1.25 No_Entry_Calls_In_Elaboration_Code ............................ 142
  5.1.26 NoEnumeration_Maps ............................................. 143
<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>5.1.27 No_Exception_Handlers</td>
<td>143</td>
</tr>
<tr>
<td>5.1.28 No_Exception_Propagation</td>
<td>143</td>
</tr>
<tr>
<td>5.1.29 No_Exception_Registration</td>
<td>143</td>
</tr>
<tr>
<td>5.1.30 No_Exceptions</td>
<td>143</td>
</tr>
<tr>
<td>5.1.31 No_Finalization</td>
<td>143</td>
</tr>
<tr>
<td>5.1.32 No_Fixed_Point</td>
<td>144</td>
</tr>
<tr>
<td>5.1.33 No_Floating_Point</td>
<td>144</td>
</tr>
<tr>
<td>5.1.34 No_Implicit_Conditionals</td>
<td>144</td>
</tr>
<tr>
<td>5.1.35 No_Implicit_Dynamic_Code</td>
<td>144</td>
</tr>
<tr>
<td>5.1.36 No_Implicit_Heap_Allocations</td>
<td>144</td>
</tr>
<tr>
<td>5.1.37 No_Implicit_Protected_Object_Allocations</td>
<td>145</td>
</tr>
<tr>
<td>5.1.38 No_Implicit_Task_Allocations</td>
<td>145</td>
</tr>
<tr>
<td>5.1.39 No_Initialize_Scalars</td>
<td>145</td>
</tr>
<tr>
<td>5.1.40 No_IO</td>
<td>145</td>
</tr>
<tr>
<td>5.1.41 No_Local_Allocators</td>
<td>145</td>
</tr>
<tr>
<td>5.1.42 No_Local_Protected_Objects</td>
<td>145</td>
</tr>
<tr>
<td>5.1.43 No_Local_Tagged_Types</td>
<td>145</td>
</tr>
<tr>
<td>5.1.44 No_Local_Timing_Events</td>
<td>145</td>
</tr>
<tr>
<td>5.1.45 No_Long_Long_Integers</td>
<td>145</td>
</tr>
<tr>
<td>5.1.46 No_Multiple_Elaboration</td>
<td>145</td>
</tr>
<tr>
<td>5.1.47 No_Nested_Finalization</td>
<td>146</td>
</tr>
<tr>
<td>5.1.48 No_Protected_Type_Allocators</td>
<td>146</td>
</tr>
<tr>
<td>5.1.49 No_Protected_Types</td>
<td>146</td>
</tr>
<tr>
<td>5.1.50 No_Recursion</td>
<td>146</td>
</tr>
<tr>
<td>5.1.51 No_Reentrancy</td>
<td>146</td>
</tr>
<tr>
<td>5.1.52 No_Relative_Delay</td>
<td>146</td>
</tr>
<tr>
<td>5.1.53 No_Requeue_Statements</td>
<td>146</td>
</tr>
<tr>
<td>5.1.54 No_Secondary_Stack</td>
<td>146</td>
</tr>
<tr>
<td>5.1.55 No_Select_Statements</td>
<td>146</td>
</tr>
<tr>
<td>5.1.56 No_Specific_Termination_Handlers</td>
<td>146</td>
</tr>
<tr>
<td>5.1.57 No_Specification_of_Aspect</td>
<td>147</td>
</tr>
<tr>
<td>5.1.58 No_Standard_Allocators_After_Elaboration</td>
<td>147</td>
</tr>
<tr>
<td>5.1.59 No_Standard_Storage_Pools</td>
<td>147</td>
</tr>
<tr>
<td>5.1.60 No_Stream_Optimizations</td>
<td>147</td>
</tr>
<tr>
<td>5.1.61 No_Streams</td>
<td>147</td>
</tr>
<tr>
<td>5.1.62 No_Tagged_Type_Registration</td>
<td>147</td>
</tr>
<tr>
<td>5.1.63 No_Task_Allocators</td>
<td>147</td>
</tr>
<tr>
<td>5.1.64 No_Task_At_Interrupt_Priority</td>
<td>147</td>
</tr>
<tr>
<td>5.1.65 No_Task_Attributes_Package</td>
<td>148</td>
</tr>
<tr>
<td>5.1.66 No_Task_Hierarchy</td>
<td>148</td>
</tr>
<tr>
<td>5.1.67 No_Task_Termination</td>
<td>148</td>
</tr>
<tr>
<td>5.1.68 No_Tasking</td>
<td>148</td>
</tr>
<tr>
<td>5.1.69 No_Terminate_Alternatives</td>
<td>148</td>
</tr>
<tr>
<td>5.1.70 NoUnchecked_Access</td>
<td>148</td>
</tr>
<tr>
<td>5.1.71 NoUnchecked_Conversion</td>
<td>148</td>
</tr>
<tr>
<td>5.1.72 NoUnchecked_Deallocation</td>
<td>148</td>
</tr>
<tr>
<td>5.1.73 No_Use_Of_Entity</td>
<td>148</td>
</tr>
<tr>
<td>5.1.74 Pure_Barriers</td>
<td>148</td>
</tr>
</tbody>
</table>
5.1.75 Simple Barriers ........................................... 149
5.1.76 Static Priorities ........................................ 149
5.1.77 Static Storage Size ..................................... 149

5.2 Program Unit Level Restrictions .......................... 149
5.2.1 No Elaboration Code ..................................... 149
5.2.2 No Dynamic Accessibility Checks ..................... 150
5.2.3 No Dynamic Sized Objects ............................... 151
5.2.4 No Entry Queue ........................................... 151
5.2.5 No Implementation Aspect Specifications ............. 151
5.2.6 No Implementation Attributes .......................... 151
5.2.7 No Implementation Identifiers .......................... 151
5.2.8 No Implementation Pragmas .............................. 151
5.2.9 No Implementation Restrictions ........................ 151
5.2.10 No Implementation Units ................................. 152
5.2.11 No Implicit Aliasing .................................... 152
5.2.12 No Implicit Loops ....................................... 152
5.2.13 No Obsolescent Features ............................... 152
5.2.14 No Wide Characters .................................... 152
5.2.15 Static Dispatch Tables ................................ 152
5.2.16 SPARK 05 ................................................. 152

6 Implementation Advice ........................................ 153
6.1 RM 1.1.3(20): Error Detection ............................ 153
6.2 RM 1.1.3(31): Child Units ................................ 153
6.3 RM 1.1.5(12): Bounded Errors ............................ 153
6.4 RM 2.8(16): Pragmas ...................................... 153
6.5 RM 2.8(17-19): Pragmas ................................... 154
6.6 RM 3.5.2(5): Alternative Character Sets ............... 154
6.7 RM 3.5.4(28): Integer Types ............................... 155
6.8 RM 3.5.4(29): Integer Types ............................... 155
6.9 RM 3.5.5(8): Enumeration Values .......................... 155
6.10 RM 3.5.7(17): Float Types ................................ 155
6.11 RM 3.6.2(11): Multidimensional Arrays .................. 156
6.12 RM 9.6(30-31): Duration'Small ......................... 156
6.13 RM 10.2.1(12): Consistent Representation .......... 156
6.14 RM 11.4.1(19): Exception Information .................. 156
6.15 RM 11.5(28): Suppression of Checks .................... 157
6.16 RM 13.1 (21-24): Representation Clauses ............... 157
6.17 RM 13.2(6-8): Packed Types .............................. 157
6.18 RM 13.3(14-19): Address Clauses ........................ 158
6.19 RM 13.3(29-35): Alignment Clauses ...................... 158
6.20 RM 13.3(42-43): Size Clauses ............................. 159
6.21 RM 13.3(50-56): Size Clauses ............................. 159
6.22 RM 13.3(71-73): Component Size Clauses ............... 160
6.23 RM 13.4(9-10): Enumeration Representation Clauses ...... 160
6.24 RM 13.5.1(17-22): Record Representation Clauses ...... 160
6.25 RM 13.5.2(5): Storage Place Attributes ................. 161
6.26 RM 13.5.3(7-8): Bit Ordering ............................ 161
<table>
<thead>
<tr>
<th>Section</th>
<th>RM</th>
<th>Topic</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>6.27</td>
<td>RM 13.7(37): Address as Private</td>
<td>161</td>
<td></td>
</tr>
<tr>
<td>6.28</td>
<td>RM 13.7.1(16): Address Operations</td>
<td>161</td>
<td></td>
</tr>
<tr>
<td>6.29</td>
<td>RM 13.9(14-17): Unchecked Conversion</td>
<td>161</td>
<td></td>
</tr>
<tr>
<td>6.30</td>
<td>RM 13.11(23-25): Implicit Heap Usage</td>
<td>162</td>
<td></td>
</tr>
<tr>
<td>6.31</td>
<td>RM 13.11.2(17): Unchecked Deallocation</td>
<td>162</td>
<td></td>
</tr>
<tr>
<td>6.32</td>
<td>RM 13.13.2(1.6): Stream Oriented Attributes</td>
<td>162</td>
<td></td>
</tr>
<tr>
<td>6.33</td>
<td>RM A.1(52): Names of Predefined Numeric Types</td>
<td>163</td>
<td></td>
</tr>
<tr>
<td>6.36</td>
<td>RM A.5.2(46-47): Random Number Generation</td>
<td>163</td>
<td></td>
</tr>
<tr>
<td>6.37</td>
<td>RM A.10.7(23): Get_Immediate</td>
<td>164</td>
<td></td>
</tr>
<tr>
<td>6.38</td>
<td>RM A.18: Containers</td>
<td>164</td>
<td></td>
</tr>
<tr>
<td>6.39</td>
<td>RM B.1(39-41): Pragma Export</td>
<td>164</td>
<td></td>
</tr>
<tr>
<td>6.40</td>
<td>RM B.2(12-13): Package Interfaces</td>
<td>165</td>
<td></td>
</tr>
<tr>
<td>6.41</td>
<td>RM B.3(63-71): Interfacing with C</td>
<td>165</td>
<td></td>
</tr>
<tr>
<td>6.42</td>
<td>RM B.4(95-98): Interfacing with COBOL</td>
<td>166</td>
<td></td>
</tr>
<tr>
<td>6.43</td>
<td>RM B.5(22-26): Interfacing with Fortran</td>
<td>166</td>
<td></td>
</tr>
<tr>
<td>6.46</td>
<td>RM C.3(28): Interrupt Support</td>
<td>168</td>
<td></td>
</tr>
<tr>
<td>6.47</td>
<td>RM C.3.1(20-21): Protected Procedure Handlers</td>
<td>168</td>
<td></td>
</tr>
<tr>
<td>6.48</td>
<td>RM C.3.2(25): Package Interrupts</td>
<td>168</td>
<td></td>
</tr>
<tr>
<td>6.49</td>
<td>RM C.4(14): Pre-elaboration Requirements</td>
<td>168</td>
<td></td>
</tr>
<tr>
<td>6.50</td>
<td>RM C.5(8): Pragma Discard_Names</td>
<td>168</td>
<td></td>
</tr>
<tr>
<td>6.51</td>
<td>RM C.7.2(30): The Package Task_Attributes</td>
<td>168</td>
<td></td>
</tr>
<tr>
<td>6.52</td>
<td>RM D.3(17): Locking Policies</td>
<td>169</td>
<td></td>
</tr>
<tr>
<td>6.53</td>
<td>RM D.4(16): Entry Queuing Policies</td>
<td>169</td>
<td></td>
</tr>
<tr>
<td>6.54</td>
<td>RM D.6(9-10): Preemptive Abort</td>
<td>169</td>
<td></td>
</tr>
<tr>
<td>6.55</td>
<td>RM D.7(21): Tasking Restrictions</td>
<td>169</td>
<td></td>
</tr>
<tr>
<td>6.56</td>
<td>RM D.8(47-49): Monotonic Time</td>
<td>169</td>
<td></td>
</tr>
<tr>
<td>6.57</td>
<td>RM E.5(28-29): Partition Communication Subsystem</td>
<td>170</td>
<td></td>
</tr>
<tr>
<td>6.58</td>
<td>RM F(7): COBOL Support</td>
<td>170</td>
<td></td>
</tr>
<tr>
<td>6.59</td>
<td>RM F.1(2): Decimal Radix Support</td>
<td>170</td>
<td></td>
</tr>
<tr>
<td>6.60</td>
<td>RM G: Numerics</td>
<td>170</td>
<td></td>
</tr>
<tr>
<td>6.61</td>
<td>RM G.1.1(56-58): Complex Types</td>
<td>171</td>
<td></td>
</tr>
<tr>
<td>6.62</td>
<td>RM G.1.2(49): Complex Elementary Functions</td>
<td>171</td>
<td></td>
</tr>
<tr>
<td>6.63</td>
<td>RM G.2.4(19): Accuracy Requirements</td>
<td>172</td>
<td></td>
</tr>
<tr>
<td>6.64</td>
<td>RM G.2.6(15): Complex Arithmetic Accuracy</td>
<td>172</td>
<td></td>
</tr>
<tr>
<td>6.65</td>
<td>RM H.6(15/2): Pragma Partition_Elaboration_Policy</td>
<td>172</td>
<td></td>
</tr>
</tbody>
</table>

7 Implementation Defined Characteristics . . . 173
8 Intrinsic Subprograms ................. 192
  8.1 Intrinsic Operators .................. 192
  8.2 Compilation_ISO_Date ................. 192
  8.3 Compilation_Date .................... 192
  8.4 Compilation_Time .................... 192
  8.5 Enclosing_Entity .................... 193
  8.6 Exception_Information ............... 193
  8.7 Exception_Message ................... 193
  8.8 Exception_Name ...................... 193
  8.9 File ................................ 193
  8.10 Line ................................ 193
  8.11 Shifts and Rotates .................. 194
  8.12 Source_Location ..................... 194

9 Representation Clauses and Pragmas ...... 195
  9.1 Alignment Clauses .................... 195
  9.2 Size Clauses .......................... 196
  9.3 Storage_Size Clauses ................. 197
  9.4 Size of Variant Record Objects ........ 198
  9.5 Biased Representation ................ 200
  9.6 Value_Size and Object_Size Clauses .... 200
  9.7 Component_Size Clauses .............. 203
  9.8 Bit_Order Clauses .................... 204
  9.9 Effect of Bit_Order on Byte Ordering .. 205
  9.10 Pragma Pack for Arrays .............. 209
  9.11 Pragma Pack for Records ............. 211
  9.12 Record Representation Clauses ....... 212
  9.13 Handling of Records with Holes ....... 213
  9.14 Enumeration Clauses ................. 214
  9.15 Address Clauses ..................... 215
  9.16 Use of Address Clauses for Memory-Mapped I/O .... 220
  9.17 Effect of Convention on Representation ...... 220
  9.18 Conventions and Anonymous Access Types .... 221
  9.19 Determining the Representations chosen by GNAT .......... 223

10 Standard Library Routines ............. 226

11 The Implementation of Standard I/O ..... 237
  11.1 Standard I/O Packages ................ 237
  11.2 FORM Strings ......................... 238
  11.3 Direct_IO ............................. 238
  11.4 Sequential_IO .......................... 238
  11.5 Text_IO ............................... 239
    11.5.1 Stream Pointer Positioning ....... 240
    11.5.2 Reading and Writing Non-Regular Files ...... 240
    11.5.3 Get_Immediate ..................... 241
    11.5.4 Treating Text_IO Files as Streams .......... 241
12 The GNAT Library .................................. 255

12.2 Ada.Characters.Wide_Latin_1 (a-cwila1.ads) .......... 255
12.4 Ada.Characters.Wide_Wide_Latin_1 (a-chzl1.ads) .... 255
12.5 Ada.Characters.Wide_Wide_Latin_9 (a-chzla9.ads) ... 256
12.6 Ada.Containers.Bounded_Holders (a-coboho.ads) ... 256
12.7 Ada.Command_Line.Environment (a-colien.ads) ..... 256
12.8 Ada.Command_Line.Remove (a-colire.ads) ........... 256
12.9 Ada.Command_Line.Response_File (a-clrefi.ads) ... 256
12.10 Ada.Direct_IO.C_Streams (a-diocst.ads) .......... 256
12.11 Ada.Exceptions.Is_Null_Occurrence (a-einuoc.ads) . 256
12.12 Ada.Exceptions.Last_Chance_Handler (a-elchha.ads) 256
12.13 Ada.Exceptions.Traceback (a-exctra.ads) ......... 257
12.14 Ada.Sequential_IO.C_Streams (a-siocst.ads) .......... 257
12.15 Ada.Strings.Wide_Direct_IO.C_Streams (a-ssicst.ads) 257
12.16 Ada.Strings.Unbounded.Text_IO (a-suteio.ads) ... 257
12.17 Ada.Strings.Wide_Unbounded.Wide_Text_IO (a-swuwti.ads) ................................................. 257
12.18 Ada.Strings.Wide_Wide_Unbounded.Wide_Wide_Text_IO
(a-szuzti.ads) ......................................... 257
12.19 Ada.Task_Initialization (a-tasini.ads) ............. 257
12.20 Ada.Text_IO.C_Streams (a-tiocst.ads) ............ 257
12.21 Ada.Text_IO.Reset_Standard_Files (a-tirsfi.ads) .... 257
12.22 Ada.Wide_Characters.Unicode (a-wichun.ads) .... 258
12.23 Ada.Wide_Text_IO.C_Streams (a-wtcstr.ads) ....... 258
12.24 Ada.Wide_Text_IO.Reset_Standard_Files (a-wrstfi.ads)
........................................................................... 258
12.25 Ada.Wide_Wide_Characters.Unicode (a-zchuni.ads) . 258
12.26 Ada.Wide_Wide_Text_IO.C_Streams (a-ztcstr.ads) ..... 258

11.5.5 Text_IO Extensions ................................. 241
11.5.6 Text_IO Facilities for Unbounded Strings ........ 241
11.6 Wide_Text_IO ........................................ 242
11.6.1 Stream Pointer Positioning ...................... 244
11.6.2 Reading and Writing Non-Regular Files ........ 244
11.7 Wide_Wide_Text_IO .................................. 245
11.7.1 Stream Pointer Positioning ..................... 246
11.7.2 Reading and Writing Non-Regular Files ........ 246
11.8 Stream_IO ........................................... 246
11.9 Text Translation .................................... 247
11.10 Shared Files ....................................... 247
11.11 Filenames encoding ................................ 248
11.12 File content encoding ............................. 248
11.13 Open Modes ........................................ 249
11.14 Operations on C Streams .......................... 249
11.15 Interfacing to C Streams .......................... 252
12.27 Ada.Wide_Wide_Text_IO.Reset_Standard_Files
(a-zrstfi.ads)........................................ 258
12.28 GNAT.Altivec (g-altive.ads)............................... 258
12.29 GNAT.Altivec.Conversions (g-altcon.ads)............. 258
12.30 GNAT.Altivec.Vector_Operations (g-alveop.ads)...... 258
12.31 GNAT.Altivec.Vector_TYPES (g-alvety.ads)........... 259
12.32 GNAT.Altivec.Vector_VIEWS (g-alvevi.ads)........... 259
12.33 GNAT.Array_Split (g-arrspl.ads).......................... 259
12.34 GNAT.AWK (g-awk.ads)..................................... 259
12.35 GNAT.Binary_Search (g-binsea.ads).................... 259
12.36 GNAT.Bind_Environment (g-binenv.ads).................. 259
12.37 GNAT.Branch_Prediction (g-brapre.ads)................. 259
12.38 GNAT.Bounded_Buffers (g-bobuf.ads)................... 259
12.39 GNAT.Bounded_Mailboxes (g-boumai.ads)................ 259
12.40 GNAT.Bubble_Sort (g-bubsor.ads)....................... 259
12.41 GNAT.Bubble_Sort_A (g-busora.ads)..................... 260
12.42 GNAT.Bubble_Sort_G (g-busorg.ads)..................... 260
12.43 GNAT.Byte_Order_Mark (g-byorma.ads).................. 260
12.44 GNAT.Byte_Swapping (g-bytswa.ads)...................... 260
12.45 GNAT.Calendar (g-calend.ads)............................ 260
12.46 GNAT.Calendar.Time_IO (g-catiiio.ads).................. 260
12.47 GNAT.CRC32 (g-crc32.ads)............................... 260
12.48 GNAT.Case_Util (g-casuti.ads)........................... 260
12.49 GNAT.CGI (g-cgi.ads)..................................... 260
12.50 GNAT.CGI.Cookie (g-cgicoo.ads).......................... 261
12.51 GNAT.CGI.Debug (g-cgideb.ads).......................... 261
12.52 GNAT.Command_Line (g-comlin.ads)...................... 261
12.53 GNAT.Compiler_Version (g-comver.ads)................... 261
12.54 GNAT.Ctrl_C (g-ctrl_c.ads).............................. 261
12.55 GNAT.Current_Exception (g-curexc.ads).................. 261
12.56 GNAT.Debug_Pools (g-debpoo.ads)......................... 261
12.57 GNAT.Debug_Utilsites (g-debuti.ads).................... 261
12.58 GNAT.Decode_String (g-decstr.ads).................... 261
12.59 GNAT.Decode_UTF8_String (g-deutst.ads)............... 262
12.60 GNAT.Directory_Operations (g-dirop.e.ads)............ 262
12.61 GNAT.Directory_Operations.Iteration (g-dioit.ad)...... 262
12.62 GNAT.Dynamic_HTables (g-dynhta.ads)................... 262
12.63 GNAT.Dynamic_Tables (g-dyntab.ads)..................... 262
12.64 GNAT.Encode_String (g-encstr.ads)..................... 262
12.65 GNAT.Encode_UTF8_String (g-enutst.ads)............... 262
12.66 GNAT.Exception_Actions (g-exact.ad)................... 262
12.67 GNAT.Exception_Traces (g-exutra.ads).................. 262
12.68 GNAT.Exceptions (g-except.ads)........................ 263
12.69 GNAT.Expect (g-expect.ads).............................. 263
12.70 GNATEXPECT.TTY (g-exptty.ads).......................... 263
12.71 GNAT.Float_Control (g-flocon.ads)..................... 263
12.72 GNAT.Formatted_String (g-forstr.ads)................... 263
12.73 GNAT.Generic_Fast_Math_Functions (g-gfmafu.ads) ... 263
12.74 GNAT.Heap_Sort (g-heasor.ads) ......................... 263
12.75 GNAT.Heap_Sort_A (g-hesora.ads) ....................... 264
12.76 GNAT.Heap_Sort_G (g-hesorg.ads) ....................... 264
12.77 GNAT.HTable (g-htable.ads) ......................... 264
12.78 GNAT.IO (g-i.o.ads) .................................. 264
12.79 GNAT.IO_Aux (g-i.o_aux.ads) ....................... 264
12.80 GNAT.Lock_Files (g-locfil.ads) ..................... 264
12.81 GNAT.MBBS_Discrete_Random (g-mbdira.ads) ............ 264
12.82 GNAT.MBBS_Float_Random (g-mbf lr.a.ds) ............. 264
12.83 GNAT.MD5 (g-md5.ads) .................................. 265
12.84 GNAT.Memory_Dump (g-memdum.ads) ............... 265
12.85 GNAT.Most_Recent_Exception (g-moreex.ads) .......... 265
12.86 GNAT.OS_Lib (g-os_lib.ads) ....................... 265
12.87 GNAT.Perfect_Hash_Generators (g-pezage.ads) ........ 265
12.88 GNAT.Random_Numbers (g-rannum.ads) ............ 265
12.89 GNAT.Regexp (g-regexp.ads) ..................... 265
12.90 GNAT.Registry (g-regist.ads) ................... 265
12.91 GNAT.Regpat (g-regpat.ads) .................. 266
12.92 GNAT.Rewrite_Data (g-rewdat.ads) ............... 266
12.93 GNAT.Secondary_Stack_Info (g-sestin.ads) .......... 266
12.94 GNAT.Semaphores (g-semaph.ads) .................. 266
12.95 GNAT.Serial_Communications (g-sercom.ads) ............ 266
12.96 GNAT.SHA1 (g-sha1.ads) .................................. 266
12.97 GNAT.SHA224 (g-sha224.ads) ................................. 266
12.98 GNAT.SHA256 (g-sha256.ads) ................................. 266
12.99 GNAT.SHA384 (g-sha384.ads) ................................. 266
12.100 GNAT.SHA512 (g-sha512.ads) ............................... 267
12.101 GNAT.Signals (g-signal.ads) ............................. 267
12.102 GNAT.Sockets (g-socket.ads) ...................... 267
12.103 GNAT.Source_Info (g-souinf.ads) ................. 267
12.104 GNAT.Spelling_Checker (g-speche.ads) ............. 267
12.105 GNAT.Spelling_Checker_Generic (g-spchge.ads) ....... 267
12.106 GNAT.Spitbol.Patterns (g-spipat.ads) ............ 267
12.107 GNAT.Spitbol (g-spitbo.ads) ................... 267
12.108 GNAT.Spitbol.Table_Boolean (g-sptabo.ads) ........... 267
12.109 GNAT.Spitbol.Table_Integer (g-sptain.ads) ........... 268
12.110 GNAT.Spitbol.Table_VString (g-sptavs.ads) ........... 268
12.111 GNAT.SSE (g-sse.ads) .................................. 268
12.112 GNAT.SSE.Vector_Types (g-ssvety.ads) .............. 268
12.113 GNAT.String_Hash (g-strhas.ads) ................... 268
12.114 GNAT.Strings (g-string.ads) .............................. 268
12.115 GNAT.String_Split (g-strspl.ads) ..................... 268
12.116 GNAT.Table (g-table.ads) ............................ 268
12.117 GNAT.Task_Lock (g-tasloc.ads) ..................... 268
12.118 GNAT.Time_Stamp (g-timsta.ads) ................... 269
12.119 GNAT.Threads (g-thread.ads) ..................... 269
12.120 GNAT.Traceback (g-traceb.ads) ..................... 269
<table>
<thead>
<tr>
<th>Section</th>
<th>Description</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>12.121</td>
<td>GNAT.Traceback.Symbolic (g-trasym.ads)</td>
<td>269</td>
</tr>
<tr>
<td>12.122</td>
<td>GNAT.UTF_32 (g-utf_32.ads)</td>
<td>269</td>
</tr>
<tr>
<td>12.123</td>
<td>GNAT.UTF_32_Spelling_Checker (g-u3spch.ads)</td>
<td>269</td>
</tr>
<tr>
<td>12.124</td>
<td>GNAT.Wide_Spelling_Checker (g-wispch.ads)</td>
<td>269</td>
</tr>
<tr>
<td>12.125</td>
<td>GNAT.Wide_String_Split (g-wistsp.ads)</td>
<td>269</td>
</tr>
<tr>
<td>12.126</td>
<td>GNAT.Wide_Wide_Spelling_Checker (g-zspche.ads)</td>
<td>269</td>
</tr>
<tr>
<td>12.127</td>
<td>GNAT.Wide_Wide_String_Split (g-zistsp.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.128</td>
<td>Interfaces.C.Extensions (i-cexten.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.129</td>
<td>Interfaces.C.Streams (i-cstrea.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.130</td>
<td>Interfaces.Packed_Decimal (i-pacdec.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.131</td>
<td>Interfaces.VxWorks (i-vxwork.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.132</td>
<td>Interfaces.VxWorks.Int_Connection (i-vxinco.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.133</td>
<td>Interfaces.VxWorks.IO (i-vxwoio.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.134</td>
<td>System.Address_Image (s-addima.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.135</td>
<td>System.Assertions (s-assert.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.136</td>
<td>System.Atomic_Counters (s-atocou.ads)</td>
<td>270</td>
</tr>
<tr>
<td>12.137</td>
<td>System.Memory (s-memory.ads)</td>
<td>271</td>
</tr>
<tr>
<td>12.138</td>
<td>System.Multiprocessors (s-multip.ads)</td>
<td>271</td>
</tr>
<tr>
<td>12.139</td>
<td>System.Multiprocessors.Dispatching_Domains</td>
<td>271</td>
</tr>
<tr>
<td></td>
<td>(s-mudido.ads)</td>
<td></td>
</tr>
<tr>
<td>12.140</td>
<td>System.Partition_Interface (s-parint.ads)</td>
<td>271</td>
</tr>
<tr>
<td>12.141</td>
<td>System.Pool_Global (s-pooglo.ads)</td>
<td>271</td>
</tr>
<tr>
<td>12.142</td>
<td>System.Pool_Local (s-pooloc.ads)</td>
<td>271</td>
</tr>
<tr>
<td>12.143</td>
<td>System.Restrictions (s-restri.ads)</td>
<td>271</td>
</tr>
<tr>
<td>12.144</td>
<td>System.Rident (s-rident.ads)</td>
<td>271</td>
</tr>
<tr>
<td>12.145</td>
<td>System.Strings.Stream_Ops (s-ststop.ads)</td>
<td>272</td>
</tr>
<tr>
<td>12.146</td>
<td>System.Unsigned_Types (s-unstyp.ads)</td>
<td>272</td>
</tr>
<tr>
<td>12.147</td>
<td>System.Wch_Cnv (s-wchcnv.ads)</td>
<td>272</td>
</tr>
<tr>
<td>12.148</td>
<td>System.Wch_Con (s-wchcon.ads)</td>
<td>272</td>
</tr>
</tbody>
</table>

13 Interfacing to Other Languages ............ 273
13.1 Interfacing to C .................................. 273
13.2 Interfacing to C++ ................................. 274
13.3 Interfacing to COBOL ............................... 274
13.4 Interfacing to Fortran ............................. 275
13.5 Interfacing to non-GNAT Ada code .................. 275

14 Specialized Needs Annexes ................. 276
15 Implementation of Specific Ada Features 

15.1 Machine Code Insertions .................................. 277
15.2 GNAT Implementation of Tasking ...................... 279
  15.2.1 Mapping Ada Tasks onto the Underlying Kernel Threads
  ................................................................. 279
  15.2.2 Ensuring Compliance with the Real-Time Annex .... 280
  15.2.3 Support for Locking Policies .......................... 280
15.3 GNAT Implementation of Shared Passive Packages .... 281
15.4 Code Generation for Array Aggregates ................. 282
  15.4.1 Static constant aggregates with static bounds .... 282
  15.4.2 Constant aggregates with unconstrained nominal types
  ................................................................. 283
  15.4.3 Aggregates with static bounds ...................... 283
  15.4.4 Aggregates with nonstatic bounds .................. 283
  15.4.5 Aggregates in assignment statements .............. 283
15.5 The Size of Discriminated Records with Default Discriminants
  ........................................................................ 284
15.6 Image Values For Nonscalar Types ................. 285
15.7 Strict Conformance to the Ada Reference Manual .... 285

16 Implementation of Ada 2012 Features .... 287

17 Security Hardening Features ............... 304
  17.1 Register Scrubbing ........................................ 304
  17.2 Stack Scrubbing ............................................ 304
  17.3 Hardened Conditionals .................................. 306
  17.4 Hardened Booleans ....................................... 308
  17.5 Control Flow Redundancy ............................... 308

18 Obsolescent Features ......................... 312
  18.1 pragma No_Run_Time ...................................... 312
  18.2 pragma Ravenscar ......................................... 312
  18.3 pragma Restricted_Run_Time ............................ 312
  18.4 pragma Task_Info ......................................... 312
  18.5 package System.Task_Info (s-tasinf.ads) .......... 312

19 Compatibility and Porting Guide ............ 313
  19.1 Writing Portable Fixed-Point Declarations ............ 313
  19.2 Compatibility with Ada 83 ............................. 314
    19.2.1 Legal Ada 83 programs that are illegal in Ada 95
    ............................................................... 314
    19.2.2 More deterministic semantics .................... 316
    19.2.3 Changed semantics .................................. 316
    19.2.4 Other language compatibility issues ............. 316
  19.3 Compatibility between Ada 95 and Ada 2005 .......... 317
  19.4 Implementation-dependent characteristics .......... 318
    19.4.1 Implementation-defined pragmas .................. 318
19.4.2 Implementation-defined attributes .................................. 318
19.4.3 Libraries ........................................................................... 318
19.4.4 Elaboration order ................................................................. 318
19.4.5 Target-specific aspects ......................................................... 319
19.5 Compatibility with Other Ada Systems ................................. 319
19.6 Representation Clauses ......................................................... 320
19.7 Compatibility with HP Ada 83 ............................................... 321

20 GNU Free Documentation License ................................. 322

Index ......................................................................................... 329
‘GNAT, The GNU Ada Development Environment’
GCC version 13.0.0
AdaCore
Permission is granted to copy, distribute and/or modify this document under the terms of
the GNU Free Documentation License, Version 1.3 or any later version published by the Free
Software Foundation; with no Invariant Sections, with the Front-Cover Texts being “GNAT
Reference Manual”, and with no Back-Cover Texts. A copy of the license is included in the
section entitled [GNU Free Documentation License], page 321.
1 About This Guide

This manual contains useful information in writing programs using the GNAT compiler. It includes information on implementation dependent characteristics of GNAT, including all the information required by Annex M of the Ada language standard.

GNAT implements Ada 95, Ada 2005 and Ada 2012, and it may also be invoked in Ada 83 compatibility mode. By default, GNAT assumes Ada 2012, but you can override with a compiler switch to explicitly specify the language version. (Please refer to the ‘GNAT User’s Guide’ for details on these switches.) Throughout this manual, references to ‘Ada’ without a year suffix apply to all the Ada versions of the language.

Ada is designed to be highly portable. In general, a program will have the same effect even when compiled by different compilers on different platforms. However, since Ada is designed to be used in a wide variety of applications, it also contains a number of system dependent features to be used in interfacing to the external world.

Note: Any program that makes use of implementation-dependent features may be non-portable. You should follow good programming practice and isolate and clearly document any sections of your program that make use of these features in a non-portable manner.

1.1 What This Reference Manual Contains

This reference manual contains the following chapters:

* [Implementation Defined Pragmas], page 4, lists GNAT implementation-dependent pragmas, which can be used to extend and enhance the functionality of the compiler.
* [Implementation Defined Attributes], page 115, lists GNAT implementation-dependent attributes, which can be used to extend and enhance the functionality of the compiler.
* [Standard and Implementation Defined Restrictions], page 138, lists GNAT implementation-dependent restrictions, which can be used to extend and enhance the functionality of the compiler.
* [Implementation Advice], page 152, provides information on generally desirable behavior which are not requirements that all compilers must follow since it cannot be provided on all systems, or which may be undesirable on some systems.
* [Implementation Defined Characteristics], page 172, provides a guide to minimizing implementation dependent features.
* [Intrinsic Subprograms], page 191, describes the intrinsic subprograms implemented by GNAT, and how they can be imported into user application programs.
* [Representation Clauses and Pragmas], page 194, describes in detail the way that GNAT represents data, and in particular the exact set of representation clauses and pragmas that is accepted.
* [Standard Library Routines], page 225, provides a listing of packages and a brief description of the functionality that is provided by Ada’s extensive set of standard library routines as implemented by GNAT.
* [The Implementation of Standard I/O], page 236, details how the GNAT implementation of the input-output facilities.
* [The GNAT Library], page 254, is a catalog of packages that complement the Ada predefined library.
Chapter 1: About This Guide

* [Interfacing to Other Languages], page 272, describes how programs written in Ada using GNAT can be interfaced to other programming languages.
* [Specialized Needs Annexes], page 275, describes the GNAT implementation of all of the specialized needs annexes.
* [Implementation of Specific Ada Features], page 276, discusses issues related to GNAT’s implementation of machine code insertions, tasking, and several other features.
* [Implementation of Ada 2012 Features], page 286, describes the status of the GNAT implementation of the Ada 2012 language standard.
* [Security Hardening Features], page 303 documents GNAT extensions aimed at security hardening.
* [Obsolescent Features], page 311 documents implementation dependent features, including pragmas and attributes, which are considered obsolescent, since there are other preferred ways of achieving the same results. These obsolescent forms are retained for backwards compatibility.
* [Compatibility and Porting Guide], page 312 presents some guidelines for developing portable Ada code, describes the compatibility issues that may arise between GNAT and other Ada compilation systems (including those for Ada 83), and shows how GNAT can expedite porting applications developed in other Ada environments.
* [GNU Free Documentation License], page 321 contains the license for this document.

This reference manual assumes a basic familiarity with the Ada 95 language, as described in the International Standard ANSI/ISO/IEC-8652:1995. It does not require knowledge of the new features introduced by Ada 2005 or Ada 2012. All three reference manuals are included in the GNAT documentation package.

1.2 Conventions

Following are examples of the typographical and graphic conventions used in this guide:
* Functions, utility program names, standard names, and classes.
* Option flags
* File names
* Variables
* ‘Emphasis’
* [optional information or parameters]
* Examples are described by text
  and then shown this way.
* Commands that are entered by the user are shown as preceded by a prompt string comprising the $ character followed by a space.

1.3 Related Information

See the following documents for further information on GNAT:
* GNAT User’s Guide for Native Platforms, which provides information on how to use the GNAT development environment.
* Ada 95 Reference Manual, the Ada 95 programming language standard.
* Ada 95 Annotated Reference Manual, which is an annotated version of the Ada 95 standard. The annotations describe detailed aspects of the design decision, and in particular contain useful sections on Ada 83 compatibility.
* Ada 2005 Annotated Reference Manual, which is an annotated version of the Ada 2005 standard. The annotations describe detailed aspects of the design decision.
* DEC Ada, Technical Overview and Comparison on DIGITAL Platforms, which contains specific information on compatibility between GNAT and DEC Ada 83 systems.
* DEC Ada, Language Reference Manual, part number AA-PYZAB-TK, which describes in detail the pragmas and attributes provided by the DEC Ada 83 compiler system.
2 Implementation Defined Pragmas

Ada defines a set of pragmas that can be used to supply additional information to the compiler. These language defined pragmas are implemented in GNAT and work as described in the Ada Reference Manual.

In addition, Ada allows implementations to define additional pragmas whose meaning is defined by the implementation. GNAT provides a number of these implementation-defined pragmas, which can be used to extend and enhance the functionality of the compiler. This section of the GNAT Reference Manual describes these additional pragmas.

Note that any program using these pragmas might not be portable to other compilers (although GNAT implements this set of pragmas on all platforms). Therefore if portability to other compilers is an important consideration, the use of these pragmas should be minimized.

2.1 Pragma Abort_Defer

Syntax:

```
pragma Abort_Defer;
```

This pragma must appear at the start of the statement sequence of a handled sequence of statements (right after the `begin`). It has the effect of deferring aborts for the sequence of statements (but not for the declarations or handlers, if any, associated with this statement sequence). This can also be useful for adding a polling point in Ada code, where asynchronous abort of tasks is checked when leaving the statement sequence, and is lighter than, for example, using `delay 0.0;`, since with zero-cost exception handling, propagating exceptions (implicitly used to implement task abort) cannot be done reliably in an asynchronous way.

An example of usage would be:

```
-- Add a polling point to check for task aborts
begin
  pragma Abort_Defer;
end;
```

2.2 Pragma Abstract_State

Syntax:

```
pragma Abstract_State (ABSTRACT_STATE_LIST);
```

```
ABSTRACT_STATE_LIST ::= 
  null ] STATE_NAME_WITH_OPTIONS 
  | (STATE_NAME_WITH_OPTIONS {, STATE_NAME_WITH_OPTIONS} )

STATE_NAME_WITH_OPTIONS ::= 
  STATE_NAME 
  | (STATE_NAME with OPTION_LIST)
```
OPTION_LIST ::= OPTION {, OPTION}

OPTION ::= SIMPLE_OPTION | NAME_VALUE_OPTION

SIMPLE_OPTION ::= Ghost | Synchronous

NAME_VALUE_OPTION ::= Part_Of => ABSTRACT_STATE | External [=> EXTERNAL_PROPERTY_LIST]

EXTERNAL_PROPERTY_LIST ::= EXTERNAL_PROPERTY | (EXTERNAL_PROPERTY {, EXTERNAL_PROPERTY} )


STATE_NAME ::= defining_identifier

ABSTRACT_STATE ::= name

For the semantics of this pragma, see the entry for aspect Abstract_State in the SPARK 2014 Reference Manual, section 7.1.4.

2.3 Pragma Ada_83

Syntax:
pragma Ada_83;

A configuration pragma that establishes Ada 83 mode for the unit to which it applies, regardless of the mode set by the command line switches. In Ada 83 mode, GNAT attempts to be as compatible with the syntax and semantics of Ada 83, as defined in the original Ada 83 Reference Manual as possible. In particular, the keywords added by Ada 95 and Ada 2005 are not recognized, optional package bodies are allowed, and generics may name types with unknown discriminants without using the (<> notation. In addition, some but not all of the additional restrictions of Ada 83 are enforced.

Ada 83 mode is intended for two purposes. Firstly, it allows existing Ada 83 code to be compiled and adapted to GNAT with less effort. Secondly, it aids in keeping code backwards compatible with Ada 83. However, there is no guarantee that code that is processed correctly by GNAT in Ada 83 mode will in fact compile and execute with an Ada 83 compiler, since GNAT does not enforce all the additional checks required by Ada 83.
2.4 Pragma Ada_95

Syntax:

```
pragma Ada_95;
```

A configuration pragma that establishes Ada 95 mode for the unit to which it applies, regardless of the mode set by the command line switches. This mode is set automatically for the Ada and System packages and their children, so you need not specify it in these contexts. This pragma is useful when writing a reusable component that itself uses Ada 95 features, but which is intended to be usable from either Ada 83 or Ada 95 programs.

2.5 Pragma Ada_05

Syntax:

```
pragma Ada_05;
pragma Ada_05 (local_NAME);
```

A configuration pragma that establishes Ada 2005 mode for the unit to which it applies, regardless of the mode set by the command line switches. This pragma is useful when writing a reusable component that itself uses Ada 2005 features, but which is intended to be usable from either Ada 83 or Ada 95 programs.

The one argument form (which is not a configuration pragma) is used for managing the transition from Ada 95 to Ada 2005 in the run-time library. If an entity is marked as Ada_2005 only, then referencing the entity in Ada_83 or Ada_95 mode will generate a warning. In addition, in Ada_83 or Ada_95 mode, a preference rule is established which does not choose such an entity unless it is unambiguously specified. This avoids extra subprograms marked this way from generating ambiguities in otherwise legal pre-Ada_2005 programs. The one argument form is intended for exclusive use in the GNAT run-time library.

2.6 Pragma Ada_2005

Syntax:

```
pragma Ada_2005;
```

This configuration pragma is a synonym for pragma Ada_05 and has the same syntax and effect.

2.7 Pragma Ada_12

Syntax:

```
pragma Ada_12;
pragma Ada_12 (local_NAME);
```

A configuration pragma that establishes Ada 2012 mode for the unit to which it applies, regardless of the mode set by the command line switches. This mode is set automatically for the Ada and System packages and their children, so you need not specify it in these contexts. This pragma is useful when writing a reusable component that itself uses Ada 2012 features, but which is intended to be usable from Ada 83, Ada 95, or Ada 2005 programs.

The one argument form, which is not a configuration pragma, is used for managing the transition from Ada 2005 to Ada 2012 in the run-time library. If an entity is marked
as Ada_2012 only, then referencing the entity in any pre-Ada_2012 mode will generate a warning. In addition, in any pre-Ada_2012 mode, a preference rule is established which does not choose such an entity unless it is unambiguously specified. This avoids extra subprograms marked this way from generating ambiguities in otherwise legal pre-Ada_2012 programs. The one argument form is intended for exclusive use in the GNAT run-time library.

2.8 Pragma Ada_2012
Syntax:
pragma Ada_2012;
This configuration pragma is a synonym for pragma Ada_12 and has the same syntax and effect.

2.9 Pragma Ada_2022
Syntax:
pragma Ada_2022;
pragma Ada_2022 (local_NAME);
A configuration pragma that establishes Ada 2022 mode for the unit to which it applies, regardless of the mode set by the command line switches. This mode is set automatically for the Ada and System packages and their children, so you need not specify it in these contexts. This pragma is useful when writing a reusable component that itself uses Ada 2022 features, but which is intended to be usable from Ada 83, Ada 95, Ada 2005 or Ada 2012 programs.
The one argument form, which is not a configuration pragma, is used for managing the transition from Ada 2012 to Ada 2022 in the run-time library. If an entity is marked as Ada_2022 only, then referencing the entity in any pre-Ada_2022 mode will generate a warning. In addition, in any pre-Ada_2012 mode, a preference rule is established which does not choose such an entity unless it is unambiguously specified. This avoids extra subprograms marked this way from generating ambiguities in otherwise legal pre-Ada_2022 programs. The one argument form is intended for exclusive use in the GNAT run-time library.

2.10 Pragma Aggregate_Individually_Assign
Syntax:
pragma Aggregate_Individually_Assign;
Where possible, GNAT will store the binary representation of a record aggregate in memory for space and performance reasons. This configuration pragma changes this behavior so that record aggregates are instead always converted into individual assignment statements.

2.11 Pragma Allow_Integer_Address
Syntax:
pragma Allow_Integer_Address;
In almost all versions of GNAT, `System.Address` is a private type in accordance with the implementation advice in the RM. This means that integer values, in particular integer literals, are not allowed as address values. If the configuration pragma `Allow_Integer_Address` is given, then integer expressions may be used anywhere a value of type `System.Address` is required. The effect is to introduce an implicit unchecked conversion from the integer value to type `System.Address`. The reverse case of using an address where an integer type is required is handled analogously. The following example compiles without errors:

```pascal
pragma Allow_Integer_Address;
with System; use System;
package AddrAsInt is
  X : Integer;
  Y : Integer;
  for X'Address use 16#1240#;
  for Y use at 16#3230#;
  m : Address := 16#4000#;
  n : constant Address := 4000;
  p : constant Address := Address (X + Y);
  v : Integer := Y'Address;
  w : constant Integer := Integer (Y'Address);
  type R is new integer;
  RR : R := 1000;
  Z : Integer;
  for Z'Address use RR;
end AddrAsInt;
```

Note that pragma `Allow_Integer_Address` is ignored if `System.Address` is not a private type. In implementations of GNAT where `System.Address` is a visible integer type, this pragma serves no purpose but is ignored rather than rejected to allow common sets of sources to be used in the two situations.

### 2.12 Pragma Annotate

Syntax:

```
pragma Annotate (IDENTIFIER [, IDENTIFIER {, ARG}] [, entity => local_NAME]);
```

ARG ::= NAME | EXPRESSION

This pragma is used to annotate programs. IDENTIFIER identifies the type of annotation. GNAT verifies that it is an identifier, but does not otherwise analyze it. The second optional identifier is also left unanalyzed, and by convention is used to control the action of the tool to which the annotation is addressed. The remaining ARG arguments can be either string literals or more generally expressions. String literals (and concatenations of string literals) are assumed to be either of type `Standard.String` or else `Wide_String` or `Wide_Wide_String` depending on the character literals they contain. All other kinds of arguments are analyzed as expressions, and must be unambiguous. The last argument if present must have the identifier `Entity` and GNAT verifies that a local name is given.

The analyzed pragma is retained in the tree, but not otherwise processed by any part of the GNAT compiler, except to generate corresponding note lines in the generated ALI file.
For the format of these note lines, see the compiler source file lib-writ.ads. This pragma is intended for use by external tools, including ASIS. The use of pragma Annotate does not affect the compilation process in any way. This pragma may be used as a configuration pragma.

### 2.13 Pragma Assert

**Syntax:**

```ada
pragma Assert (  
    boolean_EXPRESSION  
[, string_EXPRESSION])
```

The effect of this pragma depends on whether the corresponding command line switch is set to activate assertions. The pragma expands into code equivalent to the following:

```ada
if assertions-enabled then  
    if not boolean_EXPRESSION then  
        System.Assertions.Raise_Asert_Failure  
            (string_EXPRESSION);  
    end if;  
end if;
```

The string argument, if given, is the message that will be associated with the exception occurrence if the exception is raised. If no second argument is given, the default message is `file:nnn`, where `file` is the name of the source file containing the assert, and `nnn` is the line number of the assert.

Note that, as with the `if` statement to which it is equivalent, the type of the expression is either `Standard.Boolean`, or any type derived from this standard type.

Assert checks can be either checked or ignored. By default they are ignored. They will be checked if either the command line switch `-gnata` is used, or if an `Assertion_Policy` or `Check_Policy` pragma is used to enable `Assert_Checks`.

If assertions are ignored, then there is no run-time effect (and in particular, any side effects from the expression will not occur at run time). (The expression is still analyzed at compile time, and may cause types to be frozen if they are mentioned here for the first time).

If assertions are checked, then the given expression is tested, and if it is `False` then `System.Assertions.Raise_Asert_Failure` is called which results in the raising of `Assert_Failure` with the given message.

You should generally avoid side effects in the expression arguments of this pragma, because these side effects will turn on and off with the setting of the assertions mode, resulting in assertions that have an effect on the program. However, the expressions are analyzed for semantic correctness whether or not assertions are enabled, so turning assertions on and off cannot affect the legality of a program.

Note that the implementation defined policy `DISABLE`, given in a pragma `Assertion_Policy`, can be used to suppress this semantic analysis.

Note: this is a standard language-defined pragma in versions of Ada from 2005 on. In GNAT, it is implemented in all versions of Ada, and the DISABLE policy is an implementation-defined addition.
2.14 Pragma Assert_And_Cut

Syntax:

```plaintext
pragma Assert_And_Cut (boolean_EXPRESSION [, string_EXPRESSION]);
```

The effect of this pragma is identical to that of pragma Assert, except that in an Assertion_Policy pragma, the identifier Assert_And_Cut is used to control whether it is ignored or checked (or disabled).

The intention is that this be used within a subprogram when the given test expression sums up all the work done so far in the subprogram, so that the rest of the subprogram can be verified (informally or formally) using only the entry preconditions, and the expression in this pragma. This allows dividing up a subprogram into sections for the purposes of testing or formal verification. The pragma also serves as useful documentation.

2.15 Pragma Assertion_Policy

Syntax:

```plaintext
pragma Assertion_Policy (CHECK | DISABLE | IGNORE | SUPPRESSIBLE);
```

```plaintext
pragma Assertion_Policy ( 
    ASSERTION_KIND => POLICY_IDENTIFIER 
    {, ASSERTION_KIND => POLICY_IDENTIFIER});
```

```plaintext
ASSERTION_KIND ::= RM_ASSERTION_KIND | ID_ASSERTION_KIND
```

```plaintext
RM_ASSERTION_KIND ::= Assert | Static_Predicate | Dynamic_Predicate | Pre | Pre’Class | Post | Post’Class | Type_Invariant | Type_Invariant’Class | Default_Initial.Condition
```

```plaintext
ID_ASSERTION_KIND ::= Assertions | Assert_And_Cut | Assume | Contract_Cases | Debug | Ghost | Initial.Condition | Invariant | Invariant’Class | Loop_Invariant
```
Chapter 2: Implementation Defined Pragmas

POLICY_IDENTIFIER ::= Check | Disable | Ignore | Suppressible

This is a standard Ada 2012 pragma that is available as an implementation-defined pragma in earlier versions of Ada. The assertion kinds RM_ASSERTION_KIND are those defined in the Ada standard. The assertion kinds ID_ASSERTION_KIND are implementation defined additions recognized by the GNAT compiler.

The pragma applies in both cases to pragmas and aspects with matching names, e.g. Pre applies to the Pre aspect, and Precondition applies to both the Precondition pragma and the aspect Precondition. Note that the identifiers for pragmas Pre'Class and Post'Class are Pre'Class and Post'Class (not Pre.Class and Post.Class), since these pragmas are intended to be identical to the corresponding aspects.

If the policy is CHECK, then assertions are enabled, i.e. the corresponding pragma or aspect is activated. If the policy is IGNORE, then assertions are ignored, i.e. the corresponding pragma or aspect is deactivated. This pragma overrides the effect of the ‘-gnata’ switch on the command line. If the policy is SUPPRESSIBLE, then assertions are enabled by default, however, if the ‘-gnatp’ switch is specified all assertions are ignored.

The implementation defined policy DISABLE is like IGNORE except that it completely disables semantic checking of the corresponding pragma or aspect. This is useful when the pragma or aspect argument references subprograms in a with’ed package which is replaced by a dummy package for the final build.

The implementation defined assertion kind Assertions applies to all assertion kinds. The form with no assertion kind given implies this choice, so it applies to all assertion kinds (RM defined, and implementation defined).

The implementation defined assertion kind StatementAssertions applies to Assert, Assert_And_Cut, Assume, Loop_Invariant, and Loop_Variant.

2.16 Pragma Assume

Syntax:

    pragma Assume (  
        boolean_EXPRESSION  
        [, string_EXPRESSION])

The effect of this pragma is identical to that of pragma Assert, except that in an Assertion_Policy pragma, the identifier Assume is used to control whether it is ignored or checked (or disabled).

The intention is that this be used for assumptions about the external environment. So you cannot expect to verify formally or informally that the condition is met, this must be established by examining things outside the program itself. For example, we may have code that depends on the size of Long_Long_Integer being at least 64. So we could write:
pragma Assume (Long_Long_Integer’Size >= 64);

This assumption cannot be proved from the program itself, but it acts as a useful run-time check that the assumption is met, and documents the need to ensure that it is met by reference to information outside the program.

### 2.17 Pragma Assume_No_Invalid_Values

**Syntax:**

```
pragma Assume_No_Invalid_Values (On | Off);
```

This is a configuration pragma that controls the assumptions made by the compiler about the occurrence of invalid representations (invalid values) in the code.

The default behavior (corresponding to an Off argument for this pragma), is to assume that values may in general be invalid unless the compiler can prove they are valid. Consider the following example:

```pascal
V1 : Integer range 1 .. 10;
V2 : Integer range 11 .. 20;
...
for J in V2 .. V1 loop
  ...
end loop;
```

if V1 and V2 have valid values, then the loop is known at compile time not to execute since the lower bound must be greater than the upper bound. However in default mode, no such assumption is made, and the loop may execute. If `Assume_No_Invalid_Values (On)` is given, the compiler will assume that any occurrence of a variable other than in an explicit ‘Valid test always has a valid value, and the loop above will be optimized away.

The use of `Assume_No_Invalid_Values (On)` is appropriate if you know your code is free of uninitialized variables and other possible sources of invalid representations, and may result in more efficient code. A program that accesses an invalid representation with this pragma in effect is erroneous, so no guarantees can be made about its behavior.

It is peculiar though permissible to use this pragma in conjunction with validity checking (`-gnatVa`). In such cases, accessing invalid values will generally give an exception, though formally the program is erroneous so there are no guarantees that this will always be the case, and it is recommended that these two options not be used together.

### 2.18 Pragma Async_Readers

**Syntax:**

```
pragma Async_Readers [ (static_boolean_EXPRESSION) ];
```

For the semantics of this pragma, see the entry for aspect `Async_Readers` in the SPARK 2014 Reference Manual, section 7.1.2.

### 2.19 Pragma Async_Writers

**Syntax:**

```
pragma Async_Writers [ (static_boolean_EXPRESSION) ];
```
Chapter 2: Implementation Defined Pragmas

For the semantics of this pragma, see the entry for aspect Async_Writers in the SPARK 2014 Reference Manual, section 7.1.2.

2.20 Pragma Attribute_Definition

Syntax:

```
pragma Attribute_Definition
  ([Attribute =>] ATTRIBUTE_DESIGNATOR,
   [Entity =>] LOCAL_NAME,
   [Expression =>] EXPRESSION | NAME);
```

If \texttt{Attribute} is a known attribute name, this pragma is equivalent to the attribute definition clause:

```
for Entity'Attribute use Expression;
```

If \texttt{Attribute} is not a recognized attribute name, the pragma is ignored, and a warning is emitted. This allows source code to be written that takes advantage of some new attribute, while remaining compilable with earlier compilers.

2.21 Pragma C_Pass_By_Copy

Syntax:

```
pragma C_Pass_By_Copy
  ([Max_Size =>] static_integer_EXPRESSION);
```

Normally the default mechanism for passing C convention records to C convention subprograms is to pass them by reference, as suggested by RM B.3(69). Use the configuration pragma \texttt{C_Pass_By_Copy} to change this default, by requiring that record formal parameters be passed by copy if all of the following conditions are met:

* The size of the record type does not exceed the value specified for \texttt{Max_Size}.
* The record type has \texttt{Convention C}.
* The formal parameter has this record type, and the subprogram has a foreign (non-Ada) convention.

If these conditions are met the argument is passed by copy; i.e., in a manner consistent with what C expects if the corresponding formal in the C prototype is a struct (rather than a pointer to a struct).

You can also pass records by copy by specifying the convention \texttt{C_Pass_By_Copy} for the record type, or by using the extended \texttt{Import} and \texttt{Export} pragmas, which allow specification of passing mechanisms on a parameter by parameter basis.

2.22 Pragma Check

Syntax:

```
pragma Check (
  [Name =>] CHECK_KIND,
  [Check =>] Boolean_EXPRESSION
  [, [Message =>] string_EXPRESSION] );
```
This pragma is similar to the predefined pragma `Assert` except that an extra identifier argument is present. In conjunction with pragma `Check_Policy`, this can be used to define groups of assertions that can be independently controlled. The identifier `Assertion` is special, it refers to the normal set of pragma `Assert` statements.

Checks introduced by this pragma are normally deactivated by default. They can be activated either by the command line option `-gnata`, which turns on all checks, or individually controlled using pragma `Check_Policy`.

The identifiers `Assertions` and `Statement_AAssertions` are not permitted as check kinds, since this would cause confusion with the use of these identifiers in `Assertion_Policy` and `Check_Policy` pragmas, where they are used to refer to sets of assertions.

### 2.23 Pragma Check_Float_Overflow

Syntax:

```
pragma Check_Float_Overflow;
```

In Ada, the predefined floating-point types (`Short_Float`, `Float`, `Long_Float`, `Long_Long_Float`) are defined to be ‘unconstrained’. This means that even though each has a well-defined base range, an operation that delivers a result outside this base range is not required to raise an exception. This implementation permission accommodates the notion of infinities in IEEE floating-point, and corresponds to the efficient execution mode on most machines. GNAT will not raise overflow exceptions on these machines; instead it will generate infinities and NaN’s as defined in the IEEE standard.

Generating infinities, although efficient, is not always desirable. Often the preferable approach is to check for overflow, even at the (perhaps considerable) expense of run-time performance. This can be accomplished by defining your own constrained floating-point subtypes – i.e., by supplying explicit range constraints – and indeed such a subtype can have the same base range as its base type. For example:

```
subtype My_Float is Float range Float’Range;
```

Here `My_Float` has the same range as `Float` but is constrained, so operations on `My_Float` values will be checked for overflow against this range.

This style will achieve the desired goal, but it is often more convenient to be able to simply use the standard predefined floating-point types as long as overflow checking could be guaranteed. The `Check_Float_Overflow` configuration pragma achieves this effect. If a unit is compiled subject to this configuration pragma, then all operations on predefined floating-point types including operations on base types of these floating-point types will be treated as though those types were constrained, and overflow checks will be generated. The `Constraint_Error` exception is raised if the result is out of range.

This mode can also be set by use of the compiler switch ‘-gnateF’.
2.24 Pragma Check_Name

Syntax:

    pragma Check_Name (check_name_IDENTIFIER);

This is a configuration pragma that defines a new implementation defined check name (unless IDENTIFIER matches one of the predefined check names, in which case the pragma has no effect). Check names are global to a partition, so if two or more configuration pragmas are present in a partition mentioning the same name, only one new check name is introduced.

An implementation defined check name introduced with this pragma may be used in only three contexts: \texttt{pragma Suppress}, \texttt{pragma Unsuppress}, and as the prefix of a \texttt{Check\_Name'Enabled} attribute reference. For any of these three cases, the check name must be visible. A check name is visible if it is in the configuration pragmas applying to the current unit, or if it appears at the start of any unit that is part of the dependency set of the current unit (e.g., units that are mentioned in \texttt{with} clauses).

Check names introduced by this pragma are subject to control by compiler switches (in particular \texttt{-gnatp}) in the usual manner.

2.25 Pragma Check_Policy

Syntax:

    pragma Check_Policy
    ([Name =>] CHECK_KIND,
     [Policy =>] POLICY_IDENTIFIER);

    pragma Check_Policy (
        CHECK_KIND => POLICY_IDENTIFIER
    , CHECK_KIND => POLICY_IDENTIFIER));

    ASSERTION_KIND ::= RM_ASSERTION_KIND | ID_ASSERTION_KIND

    CHECK_KIND ::= IDENTIFIER | Pre'Class | Post'Class | Type_Invariant'Class | Invariant'Class

The identifiers Name and Policy are not allowed as CHECK_KIND values. This avoids confusion between the two possible syntax forms for this pragma.

    POLICY_IDENTIFIER ::= ON | OFF | CHECK | DISABLE | IGNORE

This pragma is used to set the checking policy for assertions (specified by aspects or pragmas), the \texttt{Debug} pragma, or additional checks to be checked using the \texttt{Check} pragma. It may appear either as a configuration pragma, or within a declarative part of package. In the latter case, it applies from the point where it appears to the end of the declarative region (like \texttt{pragma Suppress}).
The `Check_Policy` pragma is similar to the predefined `Assertion_Policy` pragma, and if the check kind corresponds to one of the assertion kinds that are allowed by `Assertion_Policy`, then the effect is identical.

If the first argument is `Debug`, then the policy applies to `Debug` pragmas, disabling their effect if the policy is `OFF`, `DISABLE`, or `IGNORE`, and allowing them to execute with normal semantics if the policy is `ON` or `CHECK`. In addition if the policy is `DISABLE`, then the procedure call in `Debug` pragmas will be totally ignored and not analyzed semantically.

Finally the first argument may be some other identifier than the above possibilities, in which case it controls a set of named assertions that can be checked using pragma `Check`. For example, if the pragma:

```
pragma Check_Policy (Critical_Error, OFF);
```

is given, then subsequent `Check` pragmas whose first argument is also `Critical_Error` will be disabled.

The check policy is `OFF` to turn off corresponding checks, and `ON` to turn on corresponding checks. The default for a set of checks for which no `Check_Policy` is given is `OFF` unless the compiler switch `-gnata` is given, which turns on all checks by default.

The check policy settings `CHECK` and `IGNORE` are recognized as synonyms for `ON` and `OFF`. These synonyms are provided for compatibility with the standard `Assertion_Policy` pragma. The check policy setting `DISABLE` causes the second argument of a corresponding `Check` pragma to be completely ignored and not analyzed.

### 2.26 Pragma Comment

Syntax:

```
pragma Comment (static_string_EXPRESSION);
```

This is almost identical in effect to pragma `Ident`. It allows the placement of a comment into the object file and hence into the executable file if the operating system permits such usage. The difference is that `Comment`, unlike `Ident`, has no limitations on placement of the pragma (it can be placed anywhere in the main source unit), and if more than one pragma is used, all comments are retained.

### 2.27 Pragma Common_Object

Syntax:

```
pragma Common_Object (   
    [Internal =>] LOCAL_NAME 
    [, [External =>] EXTERNAL_SYMBOL] 
    [, [Size =>] EXTERNAL_SYMBOL] );
```

```
EXTERNAL_SYMBOL ::=  
  IDENTIFIER 
  | static_string_EXPRESSION
```

This pragma enables the shared use of variables stored in overlaid linker areas corresponding to the use of `COMMON` in Fortran. The single object `LOCAL_NAME` is assigned to the area designated by the `External` argument. You may define a record to correspond to a series of fields. The `Size` argument is syntax checked in GNAT, but otherwise ignored.
Common_Object is not supported on all platforms. If no support is available, then the code generator will issue a message indicating that the necessary attribute for implementation of this pragma is not available.

2.28 Pragma Compile_Time_Error

Syntax:

```plaintext
pragma Compile_Time_Error
  (boolean_EXPRESSION, static_string_EXPRESSION);
```

This pragma can be used to generate additional compile time error messages. It is particularly useful in generics, where errors can be issued for specific problematic instantiations. The first parameter is a boolean expression. The pragma ensures that the value of an expression is known at compile time, and has the value False. The set of expressions whose values are known at compile time includes all static boolean expressions, and also other values which the compiler can determine at compile time (e.g., the size of a record type set by an explicit size representation clause, or the value of a variable which was initialized to a constant and is known not to have been modified). If these conditions are not met, an error message is generated using the value given as the second argument. This string value may contain embedded ASCII.LF characters to break the message into multiple lines.

2.29 Pragma Compile_Time_Warning

Syntax:

```plaintext
pragma Compile_Time_Warning
  (boolean_EXPRESSION, static_string_EXPRESSION);
```

Same as pragma Compile_Time_Error, except a warning is issued instead of an error message. If switch `-gnatw_C` is used, a warning is only issued if the value of the expression is known to be True at compile time, not when the value of the expression is not known at compile time. Note that if this pragma is used in a package that is with’ed by a client, the client will get the warning even though it is issued by a with’ed package (normally warnings in with’ed units are suppressed, but this is a special exception to that rule).

One typical use is within a generic where compile time known characteristics of formal parameters are tested, and warnings given appropriately. Another use with a first parameter of True is to warn a client about use of a package, for example that it is not fully implemented.

In previous versions of the compiler, combining `-gnatwe` with Compile_Time_Warning resulted in a fatal error. Now the compiler always emits a warning. You can use [Pragma Compile_Time_Error], page 18 to force the generation of an error.

2.30 Pragma Complete_Representation

Syntax:

```plaintext
pragma Complete_Representation;
```

This pragma must appear immediately within a record representation clause. Typical placements are before the first component clause or after the last component clause. The effect is to give an error message if any component is missing a component clause. This pragma may
be used to ensure that a record representation clause is complete, and that this invariant is maintained if fields are added to the record in the future.

### 2.31 Pragma Complex_Representation

**Syntax:**

```plaintext
pragma Complex_Representation
  ([Entity =>] LOCAL_NAME);
```

The `Entity` argument must be the name of a record type which has two fields of the same floating-point type. The effect of this pragma is to force gcc to use the special internal complex representation form for this record, which may be more efficient. Note that this may result in the code for this type not conforming to standard ABI (application binary interface) requirements for the handling of record types. For example, in some environments, there is a requirement for passing records by pointer, and the use of this pragma may result in passing this type in floating-point registers.

### 2.32 Pragma Component_Alignment

**Syntax:**

```plaintext
pragma Component_Alignment (   
  [Form =>] ALIGNMENT_CHOICE  
  [, [Name =>] type_LOCAL_NAME]);
```

**ALIGNMENT_CHOICE ::=**

- `Component_Size`
- `Component_Size_4`
- `Storage_Unit`
- `Default`

Specifies the alignment of components in array or record types. The meaning of the `Form` argument is as follows:

- **‘Component_Size’**
  Aligns scalar components and subcomponents of the array or record type on boundaries appropriate to their inherent size (naturally aligned). For example, 1-byte components are aligned on byte boundaries, 2-byte integer components are aligned on 2-byte boundaries, 4-byte integer components are aligned on 4-byte boundaries and so on. These alignment rules correspond to the normal rules for C compilers on all machines except the VAX.

- **‘Component_Size_4’**
  Naturally aligns components with a size of four or fewer bytes. Components that are larger than 4 bytes are placed on the next 4-byte boundary.

- **‘Storage_Unit’**
  Specifies that array or record components are byte aligned, i.e., aligned on boundaries determined by the value of the constant `System.Storage_Unit`.

- **‘Default’**
Specifies that array or record components are aligned on default boundaries, appropriate to the underlying hardware or operating system or both. The Default choice is the same as Component_Size (natural alignment).

If the Name parameter is present, type_LOCAL_NAME must refer to a local record or array type, and the specified alignment choice applies to the specified type. The use of Component_Alignment together with a pragma Pack causes the Component_Alignment pragma to be ignored. The use of Component_Alignment together with a record representation clause is only effective for fields not specified by the representation clause.

If the Name parameter is absent, the pragma can be used as either a configuration pragma, in which case it applies to one or more units in accordance with the normal rules for configuration pragmas, or it can be used within a declarative part, in which case it applies to types that are declared within this declarative part, or within any nested scope within this declarative part. In either case it specifies the alignment to be applied to any record or array type which has otherwise standard representation.

If the alignment for a record or array type is not specified (using pragma Pack, pragma Component_Alignment, or a record rep clause), the GNAT uses the default alignment as described previously.

2.33 Pragma Constant_After_Elaboration

Syntax:

```
pragma Constant_After_Elaboration [ (static_boolean_EXPRESSION) ];
```

For the semantics of this pragma, see the entry for aspect Constant_After_Elaboration in the SPARK 2014 Reference Manual, section 3.3.1.

2.34 Pragma Contract_Cases

Syntax:

```
pragma Contract_Cases (CONTRACT_CASE {, CONTRACT_CASE});
```

```
CONTRACT_CASE ::= CASE_GUARD => CONSEQUENCE
CASE_GUARD ::= boolean_EXPRESSION | others
CONSEQUENCE ::= boolean_EXPRESSION
```

The Contract_Cases pragma allows defining fine-grain specifications that can complement or replace the contract given by a precondition and a postcondition. Additionally, the Contract_Cases pragma can be used by testing and formal verification tools. The compiler checks its validity and, depending on the assertion policy at the point of declaration of the pragma, it may insert a check in the executable. For code generation, the contract cases

```
pragma Contract_Cases (    
    Cond1 => Pred1,    
    Cond2 => Pred2);  
```

are equivalent to

```
C1 : constant Boolean := Cond1; -- evaluated at subprogram entry
```
C2 : constant Boolean := Cond2; -- evaluated at subprogram entry
pragma Precondition (((C1 and not C2) or (C2 and not C1)));
pragma Postcondition (if C1 then Pred1);
pragma Postcondition (if C2 then Pred2);

The precondition ensures that one and only one of the case guards is satisfied on entry to the subprogram. The postcondition ensures that for the case guard that was True on entry, the corresponding consequence is True on exit. Other consequence expressions are not evaluated.

A precondition \(P\) and postcondition \(Q\) can also be expressed as contract cases:

**pragma Contract_Cases \((P \Rightarrow Q)\);**

The placement and visibility rules for **Contract_Cases** pragmas are identical to those described for preconditions and postconditions.

The compiler checks that boolean expressions given in case guards and consequences are valid, where the rules for case guards are the same as the rule for an expression in **Precondition** and the rules for consequences are the same as the rule for an expression in **Postcondition**. In particular, attributes 'Old' and 'Result' can only be used within consequence expressions. The case guard for the last contract case may be **others**, to denote any case not captured by the previous cases. The following is an example of use within a package spec:

```pascal
package Math_Functions is
...
  function Sqrt (Arg : Float) return Float;
  pragma Contract_Cases (((Arg in 0.0 .. 99.0) => Sqrt'Result < 10.0,
                           Arg >= 100.0       => Sqrt'Result >= 10.0,
                           others            => Sqrt'Result = 0.0));
...
end Math_Functions;
```

The meaning of contract cases is that only one case should apply at each call, as determined by the corresponding case guard evaluating to True, and that the consequence for this case should hold when the subprogram returns.

### 2.35 Pragma Convention_Identifier

Syntax:

**pragma Convention_Identifier (**

\[[Name =>]\] IDENTIFIER,
\[[Convention =>]\] convention_IDENTIFIER**);

This pragma provides a mechanism for supplying synonyms for existing convention identifiers. The \(Name\) identifier can subsequently be used as a synonym for the given convention in other pragmas (including for example pragma **Import** or another **Convention_Identifier** pragma). As an example of the use of this, suppose you had legacy code which used Fortran77 as the identifier for Fortran. Then the pragma:

**pragma Convention_Identifier (Fortran77, Fortran);**

would allow the use of the convention identifier **Fortran77** in subsequent code, avoiding the need to modify the sources. As another example, you could use this to parameterize
convention requirements according to systems. Suppose you needed to use Stdcall on windows systems, and C on some other system, then you could define a convention identifier Library and use a single Convention_Identifier pragma to specify which convention would be used system-wide.

2.36 Pragma CPP_Class

Syntax:

\[
\text{pragma CPP\_Class ([Entity =>] LOCAL\_NAME);}
\]

The argument denotes an entity in the current declarative region that is declared as a record type. It indicates that the type corresponds to an externally declared C++ class type, and is to be laid out the same way that C++ would lay out the type. If the C++ class has virtual primitives then the record must be declared as a tagged record type.

Types for which CPP_Class is specified do not have assignment or equality operators defined (such operations can be imported or declared as subprograms as required). Initialization is allowed only by constructor functions (see pragma CPP_Constructor). Such types are implicitly limited if not explicitly declared as limited or derived from a limited type, and an error is issued in that case.

See [Interfacing to C++], page 274 for related information.

Note: Pragma CPP_Class is currently obsolete. It is supported for backward compatibility but its functionality is available using pragma Import with Convention = CPP.

2.37 Pragma CPP_Constructor

Syntax:

\[
\text{pragma CPP\_Constructor ([Entity =>] LOCAL\_NAME}
\]

\[
[, [External\_Name =>] static\_string\_EXPRESSION ]
\]

\[
[, [Link\_Name =>] static\_string\_EXPRESSION ]);
\]

This pragma identifies an imported function (imported in the usual way with pragma Import) as corresponding to a C++ constructor. If External_Name and Link_Name are not specified then the Entity argument is a name that must have been previously mentioned in a pragma Import with Convention = CPP. Such name must be of one of the following forms:

* ‘function’ Fname ‘return’ T'
* ‘function’ Fname ‘return’ T'Class
* ‘function’ Fname (... ) ‘return’ T'
* ‘function’ Fname (... ) ‘return’ T'Class

where T is a limited record type imported from C++ with pragma Import and Convention = CPP.

The first two forms import the default constructor, used when an object of type T is created on the Ada side with no explicit constructor. The latter two forms cover all the non-default constructors of the type. See the GNAT User’s Guide for details.

If no constructors are imported, it is impossible to create any objects on the Ada side and the type is implicitly declared abstract.
Pragma **CPP_Constructor** is intended primarily for automatic generation using an automatic binding generator tool (such as the `-fdump-ada-spec` GCC switch). See [Interfacing to C++], page 274 for more related information.

Note: The use of functions returning class-wide types for constructors is currently obsolete. They are supported for backward compatibility. The use of functions returning the type T leave the Ada sources more clear because the imported C++ constructors always return an object of type T; that is, they never return an object whose type is a descendant of type T.

### 2.38 Pragma **CPP_Virtual**

This pragma is now obsolete and, other than generating a warning if warnings on obsolescent features are enabled, is completely ignored. It is retained for compatibility purposes. It used to be required to ensure compatibility with C++, but is no longer required for that purpose because GNAT generates the same object layout as the G++ compiler by default.

See [Interfacing to C++], page 274 for related information.

### 2.39 Pragma **CPP_Vtable**

This pragma is now obsolete and, other than generating a warning if warnings on obsolescent features are enabled, is completely ignored. It used to be required to ensure compatibility with C++, but is no longer required for that purpose because GNAT generates the same object layout as the G++ compiler by default.

See [Interfacing to C++], page 274 for related information.

### 2.40 Pragma **CPU**

Syntax:

```ada
pragma CPU (EXPRESSION);
```

This pragma is standard in Ada 2012, but is available in all earlier versions of Ada as an implementation-defined pragma. See Ada 2012 Reference Manual for details.

### 2.41 Pragma **Deadline_Floor**

Syntax:

```ada
pragma Deadline_Floor (time_span_EXPRESSION);
```

This pragma applies only to protected types and specifies the floor deadline inherited by a task when the task enters a protected object. It is effective only when the EDF scheduling policy is used.

### 2.42 Pragma **Debug**

Syntax:

```ada
pragma Debug ([CONDITION, ]PROCEDURE_CALL_WITHOUT_SEMICOLON);
```

```
    PROCEDURE_CALL WITHOUT SEMICOLON ::=
      PROCEDURE_NAME
      | PROCEDURE_PREFIX ACTUAL_PARAMETER_PART
```
The procedure call argument has the syntactic form of an expression, meeting the syntactic requirements for pragmas.

If debug pragmas are not enabled or if the condition is present and evaluates to False, this pragma has no effect. If debug pragmas are enabled, the semantics of the pragma is exactly equivalent to the procedure call statement corresponding to the argument with a terminating semicolon. Pragmas are permitted in sequences of declarations, so you can use pragma `Debug` to intersperse calls to debug procedures in the middle of declarations. Debug pragmas can be enabled either by use of the command line switch `-gnata` or by use of the pragma `Check_Policy` with a first argument of `Debug`.

### 2.43 Pragma Debug_Policy
Syntax:
```
pragma Debug_Policy (CHECK | DISABLE | IGNORE | ON | OFF);
```
This pragma is equivalent to a corresponding `Check_Policy` pragma with a first argument of `Debug`. It is retained for historical compatibility reasons.

### 2.44 Pragma Default_Initial_Condition
Syntax:
```
pragma Default_Initial_Condition [ (null | boolean_EXPRESSION) ];
```
For the semantics of this pragma, see the entry for aspect `Default_Initial_Condition` in the SPARK 2014 Reference Manual, section 7.3.3.

### 2.45 Pragma Default_Scalar_Storage_Order
Syntax:
```
pragma Default_Scalar_Storage_Order (High_Order_First | Low_Order_First);
```
Normally if no explicit `Scalar_Storage_Order` is given for a record type or array type, then the scalar storage order defaults to the ordinary default for the target. But this default may be overridden using this pragma. The pragma may appear as a configuration pragma, or locally within a package spec or declarative part. In the latter case, it applies to all subsequent types declared within that package spec or declarative part.

The following example shows the use of this pragma:
```
pragma Default_Scalar_Storage_Order (High_Order_First);
with System; use System;
package DSS01 is
  type H1 is record
    a : Integer;
  end record;

  type L2 is record
    a : Integer;
  end record;
  for L2'Scalar_Storage_Order use Low_Order_First;
```
type L2a is new L2;

package Inner is
    type H3 is record
        a : Integer;
    end record;

    pragma Default_Scalar_Storage_Order (Low_Order_First);

    type L4 is record
        a : Integer;
    end record;
end Inner;

type H4a is new Inner.L4;

type H5 is record
    a : Integer;
end record;
end DSSO1;

In this example record types with names starting with ‘L’ have Low_Order_First scalar storage order, and record types with names starting with ‘H’ have High_Order_First. Note that in the case of H4a, the order is not inherited from the parent type. Only an explicitly set Scalar_Storage_Order gets inherited on type derivation.

If this pragma is used as a configuration pragma which appears within a configuration pragma file (as opposed to appearing explicitly at the start of a single unit), then the binder will require that all units in a partition be compiled in a similar manner, other than run-time units, which are not affected by this pragma. Note that the use of this form is discouraged because it may significantly degrade the run-time performance of the software, instead the default scalar storage order ought to be changed only on a local basis.

2.46 Pragma Default_Storage_Pool

Syntax:

    pragma Default_Storage_Pool (storage_pool_NAME | null);

This pragma is standard in Ada 2012, but is available in all earlier versions of Ada as an implementation-defined pragma. See Ada 2012 Reference Manual for details.

2.47 Pragma Depends

Syntax:

    pragma Depends (DEPENDENCY_RELATION);


    DEPENDENCY_RELATION ::= 
        null 
        | (DEPENDENCY_CLAUSE {, DEPENDENCY_CLAUSE})
DEPENDENCY_CLAUSE ::= 
    OUTPUT_LIST =>[+] INPUT_LIST 
   | NULL_DEPENDENCY_CLAUSE

NULL_DEPENDENCY_CLAUSE ::= null => INPUT_LIST

OUTPUT_LIST ::= OUTPUT | (OUTPUT {, OUTPUT})

INPUT_LIST ::= null | INPUT | (INPUT {, INPUT})

OUTPUT ::= NAME | FUNCTION_RESULT

INPUT ::= NAME

where FUNCTION_RESULT is a function Result attribute_reference

For the semantics of this pragma, see the entry for aspect Depends in the SPARK 2014 Reference Manual, section 6.1.5.

2.48 Pragma Detect_Blocking

Syntax:

pragma Detect_Blocking;

This is a standard pragma in Ada 2005, that is available in all earlier versions of Ada as an implementation-defined pragma.

This is a configuration pragma that forces the detection of potentially blocking operations within a protected operation, and to raise Program_Error if that happens.

2.49 Pragma Disable_Atomic_Synchronization

Syntax:

pragma Disable_Atomic_Synchronization [(Entity)];

Ada requires that accesses (reads or writes) of an atomic variable be regarded as synchronization points in the case of multiple tasks. Particularly in the case of multi-processors this may require special handling, e.g. the generation of memory barriers. This capability may be turned off using this pragma in cases where it is known not to be required.

The placement and scope rules for this pragma are the same as those for pragma Suppress. In particular it can be used as a configuration pragma, or in a declaration sequence where it applies till the end of the scope. If an Entity argument is present, the action applies only to that entity.

2.50 Pragma Dispatching_Domain

Syntax:

pragma Dispatching_Domain (EXPRESSION);

This pragma is standard in Ada 2012, but is available in all earlier versions of Ada as an implementation-defined pragma. See Ada 2012 Reference Manual for details.
2.51 Pragma Effective_Reads

Syntax:

\[
\text{pragma Effective\_Reads [ (static\_boolean\_EXPRESSION) ]};
\]

For the semantics of this pragma, see the entry for aspect Effective\_Reads in the SPARK 2014 Reference Manual, section 7.1.2.

2.52 Pragma Effective_Writes

Syntax:

\[
\text{pragma Effective\_Writes [ (static\_boolean\_EXPRESSION) ]};
\]

For the semantics of this pragma, see the entry for aspect Effective\_Writes in the SPARK 2014 Reference Manual, section 7.1.2.

2.53 Pragma Elaboration_Checks

Syntax:

\[
\text{pragma Elaboration\_Checks (Dynamic | Static);}
\]

This is a configuration pragma which specifies the elaboration model to be used during compilation. For more information on the elaboration models of GNAT, consult the chapter on elaboration order handling in the ‘GNAT User’s Guide’.

The pragma may appear in the following contexts:

* Configuration pragmas file
* Prior to the context clauses of a compilation unit’s initial declaration

Any other placement of the pragma will result in a warning and the effects of the offending pragma will be ignored.

If the pragma argument is Dynamic, then the dynamic elaboration model is in effect. If the pragma argument is Static, then the static elaboration model is in effect.

2.54 Pragma Eliminate

Syntax:

\[
\text{pragma Eliminate (}
\begin{align*}
[ \text{Unit\_Name} & \Rightarrow \text{IDENTIFIER | SELECTED\_COMPONENT} , \\
[ \text{Entity} & \Rightarrow \text{IDENTIFIER |}
SELECTED\_COMPONENT | \\
& \text{STRING\_LITERAL} \\
[, \text{Source\_Location} & \Rightarrow \text{SOURCE\_TRACE} ] ) ;
\end{align*}
\]

\[
\text{SOURCE\_TRACE} ::= \text{STRING\_LITERAL}
\]

This pragma indicates that the given entity is not used in the program to be compiled and built, thus allowing the compiler to eliminate the code or data associated with the named entity. Any reference to an eliminated entity causes a compile-time or link-time error.

The pragma has the following semantics, where \( U \) is the unit specified by the Unit\_Name argument and \( E \) is the entity specified by the Entity argument:
* E must be a subprogram that is explicitly declared either:
  * Within U, or
  * Within a generic package that is instantiated in U, or
  * As an instance of a generic subprogram instantiated in U.

Otherwise the pragma is ignored.

* If E is overloaded within U then, in the absence of a Source_Location argument, all overloads are eliminated.
* If E is overloaded within U and only some overloads are to be eliminated, then each overloading to be eliminated must be specified in a corresponding pragma **Eliminate** with a Source_Location argument identifying the line where the declaration appears, as described below.
* If E is declared as the result of a generic instantiation, then a Source_Location argument is needed, as described below.

Pragma **Eliminate** allows a program to be compiled in a system-independent manner, so that unused entities are eliminated but without needing to modify the source text. Normally the required set of **Eliminate** pragmas is constructed automatically using the **gnatelim** tool.

Any source file change that removes, splits, or adds lines may make the set of **Eliminate** pragmas invalid because their Source_Location argument values may get out of date.

Pragma **Eliminate** may be used where the referenced entity is a dispatching operation. In this case all the subprograms to which the given operation can dispatch are considered to be unused (are never called as a result of a direct or a dispatching call).

The string literal given for the source location specifies the line number of the declaration of the entity, using the following syntax for **SOURCE_TRACE**:

```
SOURCE_TRACE ::= SOURCE_REFERENCE [ LBRACKET SOURCE_TRACE RBRACKET ]
LBRACKET ::= ']['
RBRACKET ::= ')[
SOURCE_REFERENCE ::= FILE_NAME : LINE_NUMBER
LINE_NUMBER ::= DIGIT {DIGIT}
```

Spaces around the colon in a **SOURCE_REFERENCE** are optional.

The source trace that is given as the Source_Location must obey the following rules (or else the pragma is ignored), where U is the unit U specified by the Unit_Name argument and E is the subprogram specified by the Entity argument:

* **FILE_NAME** is the short name (with no directory information) of the Ada source file for U, using the required syntax for the underlying file system (e.g. case is significant if the underlying operating system is case sensitive). If U is a package and E is a subprogram declared in the package specification and its full declaration appears in the package body, then the relevant source file is the one for the package specification; analogously if U is a generic package.
* If \( E \) is not declared in a generic instantiation (this includes generic subprogram instances), the source trace includes only one source line reference. \texttt{LINE\_NUMBER} gives the line number of the occurrence of the declaration of \( E \) within the source file (as a decimal literal without an exponent or point).

* If \( E \) is declared by a generic instantiation, its source trace (from left to right) starts with the source location of the declaration of \( E \) in the generic unit and ends with the source location of the instantiation, given in square brackets. This approach is applied recursively with nested instantiations: the rightmost (nested most deeply in square brackets) element of the source trace is the location of the outermost instantiation, and the leftmost element (that is, outside of any square brackets) is the location of the declaration of \( E \) in the generic unit.

Examples:

```plaintext
pragma Eliminate (Pkg0, Proc);
-- Eliminate (all overloading of) Proc in Pkg0

pragma Eliminate (Pkg1, Proc,
    Source_Location => "pkg1.ads:8");
-- Eliminate overloading of Proc at line 8 in pkg1.ads

-- Assume the following file contents:
-- gen_pkg.ads
-- 1: generic
-- 2: type T is private;
-- 3: package Gen_Pkg is
-- 4:   procedure Proc(N : T);
-- ... ...
-- ... end Gen_Pkg;
--
-- q.adb
-- 1: with Gen_Pkg;
-- 2: procedure Q is
-- 3:   package Inst_Pkg is new Gen_Pkg(Integer);
-- ... -- No calls on Inst_Pkg.Proc
-- ... end Q;

-- The following pragma eliminates Inst_Pkg.Proc from Q
pragma Eliminate (Q, Proc,
    Source_Location => "gen_pkg.ads:4[q.adb:3]");
```

2.55 Pragma Enable\_Atomic\_Synchronization

Syntax:

```plaintext
pragma Enable_Atomic_Synchronization [(Entity)];
```

Ada requires that accesses (reads or writes) of an atomic variable be regarded as synchronization points in the case of multiple tasks. Particularly in the case of multi-processors this may require special handling, e.g. the generation of memory barriers. This synchro-
nization is performed by default, but can be turned off using `pragma Disable_Atomic_Synchronization`. The `Enable_Atomic_Synchronization` pragma can be used to turn it back on.

The placement and scope rules for this pragma are the same as those for `pragma Unsuppress`. In particular it can be used as a configuration pragma, or in a declaration sequence where it applies till the end of the scope. If an `Entity` argument is present, the action applies only to that entity.

### 2.56 Pragma Export_Function

Syntax:

```plaintext
pragma Export_Function (  
    [Internal =>] LOCAL_NAME  
    [, [External =>] EXTERNAL_SYMBOL]  
    [, [Parameter_Types =>] PARAMETER_TYPES]  
    [, [Result_Type =>] result_SUBTYPE_MARK]  
    [, [Mechanism =>] MECHANISM]  
    [, [Result_Mechanism =>] MECHANISM_NAME]);

EXTERNAL_SYMBOL ::=  
    IDENTIFIER  
    | static_string_EXPRESSION  
    | ""

PARAMETER_TYPES ::=  
    null  
    | TYPE_DESIGNATOR {, TYPE_DESIGNATOR}

TYPE_DESIGNATOR ::=  
    subtype_NAME  
    | subtype_Name ' Access

MECHANISM ::=  
    MECHANISM_NAME  
    | (MECHANISM_ASSOCIATION {, MECHANISM_ASSOCIATION})

MECHANISM_ASSOCIATION ::=  
    [formal_parameter_NAME =>] MECHANISM_NAME

MECHANISM_NAME ::= Value | Reference
```

Use this pragma to make a function externally callable and optionally provide information on mechanisms to be used for passing parameter and result values. We recommend, for the purposes of improving portability, this pragma always be used in conjunction with a separate pragma `Export`, which must precede the pragma `Export_Function`. GNAT does not require a separate pragma `Export`, but if none is present, `Convention Ada` is assumed, which is usually not what is wanted, so it is usually appropriate to use this pragma in conjunction with an `Export` or `Convention` pragma that specifies the desired foreign convention. Pragma
**Export Function** (and **Export**, if present) must appear in the same declarative region as the function to which they apply.

The **internal name** must uniquely designate the function to which the pragma applies. If more than one function name exists of this name in the declarative part you must use the **Parameter_Types** and **Result_Type** parameters to achieve the required unique designation. The **subtype_marks** in these parameters must exactly match the subtypes in the corresponding function specification, using positional notation to match parameters with subtype marks. The form with an 'Access attribute can be used to match an anonymous access parameter.

Special treatment is given if the EXTERNAL is an explicit null string or a static string expressions that evaluates to the null string. In this case, no external name is generated. This form still allows the specification of parameter mechanisms.

### 2.57 Pragma Export_Object

**Syntax:**

```plaintext
pragma Export_Object (  
    [Internal =>] LOCAL_NAME  
    [, [External =>] EXTERNAL_SYMBOL]  
    [, [Size    =>] EXTERNAL_SYMBOL});
```

**EXTERNAL_SYMBOL ::=**

```plaintext
IDENTIFIER  
| static_string_EXPRESSION
```

This pragma designates an object as exported, and apart from the extended rules for external symbols, is identical in effect to the use of the normal **Export** pragma applied to an object. You may use a separate Export pragma (and you probably should from the point of view of portability), but it is not required. **Size** is syntax checked, but otherwise ignored by GNAT.

### 2.58 Pragma Export_Procedure

**Syntax:**

```plaintext
pragma Export_Procedure (  
    [Internal      =>] LOCAL_NAME  
    [, [External   =>] EXTERNAL_SYMBOL]  
    [, [Parameter_Types =>] PARAMETER_TYPES]  
    [, [Mechanism  =>] MECHANISM]);
```

**EXTERNAL_SYMBOL ::=**

```plaintext
IDENTIFIER  
| static_string_EXPRESSION  
| ""
```

**PARAMETER_TYPES ::=**

```plaintext
null  
| TYPE_DESIGNATOR {, TYPE_DESIGNATOR}
```
Chapter 2: Implementation Defined Pragmas

TYPE_DESIGNATOR ::= 
    subtype_NAME 
| subtype_Name ' Access

MECHANISM ::= 
    MECHANISM_NAME 
| (MECHANISM_ASSOCIATION {, MECHANISM_ASSOCIATION})

MECHANISM_ASSOCIATION ::= 
    [formal_parameter_NAME =>] MECHANISM_NAME

MECHANISM_NAME ::= Value | Reference

This pragma is identical to Export_Function except that it applies to a procedure rather than a function and the parameters Result_Type and Result_Mechanism are not permitted. GNAT does not require a separate pragma Export, but if none is present, Convention Ada is assumed, which is usually not what is wanted, so it is usually appropriate to use this pragma in conjunction with a Export or Convention pragma that specifies the desired foreign convention.

Special treatment is given if the EXTERNAL is an explicit null string or a static string expressions that evaluates to the null string. In this case, no external name is generated. This form still allows the specification of parameter mechanisms.

2.59 Pragma Export_Valued_Procedure

Syntax:

pragma Export_Valued_Procedure ( 
    [Internal =>] LOCAL_NAME 
[, [External =>] EXTERNAL_SYMBOL] 
[, [Parameter_Types =>] PARAMETER_TYPES] 
[, [Mechanism =>] MECHANISM]);

EXTERNAL_SYMBOL ::= 
    IDENTIFIER 
| static_string_EXPRESSION 
| ""

PARAMETER_TYPES ::= 
    null 
| TYPE_DESIGNATOR {, TYPE_DESIGNATOR}

TYPE_DESIGNATOR ::= 
    subtype_NAME 
| subtype_Name ' Access

MECHANISM ::= 
    MECHANISM_NAME
(MECHANISM_ASSOCIATION {, MECHANISM_ASSOCIATION})

MECHANISM_ASSOCIATION ::= [formal_parameter_NAME =>] MECHANISM_NAME

MECHANISM_NAME ::= Value | Reference

This pragma is identical to Export_Procedure except that the first parameter of LOCAL_NAME, which must be present, must be of mode out, and externally the subprogram is treated as a function with this parameter as the result of the function. GNAT provides for this capability to allow the use of out and in out parameters in interfacing to external functions (which are not permitted in Ada functions). GNAT does not require a separate pragma Export, but if none is present, Convention Ada is assumed, which is almost certainly not what is wanted since the whole point of this pragma is to interface with foreign language functions, so it is usually appropriate to use this pragma in conjunction with a Export or Convention pragma that specifies the desired foreign convention.

Special treatment is given if the EXTERNAL is an explicit null string or a static string expressions that evaluates to the null string. In this case, no external name is generated. This form still allows the specification of parameter mechanisms.

2.60 Pragma Extend_System

Syntax:

pragma Extend_System ([Name =>] IDENTIFIER);

This pragma is used to provide backwards compatibility with other implementations that extend the facilities of package System. In GNAT, System contains only the definitions that are present in the Ada RM. However, other implementations, notably the DEC Ada 83 implementation, provide many extensions to package System.

For each such implementation accommodated by this pragma, GNAT provides a package Aux_xxx, e.g., Aux_DEC for the DEC Ada 83 implementation, which provides the required additional definitions. You can use this package in two ways. You can with it in the normal way and access entities either by selection or using a use clause. In this case no special processing is required.

However, if existing code contains references such as System.xxx where ‘xxx’ is an entity in the extended definitions provided in package System, you may use this pragma to extend visibility in System in a non-standard way that provides greater compatibility with the existing code. Pragma Extend_System is a configuration pragma whose single argument is the name of the package containing the extended definition (e.g., Aux_DEC for the DEC Ada case). A unit compiled under control of this pragma will be processed using special visibility processing that looks in package System.Aux_xxx where Aux_xxx is the pragma argument for any entity referenced in package System, but not found in package System.

You can use this pragma either to access a predefined System extension supplied with the compiler, for example Aux_DEC or you can construct your own extension unit following the above definition. Note that such a package is a child of System and thus is considered part of the implementation. To compile it you will have to use the ‘-gnatg’ switch for compiling System units, as explained in the GNAT User’s Guide.
2.61 Pragma Extensions_Allowed

Syntax:

\[ \text{pragma Extensions\_Allowed (On | Off | All)}; \]

This configuration pragma enables (via the “On” or “All” argument) or disables (via the “Off” argument) the implementation extension mode; the pragma takes precedence over the ‘-gnatX’ and ‘-gnatX0’ command switches.

If an argument of “All” is specified, the latest version of the Ada language is implemented (currently Ada 2022) and, in addition, a number of GNAT specific extensions are recognized. These extensions are listed below. An argument of “On” has the same effect except that only some, not all, of the listed extensions are enabled; those extensions are identified below.

* Constrained attribute for generic objects
  
The \texttt{Constrained} attribute is permitted for objects of generic types. The result indicates if the corresponding actual is constrained.

* Static aspect on intrinsic functions
  
The Ada 202x \texttt{Static} aspect can be specified on Intrinsic imported functions and the compiler will evaluate some of these intrinsic statically, in particular the \texttt{Shift\_Left} and \texttt{Shift\_Right} intrinsics.

An Extensions\_Allowed pragma argument of “On” enables this extension.

* [] aggregates
  
This new aggregate syntax for arrays and containers is provided under -gnatX to experiment and confirm this new language syntax.

* Additional when constructs
  
In addition to the \texttt{exit when CONDITION} control structure, several additional constructs are allowed following this format. Including \texttt{return when CONDITION}, \texttt{goto when CONDITION}, and \texttt{raise [with EXCEPTION\_MESSAGE] when CONDITION}.

Some examples:

\[ \text{return Result when Variable > 10;} \]

\[ \text{raise Program\_Error with "Element is null" when Element = null;} \]

\[ \text{goto End\_Of\_Subprogram when Variable = -1;} \]

* Casing on composite values (aka pattern matching)
  
The selector for a case statement may be of a composite type, subject to some restrictions (described below). Aggregate syntax is used for choices of such a case statement; however, in cases where a “normal” aggregate would require a discrete value, a discrete subtype may be used instead; box notation can also be used to match all values.

Consider this example:

\[ \text{type Rec is record} \]
\[ \quad \text{F1, F2 : Integer;} \]
\[ \text{end record;} \]

\[ \text{procedure Caser_1 (X : Rec) is} \]
begin
  case X is
    when (F1 => Positive, F2 => Positive) =>
      Do_This;
    when (F1 => Natural, F2 => <> | (F1 => <>, F2 => Natural)) =>
      Do_That;
    when others =>
      Do_The_Other_Thing;
  end case;
end Caser_1;

If Caser_1 is called and both components of X are positive, then Do_This will be called; otherwise, if either component is nonnegative then Do_That will be called; otherwise, Do_The_Other_Thing will be called.

If the set of values that match the choice(s) of an earlier alternative overlaps the corresponding set of a later alternative, then the first set shall be a proper subset of the second (and the later alternative will not be executed if the earlier alternative “matches”). All possible values of the composite type shall be covered. The composite type of the selector shall be an array or record type that is neither limited class-wide. Currently, a “when others =>” case choice is required; it is intended that this requirement will be relaxed at some point.

If a subcomponent’s subtype does not meet certain restrictions, then the only value that can be specified for that subcomponent in a case choice expression is a “box” component association (which matches all possible values for the subcomponent). This restriction applies if

- the component subtype is not a record, array, or discrete type; or
- the component subtype is subject to a non-static constraint or has a predicate; or
- the component type is an enumeration type that is subject to an enumeration representation clause; or
- the component type is a multidimensional array type or an array type with a nonstatic index subtype.

Support for casing on arrays (and on records that contain arrays) is currently subject to some restrictions. Non-positional array aggregates are not supported as (or within) case choices. Likewise for array type and subtype names. The current implementation exceeds compile-time capacity limits in some annoyingly common scenarios; the message generated in such cases is usually “Capacity exceeded in compiling case statement with composite selector type”.

In addition, pattern bindings are supported. This is a mechanism for binding a name to a component of a matching value for use within an alternative of a case statement. For a component association that occurs within a case choice, the expression may be followed by “is <identifier>”. In the special case of a “box” component association, the identifier may instead be provided within the box. Either of these indicates that the given identifier denotes (a constant view of) the matching subcomponent of the case selector. Binding is not yet supported for arrays or subcomponents thereof.

Consider this example (which uses type Rec from the previous example):
Chapter 2: Implementation Defined Pragmas

procedure Caser_2 (X : Rec) is
begin
  case X is
    when (F1 => Positive is Abc, F2 => Positive) =>
      Do_This (Abc)
    when (F1 => Natural is N1, F2 => <N2>) | 
      (F1 => <N2>, F2 => Natural is N1) =>
      Do_That (Param_1 => N1, Param_2 => N2);
    when others =>
      Do_The_Other_Thing;
  end case;
end Caser_2;

This example is the same as the previous one with respect to determining whether Do_This, Do_That, or Do_The_Other_Thing will be called. But for this version, Do_This takes a parameter and Do_That takes two parameters. If Do_This is called, the actual parameter in the call will be X.F1.

If Do_That is called, the situation is more complex because there are two choices for that alternative. If Do_That is called because the first choice matched (i.e., because X.F1 is nonnegative and either X.F1 or X.F2 is zero or negative), then the actual parameters of the call will be (in order) X.F1 and X.F2. If Do_That is called because the second choice matched (and the first one did not), then the actual parameters will be reversed.

Within the choice list for single alternative, each choice must define the same set of bindings and the component subtypes for for a given identifier must all statically match. Currently, the case of a binding for a nondiscrete component is not implemented.

An Extensions_Allowed pragma argument of “On” enables this extension.

* Fixed lower bounds for array types and subtypes

Unconstrained array types and subtypes can be specified with a lower bound that is fixed to a certain value, by writing an index range that uses the syntax “<lower-bound-expression> .. <>”. This guarantees that all objects of the type or subtype will have the specified lower bound.

For example, a matrix type with fixed lower bounds of zero for each dimension can be declared by the following:

  type Matrix is
    array (Natural range 0 .. <>, Natural range 0 .. <>) of Integer;

Objects of type Matrix declared with an index constraint must have index ranges starting at zero:

  M1 : Matrix (0 .. 9, 0 .. 19);
  M2 : Matrix (2 .. 11, 3 .. 22); -- Warning about bounds; will raise CE

Similarly, a subtype of String can be declared that specifies the lower bound of objects of that subtype to be 1:

  subtype String_1 is String (1 .. <>);

If a string slice is passed to a formal of subtype String_1 in a call to a subprogram S, the slice’s bounds will “slide” so that the lower bound is 1. Within S, the lower
bound of the formal is known to be 1, so, unlike a normal unconstrained String formal, there is no need to worry about accounting for other possible lower-bound values. Sliding of bounds also occurs in other contexts, such as for object declarations with an unconstrained subtype with fixed lower bound, as well as in subtype conversions.

Use of this feature increases safety by simplifying code, and can also improve the efficiency of indexing operations, since the compiler statically knows the lower bound of unconstrained array formals when the formal's subtype has index ranges with static fixed lower bounds.

An Extensions_Allowed pragma argument of “On” enables this extension.

* Prefixed-view notation for calls to primitive subprograms of untagged types

Since Ada 2005, calls to primitive subprograms of a tagged type that have a “pre-fixed view” (see RM 4.1.3(9.2)) have been allowed to be written using the form of a selected_component, with the first actual parameter given as the prefix and the name of the subprogram as a selector. This prefixed-view notation for calls is extended so as to also allow such syntax for calls to primitive subprograms of untagged types. The primitives of an untagged type T that have a prefixed view are those where the first formal parameter of the subprogram either is of type T or is an anonymous access parameter whose designated type is T. For a type that has a component that happens to have the same simple name as one of the type’s primitive subprograms, where the component is visible at the point of a selected_component using that name, preference is given to the component in a selected_component (as is currently the case for tagged types with such component names).

An Extensions_Allowed pragma argument of “On” enables this extension.

* Expression defaults for generic formal functions

The declaration of a generic formal function is allowed to specify an expression as a default, using the syntax of an expression function.

Here is an example of this feature:

```ada
generic
    type T is private;
    with function Copy (Item : T) return T is (Item); -- Defaults to Item
package Stacks is
    type Stack is limited private;
    procedure Push (S : in out Stack; X : T); -- Calls Copy on X
    function Pop (S : in out Stack) return T; -- Calls Copy to return item
private
    -- ...
end Stacks;
```

2.62 Pragma Extensions_Visible

Syntax:
pragma Extensions_Visible [ (static_boolean_EXPRESSION) ];

For the semantics of this pragma, see the entry for aspect Extensions_Visible in the SPARK 2014 Reference Manual, section 6.1.7.

2.63 Pragma External

Syntax:

    pragma External (  
        [ Convention => ] convention_IDENTIFIER,  
        [ Entity => ] LOCAL_NAME  
        [, [External_Name => ] static_string_EXPRESSION ]  
        [, [Link_Name => ] static_string_EXPRESSION ] );

This pragma is identical in syntax and semantics to pragma Export as defined in the Ada Reference Manual. It is provided for compatibility with some Ada 83 compilers that used this pragma for exactly the same purposes as pragma Export before the latter was standardized.

2.64 Pragma External_Name_Casing

Syntax:

    pragma External_Name_Casing (  
        Uppercase | Lowercase  
        [, Uppercase | Lowercase | As_Is] );

This pragma provides control over the casing of external names associated with Import and Export pragmas. There are two cases to consider:

* Implicit external names

Implicit external names are derived from identifiers. The most common case arises when a standard Ada Import or Export pragma is used with only two arguments, as in:

    pragma Import (C, C_Routine);

Since Ada is a case-insensitive language, the spelling of the identifier in the Ada source program does not provide any information on the desired casing of the external name, and so a convention is needed. In GNAT the default treatment is that such names are converted to all lower case letters. This corresponds to the normal C style in many environments. The first argument of pragma External_Name_Casing can be used to control this treatment. If Uppercase is specified, then the name will be forced to all uppercase letters. If Lowercase is specified, then the normal default of all lower case letters will be used.

This same implicit treatment is also used in the case of extended DEC Ada 83 compatible Import and Export pragmas where an external name is explicitly specified using an identifier rather than a string.

* Explicit external names

Explicit external names are given as string literals. The most common case arises when a standard Ada Import or Export pragma is used with three arguments, as in:
pragma Import (C, C_Routine, "C_routine");

In this case, the string literal normally provides the exact casing required for the external name. The second argument of pragma External_Name_Casing may be used to modify this behavior. If Uppercase is specified, then the name will be forced to all uppercase letters. If Lowercase is specified, then the name will be forced to all lowercase letters. A specification of As_Is provides the normal default behavior in which the casing is taken from the string provided.

This pragma may appear anywhere that a pragma is valid. In particular, it can be used as a configuration pragma in the gnat.adc file, in which case it applies to all subsequent compilations, or it can be used as a program unit pragma, in which case it only applies to the current unit, or it can be used more locally to control individual Import/Export pragmas.

It was primarily intended for use with OpenVMS systems, where many compilers convert all symbols to upper case by default. For interfacing to such compilers (e.g., the DEC C compiler), it may be convenient to use the pragma:

    pragma External_Name_Casing (Uppercase, Uppercase);

to enforce the upper casing of all external symbols.

2.65 Pragma Fast_Math

Syntax:

    pragma Fast_Math;

This is a configuration pragma which activates a mode in which speed is considered more important for floating-point operations than absolutely accurate adherence to the requirements of the standard. Currently the following operations are affected:

‘Complex Multiplication’

The normal simple formula for complex multiplication can result in intermediate overflows for numbers near the end of the range. The Ada standard requires that this situation be detected and corrected by scaling, but in Fast_Math mode such cases will simply result in overflow. Note that to take advantage of this you must instantiate your own version of Ada.Numerics.Generic_Complex_Types under control of the pragma, rather than use the preinstantiated versions.

2.66 Pragma Favor_Top_Level

Syntax:

    pragma Favor_Top_Level (type_NAME);

The argument of pragma Favor_Top_Level must be a named access-to-subprogram type. This pragma is an efficiency hint to the compiler, regarding the use of 'Access or 'Unrestricted_Access on nested (non-library-level) subprograms. The pragma means that nested subprograms are not used with this type, or are rare, so that the generated code should be efficient in the top-level case. When this pragma is used, dynamically generated trampolines may be used on some targets for nested subprograms. See restriction No_Implicit_Dynamic_Code.
2.67 Pragma Finalize_Storage_Only

Syntax:

```
pragma Finalize_Storage_Only (first_subtype_LOCAL_NAME);
```

The argument of pragma `Finalize_Storage_Only` must denote a local type which is derived from `Ada.Finalization.Controlled` or `Limited_Controlled`. The pragma suppresses the call to `Finalize` for declared library-level objects of the argument type. This is mostly useful for types where finalization is only used to deal with storage reclamation since in most environments it is not necessary to reclaim memory just before terminating execution, hence the name. Note that this pragma does not suppress `Finalize` calls for library-level heap-allocated objects (see pragma `No_Heap_Finalization`).

2.68 Pragma Float_Representation

Syntax:

```
pragma Float_Representation (FLOAT_REP[, float_type_LOCAL_NAME]);
```

```
FLOAT_REP ::= VAX_Float | IEEE_Float
```

In the one argument form, this pragma is a configuration pragma which allows control over the internal representation chosen for the predefined floating point types declared in the packages `Standard` and `System`. This pragma is only provided for compatibility and has no effect.

The two argument form specifies the representation to be used for the specified floating-point type. The argument must be `IEEE_Float` to specify the use of IEEE format, as follows:

* For a digits value of 6, 32-bit IEEE short format will be used.
* For a digits value of 15, 64-bit IEEE long format will be used.
* No other value of digits is permitted.

2.69 Pragma Ghost

Syntax:

```
pragma Ghost [ (static_boolean_EXPRESSION) ];
```

For the semantics of this pragma, see the entry for aspect `Ghost` in the SPARK 2014 Reference Manual, section 6.9.

2.70 Pragma Global

Syntax:

```
pragma Global (GLOBAL_SPECIFICATION);
```

```
GLOBAL_SPECIFICATION ::= null
| (GLOBAL_LIST)
| (MODED_GLOBAL_LIST {, MODED_GLOBAL_LIST})
```
MODED_GLOBAL_LIST ::= MODE_SELECTOR => GLOBAL_LIST

MODE_SELECTOR ::= In_Out | Input | Output | Proof_In
GLOBAL_LIST ::= GLOBAL_ITEM | (GLOBAL_ITEM {, GLOBAL_ITEM})
GLOBAL_ITEM ::= NAME

For the semantics of this pragma, see the entry for aspect Global in the SPARK 2014 Reference Manual, section 6.1.4.

2.71 Pragma Ident

Syntax:

pragma Ident (static_string_EXPRESSION);

This pragma is identical in effect to pragma Comment. It is provided for compatibility with other Ada compilers providing this pragma.

2.72 Pragma Ignore_Pragma

Syntax:

pragma Ignore_Pragma (pragma_IDENTIFIER);

This is a configuration pragma that takes a single argument that is a simple identifier. Any subsequent use of a pragma whose pragma identifier matches this argument will be silently ignored. This may be useful when legacy code or code intended for compilation with some other compiler contains pragmas that match the name, but not the exact implementation, of a GNAT pragma. The use of this pragma allows such pragmas to be ignored, which may be useful in CodePeer mode, or during porting of legacy code.

2.73 Pragma Implementation_Defined

Syntax:

pragma Implementation_Defined (local_NAME);

This pragma marks a previously declared entity as implementation-defined. For an overloaded entity, applies to the most recent homonym.

pragma Implementation_Defined;

The form with no arguments appears anywhere within a scope, most typically a package spec, and indicates that all entities that are defined within the package spec are Implementation_Defined.

This pragma is used within the GNAT runtime library to identify implementation-defined entities introduced in language-defined units, for the purpose of implementing the No_Implementation_Identifiers restriction.

2.74 Pragma Implemented

Syntax:

pragma Implemented (procedure_LOCAL_NAME, implementation_kind);

implementation_kind ::= By_Entry | By_Protected_Procedure | By_Any
This is an Ada 2012 representation pragma which applies to protected, task and synchronized interface primitives. The use of pragma Implemented provides a way to impose a static requirement on the overriding operation by adhering to one of the three implementation kinds: entry, protected procedure or any of the above. This pragma is available in all earlier versions of Ada as an implementation-defined pragma.

```ada
type Synch_Iface is synchronized interface;
procedure Prim_Op (Obj : in out Iface) is abstract;
pragma Implemented (Prim_Op, By_Protected_Procedure);

protected type Prot_1 is new Synch_Iface with
  procedure Prim_Op; -- Legal
end Prot_1;

protected type Prot_2 is new Synch_Iface with
  entry Prim_Op;  -- Illegal
end Prot_2;

task type Task_Typ is new Synch_Iface with
  entry Prim_Op;  -- Illegal
end Task_Typ;
```

When applied to the procedure_or_entry_NAME of a requeue statement, pragma Implemented determines the runtime behavior of the requeue. Implementation kind By_Entry guarantees that the action of requeueing will proceed from an entry to another entry. Implementation kind By_Protected_Procedure transforms the requeue into a dispatching call, thus eliminating the chance of blocking. Kind By.Any shares the behavior of By_Entry and By_Protected_Procedure depending on the target’s overriding subprogram kind.

### 2.75 Pragma Implicit_Packing

**Syntax:**

```ada
pragma Implicit_Packing;
```

This is a configuration pragma that requests implicit packing for packed arrays for which a size clause is given but no explicit pragma Pack or specification of Component_Size is present. It also applies to records where no record representation clause is present. Consider this example:

```ada
type R is array (0 .. 7) of Boolean;
for R'Size use 8;
```

In accordance with the recommendation in the RM (RM 13.3(53)), a Size clause does not change the layout of a composite object. So the Size clause in the above example is normally rejected, since the default layout of the array uses 8-bit components, and thus the array requires a minimum of 64 bits.

If this declaration is compiled in a region of code covered by an occurrence of the configuration pragma Implicit_Packing, then the Size clause in this and similar examples will cause implicit packing and thus be accepted. For this implicit packing to occur, the type in question must be an array of small components whose size is known at compile time, and the Size clause must specify the exact size that corresponds to the number of elements in the array.
multiplied by the size in bits of the component type (both single and multi-dimensioned arrays can be controlled with this pragma).

Similarly, the following example shows the use in the record case

```pascal
type r is record
  a, b, c, d, e, f, g, h : boolean;
  chr : character;
end record;
for r'size use 16;
```

Without a pragma Pack, each Boolean field requires 8 bits, so the minimum size is 72 bits, but with a pragma Pack, 16 bits would be sufficient. The use of pragma Implicit_Packing allows this record declaration to compile without an explicit pragma Pack.

### 2.76 Pragma Import_Function

Syntax:

```pascal
pragma Import_Function (  
  [Internal =>] LOCAL_NAME,  
  [, [External =>] EXTERNAL_SYMBOL]  
  [, [Parameter_Types =>] PARAMETER_TYPES]  
  [, [Result_Type =>] SUBTYPE_MARK]  
  [, [Mechanism =>] MECHANISM]  
  [, [Result_Mechanism =>] MECHANISM_NAME]);
```

**EXTERNAL_SYMBOL ::=**
- IDENTIFIER
  | static_string_EXPRESSION

**PARAMETER_TYPES ::=**
- null
  | TYPE_DESIGNATOR {, TYPE_DESIGNATOR}

**TYPE_DESIGNATOR ::=**
- subtype_NAME
  | subtype_Name ' Access

**MECHANISM ::=**
- MECHANISM_NAME
  | (MECHANISM_ASSOCIATION {, MECHANISM_ASSOCIATION})

**MECHANISM_ASSOCIATION ::=**
- [formal_parameter_NAME =>] MECHANISM_NAME

**MECHANISM_NAME ::=**
- Value
  | Reference
This pragma is used in conjunction with a pragma `Import` to specify additional information for an imported function. The pragma `Import` (or equivalent pragma `Interface`) must precede the `Import_Function` pragma and both must appear in the same declarative part as the function specification.

The `Internal` argument must uniquely designate the function to which the pragma applies. If more than one function name exists of this name in the declarative part you must use the `Parameter_Types` and `Result_Type` parameters to achieve the required unique designation. Subtype marks in these parameters must exactly match the subtypes in the corresponding function specification, using positional notation to match parameters with subtype marks. The form with an `Access` attribute can be used to match an anonymous access parameter.

You may optionally use the `Mechanism` and `Result_Mechanism` parameters to specify passing mechanisms for the parameters and result. If you specify a single mechanism name, it applies to all parameters. Otherwise you may specify a mechanism on a parameter by parameter basis using either positional or named notation. If the mechanism is not specified, the default mechanism is used.

### 2.77 Pragma Import_Object

Syntax:

```plaintext
```

```plaintext
EXTERNAL_SYMBOL ::= IDENTIFIER | static_string_EXPRESSION
```

This pragma designates an object as imported, and apart from the extended rules for external symbols, is identical in effect to the use of the normal `Import` pragma applied to an object. Unlike the subprogram case, you need not use a separate `Import` pragma, although you may do so (and probably should do so from a portability point of view). `size` is syntax checked, but otherwise ignored by GNAT.

### 2.78 Pragma Import_procedure

Syntax:

```plaintext
```

```plaintext
EXTERNAL_SYMBOL ::= IDENTIFIER | static_string_EXPRESSION
PARAMETER_TYPES ::= 
```

This pragma designates a subprogram as imported, and apart from the extended rules for external symbols, is identical in effect to the use of the normal `Import` pragma applied to a subprogram. Unlike the subprogram case, you need not use a separate `Import` pragma, although you may do so (and probably should do so from a portability point of view).
null
| TYPE_DESIGNATOR {, TYPE_DESIGNATOR}

TYPE_DESIGNATOR ::=  
  subtype_NAME  
| subtype_Name ' Access

MECHANISM ::=  
  MECHANISM_NAME  
| (MECHANISM_ASSOCIATION {, MECHANISM_ASSOCIATION})

MECHANISM_ASSOCIATION ::=  
  [formal_parameter_NAME =>] MECHANISM_NAME

MECHANISM_NAME ::= Value | Reference

This pragma is identical to `Import_Function` except that it applies to a procedure rather than a function and the parameters `Result_Type` and `Result_Mechanism` are not permitted.

### 2.79 Pragma Import_Valued_Procedure

Syntax:

```plaintext
pragma Import_Valued_Procedure (  
  [Internal =>] LOCAL_NAME  
| external => EXTERNAL_SYMBOL  
| [Parameter_Types =>] PARAMETER_TYPES  
| [Mechanism =>] MECHANISM);
```

EXTERNAL_SYMBOL ::=  
  IDENTIFIER  
| static_string_EXPRESSION

PARAMETER_TYPES ::=  
  null  
| TYPE_DESIGNATOR {, TYPE_DESIGNATOR}

TYPE_DESIGNATOR ::=  
  subtype_NAME  
| subtype_Name ' Access

MECHANISM ::=  
  MECHANISM_NAME  
| (MECHANISM_ASSOCIATION {, MECHANISM_ASSOCIATION})

MECHANISM_ASSOCIATION ::=  
  [formal_parameter_NAME =>] MECHANISM_NAME

MECHANISM_NAME ::= Value | Reference
This pragma is identical to Import_Procedure except that the first parameter of LOCAL_NAME, which must be present, must be of mode out, and externally the subprogram is treated as a function with this parameter as the result of the function. The purpose of this capability is to allow the use of out and in out parameters in interfacing to external functions (which are not permitted in Ada functions). You may optionally use the Mechanism parameters to specify passing mechanisms for the parameters. If you specify a single mechanism name, it applies to all parameters. Otherwise you may specify a mechanism on a parameter by parameter basis using either positional or named notation. If the mechanism is not specified, the default mechanism is used.

Note that it is important to use this pragma in conjunction with a separate pragma Import that specifies the desired convention, since otherwise the default convention is Ada, which is almost certainly not what is required.

### 2.80 Pragma Independent

Syntax:

```
pragma Independent (Local_NAME);
```

This pragma is standard in Ada 2012 mode (which also provides an aspect of the same name). It is also available as an implementation-defined pragma in all earlier versions. It specifies that the designated object or all objects of the designated type must be independently addressable. This means that separate tasks can safely manipulate such objects. For example, if two components of a record are independent, then two separate tasks may access these two components. This may place constraints on the representation of the object (for instance prohibiting tight packing).

### 2.81 Pragma Independent_Components

Syntax:

```
pragma Independent_Components (Local_NAME);
```

This pragma is standard in Ada 2012 mode (which also provides an aspect of the same name). It is also available as an implementation-defined pragma in all earlier versions. It specifies that the components of the designated object, or the components of each object of the designated type, must be independently addressable. This means that separate tasks can safely manipulate separate components in the composite object. This may place constraints on the representation of the object (for instance prohibiting tight packing).

### 2.82 Pragma Initial_Condition

Syntax:

```
pragma Initial_Condition (boolean_EXPRESSION);
```

For the semantics of this pragma, see the entry for aspect Initial_Condition in the SPARK 2014 Reference Manual, section 7.1.6.

### 2.83 Pragma Initialize_Scalars

Syntax:
pragma Initialize_Scalars
   [ ( TYPE_VALUE_PAIR { , TYPE_VALUE_PAIR } ) ];

TYPE_VALUE_PAIR ::= 
   SCALAR_TYPE => static_EXPRESSION

SCALAR_TYPE :=
   Short_Float
| Float
| Long_Float
| Long_Long_Flat
| Signed_8
| Signed_16
| Signed_32
| Signed_64
| Unsigned_8
| Unsigned_16
| Unsigned_32
| Unsigned_64

This pragma is similar to Normalize_Scalars conceptually but has two important differences.

First, there is no requirement for the pragma to be used uniformly in all units of a partition. In particular, it is fine to use this just for some or all of the application units of a partition, without needing to recompile the run-time library. In the case where some units are compiled with the pragma, and some without, then a declaration of a variable where the type is defined in package Standard or is locally declared will always be subject to initialization, as will any declaration of a scalar variable. For composite variables, whether the variable is initialized may also depend on whether the package in which the type of the variable is declared is compiled with the pragma.

The other important difference is that the programmer can control the value used for initializing scalar objects. This effect can be achieved in several different ways:

* At compile time, the programmer can specify the invalid value for a particular family of scalar types using the optional arguments of the pragma.

   The compile-time approach is intended to optimize the generated code for the pragma, by possibly using fast operations such as memset. Note that such optimizations require using values where the bytes all have the same binary representation.

* At bind time, the programmer has several options:

   * Initialization with invalid values (similar to Normalize_Scalars, though for Initialize_Scalars it is not always possible to determine the invalid values in complex cases like signed component fields with nonstandard sizes).

   * Initialization with high values.

   * Initialization with low values.

   * Initialization with a specific bit pattern.

See the GNAT User’s Guide for binder options for specifying these cases.
The bind-time approach is intended to provide fast turnaround for testing with different values, without having to recompile the program.

* At execution time, the programmer can specify the invalid values using an environment variable. See the GNAT User’s Guide for details.

The execution-time approach is intended to provide fast turnaround for testing with different values, without having to recompile and rebind the program.

Note that pragma `Initialize_Scalars` is particularly useful in conjunction with the enhanced validity checking that is now provided in GNAT, which checks for invalid values under more conditions. Using this feature (see description of the ‘-gnatV’ flag in the GNAT User’s Guide) in conjunction with pragma `Initialize_Scalars` provides a powerful new tool to assist in the detection of problems caused by uninitialized variables.

Note: the use of `Initialize_Scalars` has a fairly extensive effect on the generated code. This may cause your code to be substantially larger. It may also cause an increase in the amount of stack required, so it is probably a good idea to turn on stack checking (see description of stack checking in the GNAT User’s Guide) when using this pragma.

### 2.84 Pragma Initializes

**Syntax:**

```plaintext
pragma Initializes (INITIALIZATION_LIST);
```

**INITIALIZATION_LIST ::=**

```
null
| (INITIALIZATION_ITEM {, INITIALIZATION_ITEM})
```

**INITIALIZATION_ITEM ::=**

```
name [=> INPUT_LIST]
```

**INPUT_LIST ::=**

```
nul
| INPUT
| (INPUT {, INPUT})
```

**INPUT ::=**

```
name
```

For the semantics of this pragma, see the entry for aspect `Initializes` in the SPARK 2014 Reference Manual, section 7.1.5.

### 2.85 Pragma Inline_Always

**Syntax:**

```plaintext
pragma Inline_Always (NAME [, NAME]);
```

Similar to pragma `Inline` except that inlining is unconditional. `Inline_Always` instructs the compiler to inline every direct call to the subprogram or else to emit a compilation error, independently of any option, in particular ‘-gnatn’ or ‘-gnatN’ or the optimization level. It is an error to take the address or access of `NAME`. It is also an error to apply this pragma to a primitive operation of a tagged type. Thanks to such restrictions, the compiler is allowed to remove the out-of-line body of `NAME`. 
2.86 Pragma Inline_Generic

Syntax:

    pragma Inline_Generic (GNAME {, GNAME});

    GNAME ::= generic_unit_NAME | generic_instance_NAME

This pragma is provided for compatibility with Dec Ada 83. It has no effect in GNAT
(which always inlines generics), other than to check that the given names are all names of
generic units or generic instances.

2.87 Pragma Interface

Syntax:

    pragma Interface (  
        [Convention =>] convention_identifier,  
        [Entity =>] local_NAME  
                [, [External_Name =>] static_string_expression]  
                [, [Link_Name =>] static_string_expression]);

This pragma is identical in syntax and semantics to the standard Ada pragma Import.
It is provided for compatibility with Ada 83. The definition is upwards compatible both
with pragma Interface as defined in the Ada 83 Reference Manual, and also with some
extended implementations of this pragma in certain Ada 83 implementations. The only
difference between pragma Interface and pragma Import is that there is special circuitry
to allow both pragmas to appear for the same subprogram entity (normally it is illegal to
have multiple Import pragmas). This is useful in maintaining Ada 83/Ada 95 compatibility
and is compatible with other Ada 83 compilers.

2.88 Pragma Interface_Name

Syntax:

    pragma Interface_Name (  
        [Entity =>] LOCAL_NAME  
                [, [External_Name =>] static_string_EXPRESSION]  
                [, [Link_Name =>] static_string_EXPRESSION]);

This pragma provides an alternative way of specifying the interface name for an interfaced
subprogram, and is provided for compatibility with Ada 83 compilers that use the pragma
for this purpose. You must provide at least one of External_Name or Link_Name.

2.89 Pragma Interrupt_Handler

Syntax:

    pragma Interrupt_Handler (procedure_LOCAL_NAME);

This program unit pragma is supported for parameterless protected procedures as described
2.90 Pragma Interrupt_State

Syntax:

```
pragma Interrupt_State
([Name =>] value,
 [State =>] SYSTEM | RUNTIME | USER);
```

Normally certain interrupts are reserved to the implementation. Any attempt to attach an interrupt causes Program_Error to be raised, as described in RM C.3.2(22). A typical example is the SIGINT interrupt used in many systems for an Ctrl-C interrupt. Normally this interrupt is reserved to the implementation, so that Ctrl-C can be used to interrupt execution. Additionally, signals such as SIGSEGV, SIGABRT, SIGFPE and SIGILL are often mapped to specific Ada exceptions, or used to implement run-time functions such as the abort statement and stack overflow checking.

Pragma Interrupt_State provides a general mechanism for overriding such uses of interrupts. It subsumes the functionality of pragma Unreserve_All_Interrupts. Pragma Interrupt_State is not available on Windows. On all other platforms than VxWorks, it applies to signals; on VxWorks, it applies to vectored hardware interrupts and may be used to mark interrupts required by the board support package as reserved.

Interrupts can be in one of three states:

* System
  The interrupt is reserved (no Ada handler can be installed), and the Ada run-time may not install a handler. As a result you are guaranteed standard system default action if this interrupt is raised. This also allows installing a low level handler via C APIs such as sigaction(), outside of Ada control.

* Runtime
  The interrupt is reserved (no Ada handler can be installed). The run time is allowed to install a handler for internal control purposes, but is not required to do so.

* User
  The interrupt is unreserved. The user may install an Ada handler via Ada.Interrupts and pragma Interrupt_Handler or Attach_Handler to provide some other action.

These states are the allowed values of the State parameter of the pragma. The Name parameter is a value of the type Ada.Interrupts.Interrupt_ID. Typically, it is a name declared in Ada.Interrupts.Names.

This is a configuration pragma, and the binder will check that there are no inconsistencies between different units in a partition in how a given interrupt is specified. It may appear anywhere a pragma is legal.

The effect is to move the interrupt to the specified state.

By declaring interrupts to be SYSTEM, you guarantee the standard system action, such as a core dump.

By declaring interrupts to be USER, you guarantee that you can install a handler.

Note that certain signals on many operating systems cannot be caught and handled by applications. In such cases, the pragma is ignored. See the operating system documentation, or the value of the array Reserved declared in the spec of package System.OS_Interface.
Overriding the default state of signals used by the Ada runtime may interfere with an application’s runtime behavior in the cases of the synchronous signals, and in the case of the signal used to implement the `abort` statement.

### 2.91 Pragma Invariant

**Syntax:**

```
pragma Invariant ([Entity =>] private_type_LOCAL_NAME,
                  [Check =>] EXPRESSION
                  [, [Message =>] String_Expression]);
```

This pragma provides exactly the same capabilities as the `Type_Invariant` aspect defined in AI05-0146-1, and in the Ada 2012 Reference Manual. The `Type_Invariant` aspect is fully implemented in Ada 2012 mode, but since it requires the use of the aspect syntax, which is not available except in 2012 mode, it is not possible to use the `Type_Invariant` aspect in earlier versions of Ada. However the `Invariant` pragma may be used in any version of Ada. Also note that the aspect `Invariant` is a synonym in GNAT for the aspect `Type_Invariant`, but there is no pragma `Type_Invariant`.

The pragma must appear within the visible part of the package specification, after the type to which its Entity argument appears. As with the `Invariant` aspect, the Check expression is not analyzed until the end of the visible part of the package, so it may contain forward references. The Message argument, if present, provides the exception message used if the invariant is violated. If no Message parameter is provided, a default message that identifies the line on which the pragma appears is used.

It is permissible to have multiple Invariants for the same type entity, in which case they are and’d together. It is permissible to use this pragma in Ada 2012 mode, but you cannot have both an invariant aspect and an invariant pragma for the same entity.

For further details on the use of this pragma, see the Ada 2012 documentation of the `Type_Invariant` aspect.

### 2.92 Pragma Keep_Names

**Syntax:**

```
pragma Keep_Names ([On =>] enumeration_first_subtype_LOCAL_NAME);
```

The LOCAL_NAME argument must refer to an enumeration first subtype in the current declarative part. The effect is to retain the enumeration literal names for use by `Image` and `Value` even if a global `Discard_Names` pragma applies. This is useful when you want to generally suppress enumeration literal names and for example you therefore use a `Discard_Names` pragma in the `gnat.adc` file, but you want to retain the names for specific enumeration types.

### 2.93 Pragma License

**Syntax:**

```
pragma License (Unrestricted | GPL | Modified_GPL | Restricted);
```

This pragma is provided to allow automated checking for appropriate license conditions with respect to the standard and modified GPL. A pragma `License`, which is a configuration
pragma that typically appears at the start of a source file or in a separate *gnat.adc* file, specifies the licensing conditions of a unit as follows:

* **Unrestricted** This is used for a unit that can be freely used with no license restrictions. Examples of such units are public domain units, and units from the Ada Reference Manual.

* **GPL** This is used for a unit that is licensed under the unmodified GPL, and which therefore cannot be *with*ed by a restricted unit.

* **Modified GPL** This is used for a unit licensed under the GNAT modified GPL that includes a special exception paragraph that specifically permits the inclusion of the unit in programs without requiring the entire program to be released under the GPL.

* **Restricted** This is used for a unit that is restricted in that it is not permitted to depend on units that are licensed under the GPL. Typical examples are proprietary code that is to be released under more restrictive license conditions. Note that restricted units are permitted to *with* units which are licensed under the modified GPL (this is the whole point of the modified GPL).

Normally a unit with no *License* pragma is considered to have an unknown license, and no checking is done. However, standard GNAT headers are recognized, and license information is derived from them as follows.

A GNAT license header starts with a line containing 78 hyphens. The following comment text is searched for the appearance of any of the following strings.

If the string ‘GNU General Public License’ is found, then the unit is assumed to have GPL license, unless the string ‘As a special exception’ follows, in which case the license is assumed to be modified GPL.

If one of the strings ‘This specification is adapted from the Ada Semantic Interface’ or ‘This specification is derived from the Ada Reference Manual’ is found then the unit is assumed to be unrestricted.

These default actions means that a program with a restricted license pragma will automatically get warnings if a GPL unit is inappropriately *with*ed. For example, the program:

```plaintext
with Sem_Ch3;
with GNAT.Sockets;
procedure Secret_Stuff is
...
end Secret_Stuff
```

if compiled with pragma *License* (*Restricted*) in a *gnat.adc* file will generate the warning:

1. with Sem_Ch3;
   
   | >> license of withed unit "Sem_Ch3" is incompatible

2. with GNAT.Sockets;
3. procedure Secret_Stuff is

Here we get a warning on *Sem_Ch3* since it is part of the GNAT compiler and is licensed under the GPL, but no warning for *GNAT.Sockets* which is part of the GNAT run time, and is therefore licensed under the modified GPL.
2.94 Pragma Link_With

Syntax:

    pragma Link_With (static_string_EXPRESSION {,static_string_EXPRESSION});

This pragma is provided for compatibility with certain Ada 83 compilers. It has exactly the same effect as pragma Linker_Options except that spaces occurring within one of the string expressions are treated as separators. For example, in the following case:

    pragma Link_With ('-labc -ldef');

results in passing the strings -labc and -ldef as two separate arguments to the linker. In addition pragma Link_With allows multiple arguments, with the same effect as successive pragmas.

2.95 Pragma Linker_Alias

Syntax:

    pragma Linker_Alias (    
        [Entity =>] LOCAL_NAME,    
        [Target =>] static_string_EXPRESSION);

LOCAL_NAME must refer to an object that is declared at the library level. This pragma establishes the given entity as a linker alias for the given target. It is equivalent to __attribute__((alias)) in GNU C and causes LOCAL_NAME to be emitted as an alias for the symbol static_string_EXPRESSION in the object file, that is to say no space is reserved for LOCAL_NAME by the assembler and it will be resolved to the same address as static_string_EXPRESSION by the linker.

The actual linker name for the target must be used (e.g., the fully encoded name with qualification in Ada, or the mangled name in C++), or it must be declared using the C convention with pragma Import or pragma Export.

Not all target machines support this pragma. On some of them it is accepted only if pragma Weak_External has been applied to LOCAL_NAME.

    -- Example of the use of pragma Linker_Alias

package p is
    i : Integer := 1;
    pragma Export (C, i);

    new_name_for_i : Integer;
    pragma Linker_Alias (new_name_for_i, "i");
end p;

2.96 Pragma Linker_Constructor

Syntax:

    pragma Linker_Constructor (procedure_LOCAL_NAME);

procedure_LOCAL_NAME must refer to a parameterless procedure that is declared at the library level. A procedure to which this pragma is applied will be treated as an initialization routine by the linker. It is equivalent to __attribute__((constructor)) in GNU C and
causes \texttt{procedure\_LOCAL\_NAME} to be invoked before the entry point of the executable is called (or immediately after the shared library is loaded if the procedure is linked in a shared library), in particular before the Ada run-time environment is set up.

Because of these specific contexts, the set of operations such a procedure can perform is very limited and the type of objects it can manipulate is essentially restricted to the elementary types. In particular, it must only contain code to which pragma Restrictions (\texttt{No\_Elaboration\_Code}) applies.

This pragma is used by GNAT to implement auto-initialization of shared Stand Alone Libraries, which provides a related capability without the restrictions listed above. Where possible, the use of Stand Alone Libraries is preferable to the use of this pragma.

\subsection*{2.97 Pragma Linker\_Destructor}

Syntax:
\begin{verbatim}
pragma Linker\_Destructor (procedure\_LOCAL\_NAME);
\end{verbatim}

\texttt{procedure\_LOCAL\_NAME} must refer to a parameterless procedure that is declared at the library level. A procedure to which this pragma is applied will be treated as a finalization routine by the linker. It is equivalent to \texttt{\_attribute\_\_((destructor))} in GNU C and causes \texttt{procedure\_LOCAL\_NAME} to be invoked after the entry point of the executable has exited (or immediately before the shared library is unloaded if the procedure is linked in a shared library), in particular after the Ada run-time environment is shut down.

See \texttt{pragma Linker\_Constructor} for the set of restrictions that apply because of these specific contexts.

\subsection*{2.98 Pragma Linker\_Section}

Syntax:
\begin{verbatim}
pragma Linker\_Section (  
    [Entity =>] LOCAL\_NAME,  
    [Section =>] static\_string\_EXPRESSION);
\end{verbatim}

\texttt{LOCAL\_NAME} must refer to an object, type, or subprogram that is declared at the library level. This pragma specifies the name of the linker section for the given entity. It is equivalent to \texttt{\_attribute\_\_((section))} in GNU C and causes \texttt{LOCAL\_NAME} to be placed in the \texttt{static\_string\_EXPRESSION} section of the executable (assuming the linker doesn’t rename the section). GNAT also provides an implementation defined aspect of the same name.

In the case of specifying this aspect for a type, the effect is to specify the corresponding section for all library-level objects of the type that do not have an explicit linker section set. Note that this only applies to whole objects, not to components of composite objects.

In the case of a subprogram, the linker section applies to all previously declared matching overloaded subprograms in the current declarative part which do not already have a linker section assigned. The linker section aspect is useful in this case for specifying different linker sections for different elements of such an overloaded set.

Note that an empty string specifies that no linker section is specified. This is not quite the same as omitting the pragma or aspect, since it can be used to specify that one element
of an overloaded set of subprograms has the default linker section, or that one object of a
type for which a linker section is specified should has the default linker section.
The compiler normally places library-level entities in standard sections depending on the
class: procedures and functions generally go in the .text section, initialized variables in
the .data section and uninitialized variables in the .bss section.
Other, special sections may exist on given target machines to map special hardware, for
example I/O ports or flash memory. This pragma is a means to defer the final layout of the
executable to the linker, thus fully working at the symbolic level with the compiler.
Some file formats do not support arbitrary sections so not all target machines support this
pragma. The use of this pragma may cause a program execution to be erroneous if it is used
to place an entity into an inappropriate section (e.g., a modified variable into the .text
section). See also pragma Persistent_BSS.

-- Example of the use of pragma Linker_Section

package IO_Card is
  Port_A : Integer;
  pragma Volatile (Port_A);
  pragma Linker_Section (Port_A, "bss.port_a");

  Port_B : Integer;
  pragma Volatile (Port_B);
  pragma Linker_Section (Port_B, "bss.port_b");

  type Port_Type is new Integer with Linker_Section => "bss";
  PA : Port_Type with Linker_Section => "bss.PA";
  PB : Port_Type; -- ends up in linker section "bss"

  procedure Q with Linker_Section => "Qsection";
end IO_Card;

2.99 Pragma Lock_Free

Syntax: This pragma may be specified for protected types or objects. It specifies that the
implementation of protected operations must be implemented without locks. Compilation
fails if the compiler cannot generate lock-free code for the operations.
The current conditions required to support this pragma are:
  * Protected type declarations may not contain entries
  * Protected subprogram declarations may not have nonelementary parameters
In addition, each protected subprogram body must satisfy:
  * May reference only one protected component
  * May not reference nonconstant entities outside the protected subprogram scope
  * May not contain address representation items, allocators, or quantified expressions
  * May not contain delay, goto, loop, or procedure-call statements
  * May not contain exported and imported entities
May not dereferenced access values
* Function calls and attribute references must be static

If the Lock_Free aspect is specified to be True for a protected unit and the Ceiling_Locking locking policy is in effect, then the run-time actions associated with the Ceiling_Locking locking policy (described in Ada RM D.3) are not performed when a protected operation of the protected unit is executed.

### 2.100 Pragma Loop_Invariant

**Syntax:**

```
pragma Loop_Invariant ( boolean_EXPRESSION );
```

The effect of this pragma is similar to that of pragma Assert, except that in an Assertion_Policy pragma, the identifier Loop_Invariant is used to control whether it is ignored or checked (or disabled).

Loop_Invariant can only appear as one of the items in the sequence of statements of a loop body, or nested inside block statements that appear in the sequence of statements of a loop body. The intention is that it be used to represent a “loop invariant” assertion, i.e. something that is true each time through the loop, and which can be used to show that the loop is achieving its purpose.

Multiple Loop_Invariant and Loop_Variant pragmas that apply to the same loop should be grouped in the same sequence of statements.

To aid in writing such invariants, the special attribute Loop_Entry may be used to refer to the value of an expression on entry to the loop. This attribute can only be used within the expression of a Loop_Invariant pragma. For full details, see documentation of attribute Loop_Entry.

### 2.101 Pragma Loop_Optimize

**Syntax:**

```
pragma Loop_Optimize (OPTIMIZATION_HINT {, OPTIMIZATION_HINT});
```

```
OPTIMIZATION_HINT ::= Ivdep | No_Unroll | Unroll | No_Vector | Vector
```

This pragma must appear immediately within a loop statement. It allows the programmer to specify optimization hints for the enclosing loop. The hints are not mutually exclusive and can be freely mixed, but not all combinations will yield a sensible outcome.

There are five supported optimization hints for a loop:

* **Ivdep**
  
The programmer asserts that there are no loop-carried dependencies which would prevent consecutive iterations of the loop from being executed simultaneously.

* **No_Unroll**
  
The loop must not be unrolled. This is a strong hint: the compiler will not unroll a loop marked with this hint.
* Unroll
  The loop should be unrolled. This is a weak hint: the compiler will try to apply unrolling to this loop preferably to other optimizations, notably vectorization, but there is no guarantee that the loop will be unrolled.

* No_Vector
  The loop must not be vectorized. This is a strong hint: the compiler will not vectorize a loop marked with this hint.

* Vector
  The loop should be vectorized. This is a weak hint: the compiler will try to apply vectorization to this loop preferably to other optimizations, notably unrolling, but there is no guarantee that the loop will be vectorized.

These hints do not remove the need to pass the appropriate switches to the compiler in order to enable the relevant optimizations, that is to say ‘-funroll-loops’ for unrolling and ‘-ftree-vectorize’ for vectorization.

### 2.102 Pragma Loop_Variant

Syntax:

```plaintext
pragma Loop_Variant ( LOOP_VARIANT_ITEM {, LOOP_VARIANT_ITEM } );
LOOP_VARIANT_ITEM ::= CHANGE_DIRECTION => discrete_EXPRESSION
CHANGE_DIRECTION ::= Increases | Decreases
```

Loop_Variant can only appear as one of the items in the sequence of statements of a loop body, or nested inside block statements that appear in the sequence of statements of a loop body. It allows the specification of quantities which must always decrease or increase in successive iterations of the loop. In its simplest form, just one expression is specified, whose value must increase or decrease on each iteration of the loop.

In a more complex form, multiple arguments can be given which are interpreted in a nesting lexicographic manner. For example:

```plaintext
pragma Loop_Variant (Increases => X, Decreases => Y);
```

specifies that each time through the loop either X increases, or X stays the same and Y decreases. A Loop_Variant pragma ensures that the loop is making progress. It can be useful in helping to show informally or prove formally that the loop always terminates.

Loop_Variant is an assertion whose effect can be controlled using an Assertion_Policy with a check name of Loop_Variant. The policy can be Check to enable the loop variant check, Ignore to ignore the check (in which case the pragma has no effect on the program), or Disable in which case the pragma is not even checked for correct syntax.

Multiple Loop_Invariant and Loop_Variant pragmas that apply to the same loop should be grouped in the same sequence of statements.

The Loop_Entry attribute may be used within the expressions of the Loop_Variant pragma to refer to values on entry to the loop.
2.103 Pragma Machine_Attribute

Syntax:

    pragma Machine_Attribute ( 
        [Entity =>] LOCAL_NAME, 
        [Attribute_Name =>] static_string_EXPRESSION 
        [, [Info =>] static_EXPRESSION {, static_EXPRESSION}] ]
    );

Machine-dependent attributes can be specified for types and/or declarations. This pragma is semantically equivalent to __attribute__((attribute_name)) (if info is not specified) or __attribute__((attribute_name(info))) or __attribute__((attribute_name(info,...))) in GNU C, where ‘attribute_name’ is recognized by the compiler middle-end or the TARGET_ATTRIBUTE_TABLE machine specific macro. Note that a string literal for the optional parameter info or the following ones is transformed by default into an identifier, which may make this pragma unusable for some attributes. For further information see GNU Compiler Collection (GCC) Internals.

2.104 Pragma Main

Syntax:

    pragma Main 
    (MAIN_OPTION [, MAIN_OPTION]);

    MAIN_OPTION ::= 
        [Stack_Size =>] static_integer_EXPRESSION 
    | [Task_Stack_Size_Default =>] static_integer_EXPRESSION 
    | [Time_Slicing_Enabled =>] static_boolean_EXPRESSION

This pragma is provided for compatibility with OpenVMS VAX Systems. It has no effect in GNAT, other than being syntax checked.

2.105 Pragma Main_Storage

Syntax:

    pragma Main_Storage 
    (MAIN_STORAGE_OPTION [, MAIN_STORAGE_OPTION]);

    MAIN_STORAGE_OPTION ::= 
        [WORKING_STORAGE =>] static_SIMPLE_EXPRESSION 
    | [TOP_GUARD =>] static_SIMPLE_EXPRESSION

This pragma is provided for compatibility with OpenVMS VAX Systems. It has no effect in GNAT, other than being syntax checked.

2.106 Pragma Max_Queue_Length

Syntax:

    pragma Max_Entry_Queue (static_integer_EXPRESSION);

This pragma is used to specify the maximum callers per entry queue for individual protected entries and entry families. It accepts a single integer (-1 or more) as a parameter and must appear after the declaration of an entry.
A value of -1 represents no additional restriction on queue length.

2.107 Pragma No_Body

Syntax:

```
pragma No_Body;
```

There are a number of cases in which a package spec does not require a body, and in fact a body is not permitted. GNAT will not permit the spec to be compiled if there is a body around. The pragma No_Body allows you to provide a body file, even in a case where no body is allowed. The body file must contain only comments and a single No_Body pragma. This is recognized by the compiler as indicating that no body is logically present.

This is particularly useful during maintenance when a package is modified in such a way that a body needed before is no longer needed. The provision of a dummy body with a No_Body pragma ensures that there is no interference from earlier versions of the package body.

2.108 Pragma No_Caching

Syntax:

```
pragma No_Caching [(static_boolean_EXPRESSION)];
```

For the semantics of this pragma, see the entry for aspect No_Caching in the SPARK 2014 Reference Manual, section 7.1.2.

2.109 Pragma No_Component_Reordering

Syntax:

```
pragma No_Component_Reordering [(Entity => type_LOCAL_NAME)];
```

`type_LOCAL_NAME` must refer to a record type declaration in the current declarative part. The effect is to preclude any reordering of components for the layout of the record, i.e. the record is laid out by the compiler in the order in which the components are declared textually. The form with no argument is a configuration pragma which applies to all record types declared in units to which the pragma applies and there is a requirement that this pragma be used consistently within a partition.

2.110 Pragma No_Elaboration_Code_All

Syntax:

```
pragma No_Elaboration_Code_All [(program_unit_NAME)];
```

This is a program unit pragma (there is also an equivalent aspect of the same name) that establishes the restriction No_Elaboration_Code for the current unit and any extended main source units (body and subunits). It also has the effect of enforcing a transitive application of this aspect, so that if any unit is implicitly or explicitly with’ed by the current unit, it must also have the No_Elaboration_Code_All aspect set. It may be applied to package or subprogram specs or their generic versions.
2.111 Pragma No_Heap_Finalization

Syntax:

```
pragma No_Heap_Finalization [ (first_subtype_LOCAL_NAME) ];
```

Pragma No_Heap_Finalization may be used as a configuration pragma or as a type-specific pragma.

In its configuration form, the pragma must appear within a configuration file such as gnat.adc, without an argument. The pragma suppresses the call to Finalize for heap-allocated objects created through library-level named access-to-object types in cases where the designated type requires finalization actions.

In its type-specific form, the argument of the pragma must denote a library-level named access-to-object type. The pragma suppresses the call to Finalize for heap-allocated objects created through the specific access type in cases where the designated type requires finalization actions.

It is still possible to finalize such heap-allocated objects by explicitly deallocating them.

A library-level named access-to-object type declared within a generic unit will lose its No_Heap_Finalization pragma when the corresponding instance does not appear at the library level.

2.112 Pragma No_Inline

Syntax:

```
pragma No_Inline (NAME {, NAME});
```

This pragma suppresses inlining for the callable entity or the instances of the generic subprogram designated by NAME, including inlining that results from the use of pragma Inline. This pragma is always active, in particular it is not subject to the use of option ‘-gnatn’ or ‘-gnatN’. It is illegal to specify both pragma No_Inline and pragma Inline_Always for the same NAME.

2.113 Pragma No_Return

Syntax:

```
pragma No_Return (procedure_LOCAL_NAME {, procedure_LOCAL_NAME});
```

Each procedure_LOCAL_NAME argument must refer to one or more procedure declarations in the current declarative part. A procedure to which this pragma is applied may not contain any explicit return statements. In addition, if the procedure contains any implicit returns from falling off the end of a statement sequence, then execution of that implicit return will cause Program_Error to be raised.

One use of this pragma is to identify procedures whose only purpose is to raise an exception. Another use of this pragma is to suppress incorrect warnings about missing returns in functions, where the last statement of a function statement sequence is a call to such a procedure.

Note that in Ada 2005 mode, this pragma is part of the language. It is available in all earlier versions of Ada as an implementation-defined pragma.
### 2.114 Pragma No_Strict_Aliasing

Syntax:

```
pragma No_Strict_Aliasing ([Entity =>] type_LOCAL_NAME);
```

type_LOCAL_NAME must refer to an access type declaration in the current declarative part. The effect is to inhibit strict aliasing optimization for the given type. The form with no arguments is a configuration pragma which applies to all access types declared in units to which the pragma applies. For a detailed description of the strict aliasing optimization, and the situations in which it must be suppressed, see the section on Optimization and Strict Aliasing in the GNAT User's Guide.

This pragma currently has no effects on access to unconstrained array types.

### 2.115 Pragma No_Tagged_Streams

Syntax:

```
pragma No_Tagged_Streams ([Entity =>] tagged_type_LOCAL_NAME);
```

Normally when a tagged type is introduced using a full type declaration, part of the processing includes generating stream access routines to be used by stream attributes referencing the type (or one of its subtypes or derived types). This can involve the generation of significant amounts of code which is wasted space if stream routines are not needed for the type in question.

The No_Tagged_Streams pragma causes the generation of these stream routines to be skipped, and any attempt to use stream operations on types subject to this pragma will be statically rejected as illegal.

There are two forms of the pragma. The form with no arguments must appear in a declarative sequence or in the declarations of a package spec. This pragma affects all subsequent root tagged types declared in the declaration sequence, and specifies that no stream routines be generated. The form with an argument (for which there is also a corresponding aspect) specifies a single root tagged type for which stream routines are not to be generated.

Once the pragma has been given for a particular root tagged type, all subtypes and derived types of this type inherit the pragma automatically, so the effect applies to a complete hierarchy (this is necessary to deal with the class-wide dispatching versions of the stream routines).

When pragmas Discard_Names and No_Tagged_Streams are simultaneously applied to a tagged type its Expanded_Name and External_Tag are initialized with empty strings. This is useful to avoid exposing entity names at binary level but has a negative impact on the debuggability of tagged types.

### 2.116 Pragma Normalize_Scalars

Syntax:

```
pragma Normalize_Scalars;
```

This is a language defined pragma which is fully implemented in GNAT. The effect is to cause all scalar objects that are not otherwise initialized to be initialized. The initial values are implementation dependent and are as follows:
‘Standard.Character’
Objects whose root type is Standard.Character are initialized to Character’Last unless the subtype range excludes NUL (in which case NUL is used). This choice will always generate an invalid value if one exists.

‘Standard.Wide_Character’
Objects whose root type is Standard.Wide_Character are initialized to Wide_Character’Last unless the subtype range excludes NUL (in which case NUL is used). This choice will always generate an invalid value if one exists.

‘Standard.Wide_Wide_Character’
Objects whose root type is Standard.Wide_Wide_Character are initialized to the invalid value 16#FFFF_FFFF# unless the subtype range excludes NUL (in which case NUL is used). This choice will always generate an invalid value if one exists.

‘Integer types’
Objects of an integer type are treated differently depending on whether negative values are present in the subtype. If no negative values are present, then all one bits is used as the initial value except in the special case where zero is excluded from the subtype, in which case all zero bits are used. This choice will always generate an invalid value if one exists.

For subtypes with negative values present, the largest negative number is used, except in the unusual case where this largest negative number is in the subtype, and the largest positive number is not, in which case the largest positive value is used. This choice will always generate an invalid value if one exists.

‘Floating-Point Types’
Objects of all floating-point types are initialized to all 1-bits. For standard IEEE format, this corresponds to a NaN (not a number) which is indeed an invalid value.

‘Fixed-Point Types’
Objects of all fixed-point types are treated as described above for integers, with the rules applying to the underlying integer value used to represent the fixed-point value.

‘Modular types’
Objects of a modular type are initialized to all one bits, except in the special case where zero is excluded from the subtype, in which case all zero bits are used. This choice will always generate an invalid value if one exists.

‘Enumeration types’
Objects of an enumeration type are initialized to all one-bits, i.e., to the value 2 ** typ’Size - 1 unless the subtype excludes the literal whose Pos value is zero, in which case a code of zero is used. This choice will always generate an invalid value if one exists.

2.117 Pragma Obsolescent

Syntax:
This pragma can occur immediately following a declaration of an entity, including the case of a record component. If no Entity argument is present, then this declaration is the one to which the pragma applies. If an Entity parameter is present, it must either match the name of the entity in this declaration, or alternatively, the pragma can immediately follow an enumeration type declaration, where the Entity argument names one of the enumeration literals.

This pragma is used to indicate that the named entity is considered obsolescent and should not be used. Typically this is used when an API must be modified by eventually removing or modifying existing subprograms or other entities. The pragma can be used at an intermediate stage when the entity is still present, but will be removed later.

The effect of this pragma is to output a warning message on a reference to an entity thus marked that the subprogram is obsolescent if the appropriate warning option in the compiler is activated. If the Message parameter is present, then a second warning message is given containing this text. In addition, a reference to the entity is considered to be a violation of pragma Restrictions (No_Obsolescent_Features).

This pragma can also be used as a program unit pragma for a package, in which case the entity name is the name of the package, and the pragma indicates that the entire package is considered obsolescent. In this case a client withing such a package violates the restriction, and the with clause is flagged with warnings if the warning option is set.

If the Version parameter is present (which must be exactly the identifier Ada_05, no other argument is allowed), then the indication of obsolescence applies only when compiling in Ada 2005 mode. This is primarily intended for dealing with the situations in the predefined library where subprograms or packages have become defined as obsolescent in Ada 2005 (e.g., in Ada.Characters.Handling), but may be used anywhere.

The following examples show typical uses of this pragma:

```ada
package p is
  pragma Obsolescent (p, Message => "use pp instead of p");
end p;

package q is
  procedure q2;
  pragma Obsolescent ("use q2new instead");

  type R is new integer;
pragma Obsolescent
```

(Entity => R,
  Message => "use RR in Ada 2005",
  Version => Ada_05);

type M is record
  F1 : Integer;
  F2 : Integer;
  pragma Obsolescent;
  F3 : Integer;
end record;

type E is (a, bc, 'd', quack);
pragma Obsolescent (Entity => bc)
pragma Obsolescent (Entity => 'd')

function "+
  (a, b : character) return character;
pragma Obsolescent (Entity => "+");
end;

Note that, as for all pragmas, if you use a pragma argument identifier, then all subsequent
parameters must also use a pragma argument identifier. So if you specify Entity => for the
Entity argument, and a Message argument is present, it must be preceded by Message =>.

2.118 Pragma Optimize_Alignment

Syntax:

pragma Optimize_Alignment (TIME | SPACE | OFF);

This is a configuration pragma which affects the choice of default alignments for types
and objects where no alignment is explicitly specified. There is a time/space trade-off
in the selection of these values. Large alignments result in more efficient code, at the
expense of larger data space, since sizes have to be increased to match these alignments.
Smaller alignments save space, but the access code is slower. The normal choice of default
alignments for types and individual alignment promotions for objects (which is what you
get if you do not use this pragma, or if you use an argument of OFF), tries to balance these
two requirements.

Specifying SPACE causes smaller default alignments to be chosen in two cases. First any
packed record is given an alignment of 1. Second, if a size is given for the type, then the
alignment is chosen to avoid increasing this size. For example, consider:

type R is record
  X : Integer;
  Y : Character;
end record;

for R'Size use 5*8;

In the default mode, this type gets an alignment of 4, so that access to the Integer field X
are efficient. But this means that objects of the type end up with a size of 8 bytes. This
is a valid choice, since sizes of objects are allowed to be bigger than the size of the type, but it can waste space if for example fields of type R appear in an enclosing record. If the above type is compiled in `Optimize_Alignment (Space)` mode, the alignment is set to 1. However, there is one case in which SPACE is ignored. If a variable length record (that is a discriminated record with a component which is an array whose length depends on a discriminant), has a pragma Pack, then it is not in general possible to set the alignment of such a record to one, so the pragma is ignored in this case (with a warning).

Specifying SPACE also disables alignment promotions for standalone objects, which occur when the compiler increases the alignment of a specific object without changing the alignment of its type.

Specifying SPACE also disables component reordering in unpacked record types, which can result in larger sizes in order to meet alignment requirements.

Specifying TIME causes larger default alignments to be chosen in the case of small types with sizes that are not a power of 2. For example, consider:

```ada
type R is record
    A : Character;
    B : Character;
    C : Boolean;
end record;
pragma Pack (R);
for R'Size use 17;
```

The default alignment for this record is normally 1, but if this type is compiled in `Optimize_Alignment (Time)` mode, then the alignment is set to 4, which wastes space for objects of the type, since they are now 4 bytes long, but results in more efficient access when the whole record is referenced.

As noted above, this is a configuration pragma, and there is a requirement that all units in a partition be compiled with a consistent setting of the optimization setting. This would normally be achieved by use of a configuration pragma file containing the appropriate setting. The exception to this rule is that units with an explicit configuration pragma in the same file as the source unit are excluded from the consistency check, as are all predefined units. The latter are compiled by default in pragma Optimize_Alignment (Off) mode if no pragma appears at the start of the file.

### 2.119 Pragma Ordered

Syntax:

```
pragma Ordered (enumeration_first_subtype_LOCAL_NAME);
```

Most enumeration types are from a conceptual point of view unordered. For example, consider:

```ada
type Color is (Red, Blue, Green, Yellow);
```

By Ada semantics `Blue > Red` and `Green > Blue`, but really these relations make no sense; the enumeration type merely specifies a set of possible colors, and the order is unimportant.

For unordered enumeration types, it is generally a good idea if clients avoid comparisons (other than equality or inequality) and explicit ranges. (A ‘client’ is a unit where the type
is referenced, other than the unit where the type is declared, its body, and its subunits.)
For example, if code buried in some client says:

```pascal
if Current_Color < Yellow then ...
if Current_Color in Blue .. Green then ...
```
then the client code is relying on the order, which is undesirable. It makes the code hard to
read and creates maintenance difficulties if entries have to be added to the enumeration type.
Instead, the code in the client should list the possibilities, or an appropriate subtype should
be declared in the unit that declares the original enumeration type. E.g., the following
subtype could be declared along with the type Color:

```pascal
subtype RBG is Color range Red .. Green;
```
and then the client could write:

```pascal
if Current_Color in RBG then ...
if Current_Color = Blue or Current_Color = Green then ...
```
However, some enumeration types are legitimately ordered from a conceptual point of view.
For example, if you declare:

```pascal
type Day is (Mon, Tue, Wed, Thu, Fri, Sat, Sun);
```
then the ordering imposed by the language is reasonable, and clients can depend on it,
writing for example:

```pascal
if D in Mon .. Fri then ...
if D < Wed then ...
```
The pragma ‘Ordered’ is provided to mark enumeration types that are conceptually ordered,
alerting the reader that clients may depend on the ordering. GNAT provides a pragma to
mark enumerations as ordered rather than one to mark them as unordered, since in our
experience, the great majority of enumeration types are conceptually unordered.
The types Boolean, Character, Wide_Character, and Wide_Wide_Character are consid-
ered to be ordered types, so each is declared with a pragma Ordered in package Standard.
Normally pragma Ordered serves only as documentation and a guide for coding standards,
but GNAT provides a warning switch ‘-gnatw.u’ that requests warnings for inappropriate
uses (comparisons and explicit subranges) for unordered types. If this switch is used, then
any enumeration type not marked with pragma Ordered will be considered as unordered,
and will generate warnings for inappropriate uses.
Note that generic types are not considered ordered or unordered (since the template can be
instantiated for both cases), so we never generate warnings for the case of generic enumer-
ated types.
For additional information please refer to the description of the ‘-gnatw.u’ switch in the
GNAT User’s Guide.

### 2.120 Pragma Overflow_Mode

Syntax:

```pascal
pragma Overflow_Mode
  (  [General =>] MODE
    [, [Assertions =>] MODE]);
```
MODE ::= STRICT | MINIMIZED | ELIMINATED

This pragma sets the current overflow mode to the given setting. For details of the meaning of these modes, please refer to the ‘Overflow Check Handling in GNAT’ appendix in the GNAT User’s Guide. If only the General parameter is present, the given mode applies to all expressions. If both parameters are present, the General mode applies to expressions outside assertions, and the Eliminated mode applies to expressions within assertions.

The case of the MODE parameter is ignored, so MINIMIZED, Minimized and minimized all have the same effect.

The Overflow_Mode pragma has the same scoping and placement rules as pragma Suppress, so it can occur either as a configuration pragma, specifying a default for the whole program, or in a declarative scope, where it applies to the remaining declarations and statements in that scope.

The pragma Suppress (Overflow_Check) suppresses overflow checking, but does not affect the overflow mode.

The pragma Unsuppress (Overflow_Check) unsuppresses (enables) overflow checking, but does not affect the overflow mode.

2.121 Pragma Overriding_Renamings

Syntax:

```
pragma Overriding_Renamings;
```

This is a GNAT configuration pragma to simplify porting legacy code accepted by the Rational Ada compiler. In the presence of this pragma, a renaming declaration that renames an inherited operation declared in the same scope is legal if selected notation is used as in:

```
pragma Overriding_Renamings;
...
package R is
  function F (..);
  ...
  function F (..) renames R.F;
end R;
```

even though RM 8.3 (15) stipulates that an overridden operation is not visible within the declaration of the overriding operation.

2.122 Pragma Part_Of

Syntax:

```
pragma Part_Of (ABSTRACT_STATE);
```

```
ABSTRACT_STATE ::= NAME
```

For the semantics of this pragma, see the entry for aspect Part_Of in the SPARK 2014 Reference Manual, section 7.2.6.
2.123 Pragma Partition_Elaboration_Policy
Syntax:
pragma Partition_Elaboration_Policy (POLICY_IDENTIFIER);

POLICY_IDENTIFIER ::= Concurrent | Sequential

This pragma is standard in Ada 2005, but is available in all earlier versions of Ada as an implementation-defined pragma. See Ada 2012 Reference Manual for details.

2.124 Pragma Passive
Syntax:
pragma Passive [(Semaphore | No)];

Syntax checked, but otherwise ignored by GNAT. This is recognized for compatibility with DEC Ada 83 implementations, where it is used within a task definition to request that a task be made passive. If the argument Semaphore is present, or the argument is omitted, then DEC Ada 83 treats the pragma as an assertion that the containing task is passive and that optimization of context switch with this task is permitted and desired. If the argument No is present, the task must not be optimized. GNAT does not attempt to optimize any tasks in this manner (since protected objects are available in place of passive tasks).

For more information on the subject of passive tasks, see the section ‘Passive Task Optimization’ in the GNAT Users Guide.

2.125 Pragma Persistent_BSS
Syntax:
pragma Persistent_BSS [(LOCAL_NAME)]

This pragma allows selected objects to be placed in the .persistent_bss section. On some targets the linker and loader provide for special treatment of this section, allowing a program to be reloaded without affecting the contents of this data (hence the name persistent).

There are two forms of usage. If an argument is given, it must be the local name of a library-level object, with no explicit initialization and whose type is potentially persistent. If no argument is given, then the pragma is a configuration pragma, and applies to all library-level objects with no explicit initialization of potentially persistent types.

A potentially persistent type is a scalar type, or an untagged, non-discriminated record, all of whose components have no explicit initialization and are themselves of a potentially persistent type, or an array, all of whose constraints are static, and whose component type is potentially persistent.

If this pragma is used on a target where this feature is not supported, then the pragma will be ignored. See also pragma Linker_Section.

2.126 Pragma Post
Syntax:
pragma Post (Boolean_Expression);
The Post pragma is intended to be an exact replacement for the language-defined Post aspect, and shares its restrictions and semantics. It must appear either immediately following the corresponding subprogram declaration (only other pragmas may intervene), or if there is no separate subprogram declaration, then it can appear at the start of the declarations in a subprogram body (preceded only by other pragmas).

### 2.127 Pragma Postcondition

Syntax:

```plaintext
pragma Postcondition (
    [Check =>] Boolean_Expression
    [, [Message =>] String_Expression]);
```

The Postcondition pragma allows specification of automatic postcondition checks for subprograms. These checks are similar to assertions, but are automatically inserted just prior to the return statements of the subprogram with which they are associated (including implicit returns at the end of procedure bodies and associated exception handlers).

In addition, the boolean expression which is the condition which must be true may contain references to function Result in the case of a function to refer to the returned value.

Postcondition pragmas may appear either immediately following the (separate) declaration of a subprogram, or at the start of the declarations of a subprogram body. Only other pragmas may intervene (that is appear between the subprogram declaration and its postconditions, or appear before the postcondition in the declaration sequence in a subprogram body). In the case of a postcondition appearing after a subprogram declaration, the formal arguments of the subprogram are visible, and can be referenced in the postcondition expressions.

The postconditions are collected and automatically tested just before any return (implicit or explicit) in the subprogram body. A postcondition is only recognized if postconditions are active at the time the pragma is encountered. The compiler switch ‘gnata’ turns on all postconditions by default, and pragma Check_Policy with an identifier of Postcondition can also be used to control whether postconditions are active.

The general approach is that postconditions are placed in the spec if they represent functional aspects which make sense to the client. For example we might have:

```plaintext
function Direction return Integer;
pragma Postcondition
    (Direction'Result = +1
    or else
        Direction'Result = -1);
```

which serves to document that the result must be +1 or -1, and will test that this is the case at run time if postcondition checking is active.

Postconditions within the subprogram body can be used to check that some internal aspect of the implementation, not visible to the client, is operating as expected. For instance if a square root routine keeps an internal counter of the number of times it is called, then we might have the following postcondition:

```plaintext
Sqrt_Calls : Natural := 0;
```
function Sqrt (Arg : Float) return Float is
  pragma Postcondition
  (Sqrt_Calls = Sqrt_Calls'Old + 1);
  ...
end Sqrt

As this example shows, the use of the Old attribute is often useful in postconditions to refer to the state on entry to the subprogram.

Note that postconditions are only checked on normal returns from the subprogram. If an abnormal return results from raising an exception, then the postconditions are not checked. If a postcondition fails, then the exception System.Assertions.Assert_Failure is raised. If a message argument was supplied, then the given string will be used as the exception message. If no message argument was supplied, then the default message has the form “Postcondition failed at file_name:line”. The exception is raised in the context of the subprogram body, so it is possible to catch postcondition failures within the subprogram body itself.

Within a package spec, normal visibility rules in Ada would prevent forward references within a postcondition pragma to functions defined later in the same package. This would introduce undesirable ordering constraints. To avoid this problem, all postcondition pragmas are analyzed at the end of the package spec, allowing forward references.

The following example shows that this even allows mutually recursive postconditions as in:

package Parity_Functions is
  function Odd (X : Natural) return Boolean;
  pragma Postcondition
  (Odd'Result =
   (x = 1
    or else
    (x /= 0 and then Even (X - 1))));

  function Even (X : Natural) return Boolean;
  pragma Postcondition
  (Even'Result =
   (x = 0
    or else
    (x /= 1 and then Odd (X - 1))));
end Parity_Functions;

There are no restrictions on the complexity or form of conditions used within Postcondition pragmas. The following example shows that it is even possible to verify performance behavior.

package Sort is

  Performance : constant Float;
  -- Performance constant set by implementation
  -- to match target architecture behavior.
procedure Treesort (Arg : String);
   -- Sorts characters of argument using N*\log N sort
pragma Postcondition
   (Float (Clock - Clock'Old) <=
    Float (Arg'Length) *
    log (Float (Arg'Length)) *
    Performance);
end Sort;

Note: postcondition pragmas associated with subprograms that are marked as
Inline_Always, or those marked as Inline with front-end inlining (-gnatN option set)
are accepted and legality-checked by the compiler, but are ignored at run-time even if
postcondition checking is enabled.

Note that pragma Postcondition differs from the language-defined Post aspect (and cor-
responding Post pragma) in allowing multiple occurrences, allowing occurences in the body
even if there is a separate spec, and allowing a second string parameter, and the use of the
pragma identifier Check. Historically, pragma Postcondition was implemented prior to
the development of Ada 2012, and has been retained in its original form for compatibility
purposes.

2.128 Pragma Post_Class

Syntax:

    pragma Post_Class (Boolean_Expression);

The Post_Class pragma is intended to be an exact replacement for the language-defined
Post'Class aspect, and shares its restrictions and semantics. It must appear either im-
mmediately following the corresponding subprogram declaration (only other pragmas may
intervene), or if there is no separate subprogram declaration, then it can appear at the start
of the declarations in a subprogram body (preceded only by other pragmas).

Note: This pragma is called Post_Class rather than Post'Class because the latter would
not be strictly conforming to the allowed syntax for pragmas. The motivation for providing
pragmas equivalent to the aspects is to allow a program to be written using the pragmas,
and then compiled if necessary using an Ada compiler that does not recognize the pragmas
or aspects, but is prepared to ignore the pragmas. The assertion policy that controls this
pragma is Post'Class, not Post_Class.

2.129 Pragma Pre

Syntax:

    pragma Pre (Boolean_Expression);

The Pre pragma is intended to be an exact replacement for the language-defined Pre aspect,
and shares its restrictions and semantics. It must appear either immediately following the
corresponding subprogram declaration (only other pragmas may intervene), or if there is
no separate subprogram declaration, then it can appear at the start of the declarations in
a subprogram body (preceded only by other pragmas).
2.130 Pragma Precondition

Syntax:

```
pragma Precondition (  
    [Check =>] Boolean_Expression  
    [, [Message =>] String_Expression]);
```

The Precondition pragma is similar to Postcondition except that the corresponding checks take place immediately upon entry to the subprogram, and if a precondition fails, the exception is raised in the context of the caller, and the attribute ‘Result cannot be used within the precondition expression.

Otherwise, the placement and visibility rules are identical to those described for postconditions. The following is an example of use within a package spec:

```
package Math_Functions is  
  ...  
  function Sqrt (Arg : Float) return Float;  
  pragma Precondition (Arg >= 0.0)  
  ...  
end Math_Functions;
```

Precondition pragmas may appear either immediately following the (separate) declaration of a subprogram, or at the start of the declarations of a subprogram body. Only other pragmas may intervene (that is appear between the subprogram declaration and its postconditions, or appear before the postcondition in the declaration sequence in a subprogram body).

Note: precondition pragmas associated with subprograms that are marked as Inline_Always, or those marked as Inline with front-end inlining (-gnatN option set) are accepted and legality-checked by the compiler, but are ignored at run-time even if precondition checking is enabled.

Note that pragma Precondition differs from the language-defined Pre aspect (and corresponding Pre pragma) in allowing multiple occurrences, allowing occurrences in the body even if there is a separate spec, and allowing a second string parameter, and the use of the pragma identifier Check. Historically, pragma Precondition was implemented prior to the development of Ada 2012, and has been retained in its original form for compatibility purposes.

2.131 Pragma Predicate

Syntax:

```
pragma Predicate  
    ([Entity =>] type_LOCAL_NAME,  
    [Check =>] EXPRESSION);
```

This pragma (available in all versions of Ada in GNAT) encompasses both the Static_Predicate and Dynamic_Predicate aspects in Ada 2012. A predicate is regarded as static if it has an allowed form for Static_Predicate and is otherwise treated as a Dynamic_Predicate. Otherwise, predicates specified by this pragma behave exactly as described in the Ada 2012 reference manual. For example, if we have
type R is range 1 .. 10;
subtype S is R;
pragma Predicate (Entity => S, Check => S not in 4 .. 6);
subtype Q is R
pragma Predicate (Entity => Q, Check => F(Q) or G(Q));

the effect is identical to the following Ada 2012 code:

type R is range 1 .. 10;
subtype S is R with
  Static_Predicate => S not in 4 .. 6;
subtype Q is R with
  Dynamic_Predicate => F(Q) or G(Q);

Note that there are no pragmas Dynamic_Predicate or Static_Predicate. That is because these pragmas would affect legality and semantics of the program and thus do not have a neutral effect if ignored. The motivation behind providing pragmas equivalent to corresponding aspects is to allow a program to be written using the pragmas, and then compiled with a compiler that will ignore the pragmas. That doesn’t work in the case of static and dynamic predicates, since if the corresponding pragmas are ignored, then the behavior of the program is fundamentally changed (for example a membership test A in B would not take into account a predicate defined for subtype B). When following this approach, the use of predicates should be avoided.

2.132 Pragma Predicate_Failure

Syntax:

pragma Predicate_Failure ([Entity =>] type_LOCAL_NAME, [Message =>] String_Expression);

The Predicate_Failure pragma is intended to be an exact replacement for the language-defined Predicate_Failure aspect, and shares its restrictions and semantics.

2.133 Pragma Preelaborable_Initialization

Syntax:

pragma Preelaborable_Initialization (DIRECT_NAME);

This pragma is standard in Ada 2005, but is available in all earlier versions of Ada as an implementation-defined pragma. See Ada 2012 Reference Manual for details.

2.134 Pragma Prefix_Exception_Messages

Syntax:

pragma Prefix_Exception_Messages;

This is an implementation-defined configuration pragma that affects the behavior of raise statements with a message given as a static string constant (typically a string literal). In such cases, the string will be automatically prefixed by the name of the enclosing entity (giving the package and subprogram containing the raise statement). This helps to identify where messages are coming from, and this mode is automatic for the run-time library.
The pragma has no effect if the message is computed with an expression other than a static string constant, since the assumption in this case is that the program computes exactly the string it wants. If you still want the prefixing in this case, you can always call GNAT.Source_Info.Enclosing_Entity and prepend the string manually.

2.135 Pragma Pre_Class
Syntax:

    pragma Pre_Class (Boolean_Expression);

The Pre_Class pragma is intended to be an exact replacement for the language-defined Pre’Class aspect, and shares its restrictions and semantics. It must appear either immediately following the corresponding subprogram declaration (only other pragmas may intervene), or if there is no separate subprogram declaration, then it can appear at the start of the declarations in a subprogram body (preceded only by other pragmas).

Note: This pragma is called Pre_Class rather than Pre’Class because the latter would not be strictly conforming to the allowed syntax for pragmas. The motivation for providing pragmas equivalent to the aspects is to allow a program to be written using the pragmas, and then compiled if necessary using an Ada compiler that does not recognize the pragmas or aspects, but is prepared to ignore the pragmas. The assertion policy that controls this pragma is Pre’Class, not Pre_Class.

2.136 Pragma Priority_Specific_Dispatching
Syntax:

    pragma Priority_Specific_Dispatching ( 
      POLICY_IDENTIFIER, 
      first_priority_EXPRESSION, 
      last_priority_EXPRESSION)

    POLICY_IDENTIFIER ::= 
      EDF_Across_Priorities | 
      FIFO_Within_Priorities | 
      Non_Preemptive_Within_Priorities | 
      Round_Robin_Within_Priorities

This pragma is standard in Ada 2005, but is available in all earlier versions of Ada as an implementation-defined pragma. See Ada 2012 Reference Manual for details.

2.137 Pragma Profile
Syntax:

    pragma Profile (Ravenscar | Restricted | Rational | Jorvik | 
                   GNAT_Extended_Ravenscar | GNAT_Ravenscar_EDF );

This pragma is standard in Ada 2005, but is available in all earlier versions of Ada as an implementation-defined pragma. This is a configuration pragma that establishes a set of configuration pragmas that depend on the argument. Ravenscar is standard in Ada 2005. Jorvik is standard in Ada 202x. The other possibilities (Restricted,
Rational, GNAT_Extended_Ravenscar, GNAT_Ravenscar_EDF) are implementation-defined. GNAT_Extended_Ravenscar is an alias for Jorvik.

The set of configuration pragmas is defined in the following sections.

* Pragma Profile (Ravenscar)

The Ravenscar profile is standard in Ada 2005, but is available in all earlier versions of Ada as an implementation-defined pragma. This profile establishes the following set of configuration pragmas:

* Task_Dispatching_Policy (FIFO_Within_Priorities)
  [RM D.2.2] Tasks are dispatched following a preemptive priority-ordered scheduling policy.

* Locking_Policy (Ceiling_Locking)
  [RM D.3] While tasks and interrupts execute a protected action, they inherit the ceiling priority of the corresponding protected object.

* Detect_Blocking
  This pragma forces the detection of potentially blocking operations within a protected operation, and to raise Program_Error if that happens.

plus the following set of restrictions:

* Max_Entry_Queue_Length => 1
  No task can be queued on a protected entry.

* Max_Protected_Entries => 1
* Max_Task_Entries => 0
  No rendezvous statements are allowed.

* No_Abort_Statements
* No_Dynamic_Attachment
* No_Dynamic_Priorities
* No_Implicit_Heap_Allocations
* No_Local_Protected_Objects
* No_Local_Timing_Events
* No_Protected_Type_Allocators
* No_Relative_Delay
* No_Requeue_Statements
* No_Select_Statements
* No_Specific_Termination_Handlers
* No_Task_Allocators
* No_Task_Hierarchy
* No_Task_Termination
* Simple_Barriers

The Ravenscar profile also includes the following restrictions that specify that there are no semantic dependencies on the corresponding predefined packages:
* No_Dependence => Ada.Asynchronous_Task_Control  
* No_Dependence => Ada.Calendar  
* No_Dependence => Ada.Execution_Time.Group_Budget  
* No_Dependence => Ada.Execution_Time.Timers  
* No_Dependence => Ada.Task_Attributes  
* No_Dependence => System.Multiprocessors.Dispatching_Domains

This set of configuration pragmas and restrictions correspond to the definition of the ‘Ravenscar Profile’ for limited tasking, devised and published by the International Real-Time Ada Workshop, 1997. A description is also available at http://www-users.cs.york.ac.uk/~burns/ravenscar.ps.

The original definition of the profile was revised at subsequent IRTAW meetings. It has been included in the ISO Guide for the Use of the Ada Programming Language in High Integrity Systems, and was made part of the Ada 2005 standard. The formal definition given by the Ada Rapporteur Group (ARG) can be found in two Ada Issues (AI-249 and AI-305) available at http://www-ada-auth.org/cgi-bin/cvsweb.cgi/ais/ai-00249.txt and http://www-ada-auth.org/cgi-bin/cvsweb.cgi/ais/ai-00305.txt.

The above set is a superset of the restrictions provided by pragma Profile (Restricted), it includes six additional restrictions (Simple_Barriers, No_Select_Statements, No_Calendar, No_Implicit_Heap_Allocations, No_Relative_Delay and No_Task_Termination). This means that pragma Profile (Ravenscar), like the pragma Profile (Restricted), automatically causes the use of a simplified, more efficient version of the tasking run-time library.

* Pragma Profile (Jorvik)

Jorvik is the new profile added to the Ada 202x draft standard, previously implemented under the name GNAT_Extended_Ravenscar.

The No_Implicit_Heap_Allocations restriction has been replaced by No_Implicit_Task_Allocations and No_Implicit_Protected_Object_Allocations.

The Simple_Barriers restriction has been replaced by Pure_Barriers.

The Max_Protected_Entries, Max_Entry_Queue_Length, and No_Relative_Delay restrictions have been removed.

Details on the rationale for Jorvik and implications for use may be found in A New Ravenscar-Based Profile by P. Rogers, J. Ruiz, T. Gingold and P. Bernardi, in Reliable Software Technologies – Ada Europe 2017, Springer-Verlag Lecture Notes in Computer Science, Number 10300.

* Pragma Profile (GNAT_Ravenscar_EDF)

This profile corresponds to the Ravenscar profile but using EDF_Across_Priority as the Task_Scheduling_Policy.

* Pragma Profile (Restricted)

This profile corresponds to the GNAT restricted run time. It establishes the following set of restrictions:

  * No_Abort_Statements
* No_Entry_Queue
* No_Task_Hierarchy
* No_Task_Allocators
* No_DYNAMIC_PRIORITIES
* No_TERMINATE_ALTERNATIVES
* No_DYNAMIC_ATTACHMENT
* No_PROTECTED_TYPE_ALLOCATORS
* No_LOCAL_PROTECTED_OBJECTS
* No_Requeue_Statements
* No_Task_Attributes_Package
* Max_Asynchronous_Select_Nesting = 0
* Max_Task_Entries = 0
* Max_Protected_Entries = 1
* Max_Select_Alternatives = 0

This set of restrictions causes the automatic selection of a simplified version of the run

time that provides improved performance for the limited set of tasking functionality

permitted by this set of restrictions.

* Pragma Profile (Rational)
The Rational profile is intended to facilitate porting legacy code that compiles with the

Rational APEX compiler, even when the code includes non-conforming Ada constructs.
The profile enables the following three pragmas:

* pragma Implicit_Packing
* pragma Overriding_Renamings
* pragma Use_VADS_Size

2.138 Pragma Profile_Warnings

Syntax:

pragma Profile_Warnings (Ravenscar | Restricted | Rational);

This is an implementation-defined pragma that is similar in effect to pragma Profile except

that instead of generating Restrictions pragmas, it generates Restriction_Warnings

pragmas. The result is that violations of the profile generate warning messages instead of

error messages.

2.139 Pragma Propagate_Exceptions

Syntax:

pragma Propagate_Exceptions;

This pragma is now obsolete and, other than generating a warning if warnings on obsolescent

features are enabled, is ignored. It is retained for compatibility purposes. It used to be

used in connection with optimization of a now-obsolete mechanism for implementation of

exceptions.
2.140 Pragma Provide_Shift_Operators

Syntax:

```plaintext
pragma Provide_Shift_Operators (integer_first_subtype_LOCAL_NAME);
```

This pragma can be applied to a first subtype local name that specifies either an unsigned or signed type. It has the effect of providing the five shift operators (Shift_Left, Shift_Right, Shift_Right_Arithmetic, Rotate_Left and Rotate_Right) for the given type. It is similar to including the function declarations for these five operators, together with the pragma Import (Intrinsic, . . .) statements.

2.141 Pragma Psect_Object

Syntax:

```plaintext
pragma Psect_Object ( [Internal =>] LOCAL_NAME,
    [, [External =>] EXTERNAL_SYMBOL]
    [, [Size =>] EXTERNAL_SYMBOL]);
```

```plaintext
EXTERNAL_SYMBOL ::= IDENTIFIER
    | static_string_EXPRESSION
```

This pragma is identical in effect to pragma Common_Object.

2.142 Pragma Pure_Function

Syntax:

```plaintext
pragma Pure_Function ([Entity =>] function_LOCAL_NAME);
```

This pragma appears in the same declarative part as a function declaration (or a set of function declarations if more than one overloaded declaration exists, in which case the pragma applies to all entities). It specifies that the function Entity is to be considered pure for the purposes of code generation. This means that the compiler can assume that there are no side effects, and in particular that two identical calls produce the same result in the same context. It also means that the function can be used in an address clause.

Note that, quite deliberately, there are no static checks to try to ensure that this promise is met, so Pure_Function can be used with functions that are conceptually pure, even if they do modify global variables. For example, a square root function that is instrumented to count the number of times it is called is still conceptually pure, and can still be optimized, even though it modifies a global variable (the count). Memo functions are another example (where a table of previous calls is kept and consulted to avoid re-computation).

Note also that the normal rules excluding optimization of subprograms in pure units (when parameter types are descended from System.Address, or when the full view of a parameter type is limited), do not apply for the Pure_Function case. If you explicitly specify Pure_Function, the compiler may optimize away calls with identical arguments, and if that results in unexpected behavior, the proper action is not to use the pragma for subprograms that are not (conceptually) pure.

Note: Most functions in a Pure package are automatically pure, and there is no need to use pragma Pure_Function for such functions. One exception is any function that has at
least one formal of type `System.Address` or a type derived from it. Such functions are not considered pure by default, since the compiler assumes that the `Address` parameter may be functioning as a pointer and that the referenced data may change even if the address value does not. Similarly, imported functions are not considered to be pure by default, since there is no way of checking that they are in fact pure. The use of pragma `Pure_Function` for such a function will override these default assumption, and cause the compiler to treat a designated subprogram as pure in these cases.

Note: If pragma `Pure_Function` is applied to a renamed function, it applies to the underlying renamed function. This can be used to disambiguate cases of overloading where some but not all functions in a set of overloaded functions are to be designated as pure.

If pragma `Pure_Function` is applied to a library-level function, the function is also considered pure from an optimization point of view, but the unit is not a Pure unit in the categorization sense. So for example, a function thus marked is free to `with` non-pure units.

2.143 Pragma Rational

Syntax:

```
pragma Rational;
```

This pragma is considered obsolescent, but is retained for compatibility purposes. It is equivalent to:

```
pragma Profile (Rational);
```

2.144 Pragma Ravenscar

Syntax:

```
pragma Ravenscar;
```

This pragma is considered obsolescent, but is retained for compatibility purposes. It is equivalent to:

```
pragma Profile (Ravenscar);
```

which is the preferred method of setting the `Ravenscar` profile.

2.145 Pragma Refined_Depends

Syntax:

```
pragma Refined_Depends (DEPENDENCY_RELATION);
```

```
DEPENDENCY_RELATION ::= 
    null 
  | (DEPENDENCY_CLAUSE {, DEPENDENCY_CLAUSE})

DEPENDENCY_CLAUSE ::= 
    OUTPUT_LIST =>[+] INPUT_LIST 
  | NULL_DEPENDENCY_CLAUSE

NULL_DEPENDENCY_CLAUSE ::= null => INPUT_LIST
```
Chapter 2: Implementation Defined Pragmas

\[
\text{OUTPUT\_LIST ::= OUTPUT | (OUTPUT \{, OUTPUT\})}
\]

\[
\text{INPUT\_LIST ::= null | INPUT | (INPUT \{, INPUT\})}
\]

\[
\text{OUTPUT ::= NAME | FUNCTION\_RESULT}
\]

\[
\text{INPUT ::= NAME}
\]

where \text{FUNCTION\_RESULT} is a function \text{Result attribute\_reference}

For the semantics of this pragma, see the entry for aspect \text{Refined\_Depends} in the SPARK 2014 Reference Manual, section 6.1.5.

2.146 Pragma Refined\_Global

Syntax:

\[
\text{pragma Refined\_Global (GLOBAL\_SPECIFICATION);}\]

\[
\text{GLOBAL\_SPECIFICATION ::=}
\]

\[
\null | (GLOBAL\_LIST)
\]

\[
| (MODED\_GLOBAL\_LIST \{, MODED\_GLOBAL\_LIST\})
\]

\[
\text{MODED\_GLOBAL\_LIST ::= MODE\_SELECTOR => GLOBAL\_LIST}
\]

\[
\text{MODE\_SELECTOR ::= In\_Out | Input | Output | Proof\_In}
\]

\[
\text{GLOBAL\_LIST ::= GLOBAL\_ITEM | (GLOBAL\_ITEM \{, GLOBAL\_ITEM\})}
\]

\[
\text{GLOBAL\_ITEM ::= NAME}
\]

For the semantics of this pragma, see the entry for aspect \text{Refined\_Global} in the SPARK 2014 Reference Manual, section 6.1.4.

2.147 Pragma Refined\_Post

Syntax:

\[
\text{pragma Refined\_Post (boolean\_EXPRESSION);}\]

For the semantics of this pragma, see the entry for aspect \text{Refined\_Post} in the SPARK 2014 Reference Manual, section 7.2.7.

2.148 Pragma Refined\_State

Syntax:

\[
\text{pragma Refined\_State (REFINEMENT\_LIST);}\]

\[
\text{REFINEMENT\_LIST ::=}
\]

\[
(REFINEMENT\_CLAUSE \{, REFINEMENT\_CLAUSE\})
\]

\[
\text{REFINEMENT\_CLAUSE ::= state\_NAME => CONSTITUENT\_LIST}
\]

\[
\text{CONSTITUENT\_LIST ::=}
\]
null
  | CONSTITUENT
  | (CONSTITUENT {, CONSTITUENT})

CONSTITUENT ::= object_NAME | state_NAME

For the semantics of this pragma, see the entry for aspect Refined_State in the SPARK 2014 Reference Manual, section 7.2.2.

2.149 Pragma Relative_Deadline

Syntax:

pragma Relative_Deadline (time_span_EXPRESSION);

This pragma is standard in Ada 2005, but is available in all earlier versions of Ada as an implementation-defined pragma. See Ada 2012 Reference Manual for details.

2.150 Pragma Remote_Access_Type

Syntax:

pragma Remote_Access_Type ([Entity =>] formal_access_type_LOCAL_NAME);

This pragma appears in the formal part of a generic declaration. It specifies an exception to the RM rule from E.2.2(17/2), which forbids the use of a remote access to class-wide type as actual for a formal access type.

When this pragma applies to a formal access type Entity, that type is treated as a remote access to class-wide type in the generic. It must be a formal general access type, and its designated type must be the class-wide type of a formal tagged limited private type from the same generic declaration.

In the generic unit, the formal type is subject to all restrictions pertaining to remote access to class-wide types. At instantiation, the actual type must be a remote access to class-wide type.

2.151 Pragma Rename_Pragma

Syntax:

pragma Rename_Pragma ( [New_Name =>] IDENTIFIER,
                         [Renamed =>] pragma_IDENTIFIER);

This pragma provides a mechanism for supplying new names for existing pragmas. The New_Name identifier can subsequently be used as a synonym for the Renamed pragma. For example, suppose you have code that was originally developed on a compiler that supports Inline_Only as an implementation defined pragma. And suppose the semantics of pragma Inline_Only are identical to (or at least very similar to) the GNAT implementation defined pragma Inline_Always. You could globally replace Inline_Only with Inline_Always.

However, to avoid that source modification, you could instead add a configuration pragma:

pragma Rename_Pragma ( New_Name => Inline_Only,
                         Renamed => Inline_Always);
Then GNAT will treat "pragma Inline Only . . ." as if you had written "pragma In-
line Always . . .". Pragma Inline Only will not necessarily mean the same thing as the other Ada compiler; it's up to you to make sure the semantics are close enough.

### 2.152 Pragma Restricted_Run_Time

Syntax:

```
pragma Restricted_Run_Time;
```

This pragma is considered obsolescent, but is retained for compatibility purposes. It is equivalent to:

```
pragma Profile (Restricted);
```

which is the preferred method of setting the restricted run time profile.

### 2.153 Pragma Restriction_Warnings

Syntax:

```
pragma Restriction_Warnings

(restriction_IDENTIFIER {, restriction_IDENTIFIER});
```

This pragma allows a series of restriction identifiers to be specified (the list of allowed identifiers is the same as for pragma Restrictions). For each of these identifiers the compiler checks for violations of the restriction, but generates a warning message rather than an error message if the restriction is violated.

One use of this is in situations where you want to know about violations of a restriction, but you want to ignore some of these violations. Consider this example, where you want to set Ada_95 mode and enable style checks, but you want to know about any other use of implementation pragmas:

```
pragma Restriction_Warnings (No_Implementation_Pragmas);
pragma Warnings (Off, "violation of No_Implementation_Pragmas");
pragma Ada_95;
pragma Style_Checks ("2bfhkm160");
pragma Warnings (On, "violation of No_Implementation_Pragmas");
```

By including the above lines in a configuration pragmas file, the Ada_95 and Style_Checks pragmas are accepted without generating a warning, but any other use of implementation defined pragmas will cause a warning to be generated.

### 2.154 Pragma Reviewable

Syntax:

```
pragma Reviewable;
```

This pragma is an RM-defined standard pragma, but has no effect on the program being compiled, or on the code generated for the program.

To obtain the required output specified in RM H.3.1, the compiler must be run with various special switches as follows:
* ‘Where compiler-generated run-time checks remain’
  The switch ‘-gnatGL’ may be used to list the expanded code in pseudo-Ada form. Runtime checks show up in the listing either as explicit checks or operators marked with {} to indicate a check is present.

* ‘An identification of known exceptions at compile time’
  If the program is compiled with ‘-gnatwa’, the compiler warning messages will indicate all cases where the compiler detects that an exception is certain to occur at run time.

* ‘Possible reads of uninitialized variables’
  The compiler warns of many such cases, but its output is incomplete.

A supplemental static analysis tool may be used to obtain a comprehensive list of all possible points at which uninitialized data may be read.

* ‘Where run-time support routines are implicitly invoked’
  In the output from ‘-gnatGL’, run-time calls are explicitly listed as calls to the relevant run-time routine.

* ‘Object code listing’
  This may be obtained either by using the ‘-S’ switch, or the objdump utility.

* ‘Constructs known to be erroneous at compile time’
  These are identified by warnings issued by the compiler (use ‘-gnatwa’).

* ‘Stack usage information’
  Static stack usage data (maximum per-subprogram) can be obtained via the ‘-fstack-usage’ switch to the compiler. Dynamic stack usage data (per task) can be obtained via the ‘-u’ switch to gnatbind.

* ‘Object code listing of entire partition’
  This can be obtained by compiling the partition with ‘-S’, or by applying objdump to all the object files that are part of the partition.

* ‘A description of the run-time model’
  The full sources of the run-time are available, and the documentation of these routines describes how these run-time routines interface to the underlying operating system facilities.

* ‘Control and data-flow information’
  A supplemental static analysis tool may be used to obtain complete control and data-flow information, as well as comprehensive messages identifying possible problems based on this information.

### 2.155 Pragma Secondary_Stack_Size

**Syntax:**

```ada
pragma Secondary_Stack_Size (integer_EXPRESSION);
```

This pragma appears within the task definition of a single task declaration or a task type declaration (like pragma Storage_Size) and applies to all task objects of that type. The argument specifies the size of the secondary stack to be used by these task objects, and
must be of an integer type. The secondary stack is used to handle functions that return a variable-sized result, for example a function returning an unconstrained String.

Note this pragma only applies to targets using fixed secondary stacks, like VxWorks 653 and bare board targets, where a fixed block for the secondary stack is allocated from the primary stack of the task. By default, these targets assign a percentage of the primary stack for the secondary stack, as defined by `System.Parameter.Sec_Stack_Percentage`. With this pragma, an `integer_EXPRESSION` of bytes is assigned from the primary stack instead.

For most targets, the pragma does not apply as the secondary stack grows on demand: allocated as a chain of blocks in the heap. The default size of these blocks can be modified via the `-D` binder option as described in *GNAT User's Guide*.

Note that no check is made to see if the secondary stack can fit inside the primary stack.

Note the pragma cannot appear when the restriction `No_Secondary_Stack` is in effect.

### 2.156 Pragma Share_Generic

Syntax:

```
pragma Share_Generic (GNAME {, GNAME});
```

```
GNAME ::= generic_unit_NAME | generic_instance_NAME
```

This pragma is provided for compatibility with Dec Ada 83. It has no effect in GNAT (which does not implement shared generics), other than to check that the given names are all names of generic units or generic instances.

### 2.157 Pragma Shared

This pragma is provided for compatibility with Ada 83. The syntax and semantics are identical to pragma Atomic.

### 2.158 Pragma Short_Circuit_And_Or

Syntax:

```
pragma Short_Circuit_And_Or;
```

This configuration pragma causes any occurrence of the AND operator applied to operands of type `Standard.Boolean` to be short-circuited (i.e. the AND operator is treated as if it were `AND THEN`). Or is similarly treated as `OR ELSE`. This may be useful in the context of certification protocols requiring the use of short-circuited logical operators. If this configuration pragma occurs locally within the file being compiled, it applies only to the file being compiled. There is no requirement that all units in a partition use this option.

### 2.159 Pragma Short_Descriptors

Syntax:

```
pragma Short_Descriptors;
```

This pragma is provided for compatibility with other Ada implementations. It is recognized but ignored by all current versions of GNAT.
2.160 Pragma Simple_Storage_Pool_Type

Syntax:

```
pragma Simple_Storage_Pool_Type (type_LOCAL_NAME);
```

A type can be established as a ‘simple storage pool type’ by applying the representation pragma `Simple_Storage_Pool_Type` to the type. A type named in the pragma must be a library-level immutably limited record type or limited tagged type declared immediately within a package declaration. The type can also be a limited private type whose full type is allowed as a simple storage pool type.

For a simple storage pool type `SSP`, nonabstract primitive subprograms `Allocate`, `Deallocate`, and `Storage_Size` can be declared that are subtype conformant with the following subprogram declarations:

- **Procedure Allocate**
  ```
  procedure Allocate
  (Pool : in out SSP;
   Storage_Address : out System.Address;
   Size_In_Storage_Elements : System.Storage_Elements.Storage_Count;
   Alignment : System.Storage_Elements.Storage_Count);
  ```

- **Procedure Deallocate**
  ```
  procedure Deallocate
  (Pool : in out SSP;
   Storage_Address : System.Address;
   Size_In_Storage_Elements : System.Storage_Elements.Storage_Count;
   Alignment : System.Storage_Elements.Storage_Count);
  ```

- **Function Storage_Size**
  ```
  function Storage_Size (Pool : SSP)
  return System.Storage_Elements.Storage_Count;
  ```

Procedure `Allocate` must be declared, whereas `Deallocate` and `Storage_Size` are optional. If `Deallocate` is not declared, then applying an unchecked deallocation has no effect other than to set its actual parameter to null. If `Storage_Size` is not declared, then the `Storage_Size` attribute applied to an access type associated with a pool object of type `SSP` returns zero. Additional operations can be declared for a simple storage pool type (such as for supporting a mark/release storage-management discipline).

An object of a simple storage pool type can be associated with an access type by specifying the attribute [Simple_Storage_Pool], page 130. For example:

```
My_Pool : My_Simple_Storage_Pool_Type;

type Acc is access My_Data_Type;

for Acc’Simple_Storage_Pool use My_Pool;
```

See attribute [Simple_Storage_Pool], page 130 for further details.

2.161 Pragma Source_File_Name

Syntax:

```
pragma Source_File_Name (unit_NAME =>] unit_NAME,
```
Spec_File_Name => STRING_LITERAL,
[Index => INTEGER_LITERAL]);

pragma Source_File_Name (  
[Unit_Name =>] unit_NAME,  
Body_File_Name => STRING_LITERAL,  
[Index => INTEGER_LITERAL]);

Use this to override the normal naming convention. It is a configuration pragma, and so has the usual applicability of configuration pragmas (i.e., it applies to either an entire partition, or to all units in a compilation, or to a single unit, depending on how it is used). unit_name is mapped to file_name_literal. The identifier for the second argument is required, and indicates whether this is the file name for the spec or for the body.

The optional Index argument should be used when a file contains multiple units, and when you do not want to use gnatchop to separate them into multiple files (which is the recommended procedure to limit the number of recompilations that are needed when some sources change). For instance, if the source file source.ada contains

package B is
...
end B;

with B;
procedure A is
begin
...
end A;

you could use the following configuration pragmas:

pragma Source_File_Name
(B, Spec_File_Name => "source.ada", Index => 1);
pragma Source_File_Name
(A, Body_File_Name => "source.ada", Index => 2);

Note that the gnatname utility can also be used to generate those configuration pragmas.

Another form of the Source_File_Name pragma allows the specification of patterns defining alternative file naming schemes to apply to all files.

pragma Source_File_Name
(  
[Spec_File_Name =>] STRING_LITERAL  
[, [Casing =>] CASING_SPEC]  
[, [Dot_Replacement =>] STRING_LITERAL]);

pragma Source_File_Name
(  
[Body_File_Name =>] STRING_LITERAL  
[, [Casing =>] CASING_SPEC]  
[, [Dot_Replacement =>] STRING_LITERAL]);

pragma Source_File_Name
(  
[Subunit_File_Name =>] STRING_LITERAL  
[, [Casing =>] CASING_SPEC]  
[, [Dot_Replacement =>] STRING_LITERAL]);
Casing => CASING_SPEC,
[Dot_Replacement => STRING_LITERAL]);

CASING_SPEC ::= Lowercase | Uppercase | Mixedcase

The first argument is a pattern that contains a single asterisk indicating the point at which the unit name is to be inserted in the pattern string to form the file name. The second argument is optional. If present it specifies the casing of the unit name in the resulting file name string. The default is lower case. Finally the third argument allows for systematic replacement of any dots in the unit name by the specified string literal.

Note that Source_File_Name pragmas should not be used if you are using project files. The reason for this rule is that the project manager is not aware of these pragmas, and so other tools that use the project file would not be aware of the intended naming conventions. If you are using project files, file naming is controlled by Source_File_Name_Project pragmas, which are usually supplied automatically by the project manager. A pragma Source_File_Name cannot appear after a [Pragma Source_File_Name_Project], page 87.

For more details on the use of the Source_File_Name pragma, see the sections on Using Other File Names and Alternative File Naming Schemes in the GNAT User's Guide.

2.162 Pragma Source_File_Name_Project

This pragma has the same syntax and semantics as pragma Source_File_Name. It is only allowed as a stand-alone configuration pragma. It cannot appear after a [Pragma Source_File_Name], page 85, and most importantly, once pragma Source_File_Name_Project appears, no further Source_File_Name pragmas are allowed.

The intention is that Source_File_Name_Project pragmas are always generated by the Project Manager in a manner consistent with the naming specified in a project file, and when naming is controlled in this manner, it is not permissible to attempt to modify this naming scheme using Source_File_Name or Source_File_Name_Project pragmas (which would not be known to the project manager).

2.163 Pragma Source_Reference

Syntax:
pragma Source_Reference (INTEGER_LITERAL, STRING_LITERAL);

This pragma must appear as the first line of a source file. integer_literal is the logical line number of the line following the pragma line (for use in error messages and debugging information). string_literal is a static string constant that specifies the file name to be used in error messages and debugging information. This is most notably used for the output of gnatchop with the ‘-r’ switch, to make sure that the original unchopped source file is the one referred to.

The second argument must be a string literal, it cannot be a static string expression other than a string literal. This is because its value is needed for error messages issued by all phases of the compiler.

2.164 Pragma SPARK_Mode

Syntax:
pragma SPARK_Mode [(On | Off)] ;

In general a program can have some parts that are in SPARK 2014 (and follow all the rules in the SPARK Reference Manual), and some parts that are full Ada 2012. The SPARK_Mode pragma is used to identify which parts are in SPARK 2014 (by default programs are in full Ada). The SPARK_Mode pragma can be used in the following places:

* As a configuration pragma, in which case it sets the default mode for all units compiled with this pragma.
* Immediately following a library-level subprogram spec
* Immediately within a library-level package body
* Immediately following the private keyword of a library-level package spec
* Immediately following the begin keyword of a library-level package body
* Immediately within a library-level subprogram body

Normally a subprogram or package spec/body inherits the current mode that is active at the point it is declared. But this can be overridden by pragma within the spec or body as above.

The basic consistency rule is that you can’t turn SPARK_Mode back On, once you have explicitly (with a pragma) turned it Off. So the following rules apply:

If a subprogram spec has SPARK_Mode Off, then the body must also have SPARK_Mode Off.

For a package, we have four parts:

* the package public declarations
* the package private part
* the body of the package
* the elaboration code after begin

For a package, the rule is that if you explicitly turn SPARK_Mode Off for any part, then all the following parts must have SPARK_Mode Off. Note that this may require repeating a pragma SPARK_Mode (Off) in the body. For example, if we have a configuration pragma SPARK_Mode (On) that turns the mode on by default everywhere, and one particular package spec has pragma SPARK_Mode (Off), then that pragma will need to be repeated in the package body.

2.165 Pragma Static_Elaboration_Desired

Syntax:

pragma Static_Elaboration_Desired;

This pragma is used to indicate that the compiler should attempt to initialize statically the objects declared in the library unit to which the pragma applies, when these objects are initialized (explicitly or implicitly) by an aggregate. In the absence of this pragma, aggregates in object declarations are expanded into assignments and loops, even when the aggregate components are static constants. When the aggregate is present the compiler builds a static expression that requires no run-time code, so that the initialized object can be placed in read-only data space. If the components are not static, or the aggregate has more than 100 components, the compiler emits a warning that the pragma cannot be
obeyed. (See also the restriction No_Implicit_Loops, which supports static construction of larger aggregates with static components that include an others choice.)

2.166 Pragma Stream_Convert

Syntax:

```plaintext
pragma Stream_Convert (  
    [Entity =>] type_LOCAL_NAME,  
    [Read  =>] function_NAME,  
    [Write  =>] function_NAME);
```

This pragma provides an efficient way of providing user-defined stream attributes. Not only is it simpler to use than specifying the attributes directly, but more importantly, it allows the specification to be made in such a way that the predefined unit Ada.Streams is not loaded unless it is actually needed (i.e. unless the stream attributes are actually used); the use of the Stream_Convert pragma adds no overhead at all, unless the stream attributes are actually used on the designated type.

The first argument specifies the type for which stream functions are provided. The second parameter provides a function used to read values of this type. It must name a function whose argument type may be any subtype, and whose returned type must be the type given as the first argument to the pragma.

The meaning of the Read parameter is that if a stream attribute directly or indirectly specifies reading of the type given as the first parameter, then a value of the type given as the argument to the Read function is read from the stream, and then the Read function is used to convert this to the required target type.

Similarly the Write parameter specifies how to treat write attributes that directly or indirectly apply to the type given as the first parameter. It must have an input parameter of the type specified by the first parameter, and the return type must be the same as the input type of the Read function. The effect is to first call the Write function to convert to the given stream type, and then write the result type to the stream.

The Read and Write functions must not be overloaded subprograms. If necessary renamings can be supplied to meet this requirement. The usage of this attribute is best illustrated by a simple example, taken from the GNAT implementation of package Ada.Strings.Unbounded:

```plaintext
function To_Unbounded (S : String) return Unbounded_String  
renames To_Unbounded_String;

pragma Stream_Convert  
    (Unbounded_String, To_Unbounded, To_String);
```

The specifications of the referenced functions, as given in the Ada Reference Manual are:

```plaintext
function To_Unbounded_String (Source : String)  
    return Unbounded_String;

function To_String (Source : Unbounded_String)  
    return String;
```

The effect is that if the value of an unbounded string is written to a stream, then the representation of the item in the stream is in the same format that would be used for
Standard.String’Output, and this same representation is expected when a value of this
type is read from the stream. Note that the value written always includes the bounds, even
for Unbounded_String’Write, since Unbounded_String is not an array type.
Note that the Stream_Convert pragma is not effective in the case of a derived type of a
non-limited tagged type. If such a type is specified then the pragma is silently ignored, and
the default implementation of the stream attributes is used instead.

2.167 Pragma Style_Checks

Syntax:

pragma Style_Checks (string_LITERAL | ALL_CHECKS |
On | Off [, LOCAL_NAME]);

This pragma is used in conjunction with compiler switches to control the built in style
checking provided by GNAT. The compiler switches, if set, provide an initial setting for
the switches, and this pragma may be used to modify these settings, or the settings may
be provided entirely by the use of the pragma. This pragma can be used anywhere that
a pragma is legal, including use as a configuration pragma (including use in the gnat.adc
file).

The form with a string literal specifies which style options are to be activated. These are
additive, so they apply in addition to any previously set style check options. The codes for
the options are the same as those used in the ‘-gnaty’ switch to ‘gcc’ or ‘gnatmake’. For
example the following two methods can be used to enable layout checking:

* 
  pragma Style_Checks ("l");
* 
  gcc -c -gnatyl ...

The form ALL_CHECKS activates all standard checks (its use is equivalent to the use of the
gnay switch with no options. See the GNAT User’s Guide for details.)

Note: the behavior is slightly different in GNAT mode (-gnatg used). In this case, ALL-
CHECKS implies the standard set of GNAT mode style check options (i.e. equivalent to
-gnatyg).

The forms with Off and On can be used to temporarily disable style checks as shown in the
following example:

    pragma Style_Checks ("k"); -- requires keywords in lower case
    pragma Style_Checks (Off); -- turn off style checks
    NULL; -- this will not generate an error message
    pragma Style_Checks (On); -- turn style checks back on
    NULL; -- this will generate an error message

Finally the two argument form is allowed only if the first argument is On or Off. The effect
is to turn of semantic style checks for the specified entity, as shown in the following example:

    pragma Style_Checks ("r"); -- require consistency of identifier casing
    Arg : Integer;
    Rf1 : Integer := ARG; -- incorrect, wrong case
    pragma Style_Checks (Off, Arg);
    Rf2 : Integer := ARG; -- OK, no error
2.168 Pragma Subtitle

Syntax:

    pragma Subtitle ([Subtitle =>] STRING_LITERAL);

This pragma is recognized for compatibility with other Ada compilers but is ignored by GNAT.

2.169 Pragma Suppress

Syntax:

    pragma Suppress (Identifier [, [On =>] Name]);

This is a standard pragma, and supports all the check names required in the RM. It is included here because GNAT recognizes some additional check names that are implementation defined (as permitted by the RM):

- **Alignment_Check** can be used to suppress alignment checks on addresses used in address clauses. Such checks can also be suppressed by suppressing range checks, but the specific use of *Alignment_Check* allows suppression of alignment checks without suppressing other range checks. Note that *Alignment_Check* is suppressed by default on machines (such as the x86) with non-strict alignment.

- **Atomic_Synchronization** can be used to suppress the special memory synchronization instructions that are normally generated for access to *Atomic* variables to ensure correct synchronization between tasks that use such variables for synchronization purposes.

- **Duplicated_Tag_Check** Can be used to suppress the check that is generated for a duplicated tag value when a tagged type is declared.

- **Container_Checks** Can be used to suppress all checks within Ada.Containers and instances of its children, including Tampering_Check.

- **Tampering_Check** Can be used to suppress tampering check in the containers.

- **Predicate_Check** can be used to control whether predicate checks are active. It is applicable only to predicates for which the policy is Check. Unlike *Assertion_Policy*, which determines if a given predicate is ignored or checked for the whole program, the use of *Suppress* and *Unsuppress* with this check name allows a given predicate to be turned on and off at specific points in the program.

- **Validity_Check** can be used specifically to control validity checks. If *Suppress* is used to suppress validity checks, then no validity checks are performed, including those specified by the appropriate compiler switch or the *Validity_Checks* pragma.

- Additional check names previously introduced by use of the *Check_Name* pragma are also allowed.

Note that pragma Suppress gives the compiler permission to omit checks, but does not require the compiler to omit checks. The compiler will generate checks if they are essentially free, even when they are suppressed. In particular, if the compiler can prove that a certain check will necessarily fail, it will generate code to do an unconditional ‘raise’, even if checks are suppressed. The compiler warns in this case.

Of course, run-time checks are omitted whenever the compiler can prove that they will not fail, whether or not checks are suppressed.
2.170 Pragma Suppress_All

Syntax:

    pragma Suppress_All;

This pragma can appear anywhere within a unit. The effect is to apply Suppress (All_Checks) to the unit in which it appears. This pragma is implemented for compatibility with DEC Ada 83 usage where it appears at the end of a unit, and for compatibility with Rational Ada, where it appears as a program unit pragma. The use of the standard Ada pragma Suppress (All_Checks) as a normal configuration pragma is the preferred usage in GNAT.

2.171 Pragma Suppress_Debug_Info

Syntax:

    pragma Suppress_Debug_Info ([Entity =>] LOCAL_NAME);

This pragma can be used to suppress generation of debug information for the specified entity. It is intended primarily for use in debugging the debugger, and navigating around debugger problems.

2.172 Pragma Suppress_Exception_Locations

Syntax:

    pragma Suppress_Exception_Locations;

In normal mode, a raise statement for an exception by default generates an exception message giving the file name and line number for the location of the raise. This is useful for debugging and logging purposes, but this entails extra space for the strings for the messages. The configuration pragma Suppress_Exception_Locations can be used to suppress the generation of these strings, with the result that space is saved, but the exception message for such raises is null. This configuration pragma may appear in a global configuration pragma file, or in a specific unit as usual. It is not required that this pragma be used consistently within a partition, so it is fine to have some units within a partition compiled with this pragma and others compiled in normal mode without it.

2.173 Pragma Suppress_INITIALIZATION

Syntax:

    pragma Suppress_INITIALIZATION ([Entity =>] variable_or_subtype_Name);

Here variable_or_subtype_Name is the name introduced by a type declaration or subtype declaration or the name of a variable introduced by an object declaration.

In the case of a type or subtype this pragma suppresses any implicit or explicit initialization for all variables of the given type or subtype, including initialization resulting from the use of pragmas Normalize_Scalars or Initialize_Scalars.

This is considered a representation item, so it cannot be given after the type is frozen. It applies to all subsequent object declarations, and also any allocator that creates objects of the type.
If the pragma is given for the first subtype, then it is considered to apply to the base type and all its subtypes. If the pragma is given for other than a first subtype, then it applies only to the given subtype. The pragma may not be given after the type is frozen.

Note that this includes eliminating initialization of discriminants for discriminated types, and tags for tagged types. In these cases, you will have to use some non-portable mechanism (e.g. address overlays or unchecked conversion) to achieve required initialization of these fields before accessing any object of the corresponding type.

For the variable case, implicit initialization for the named variable is suppressed, just as though its subtype had been given in a pragma Suppress Initialization, as described above.

### 2.174 Pragma Task_Name

**Syntax**

\[
\text{pragma Task\_Name (string\_EXPRESSION)};
\]

This pragma appears within a task definition (like pragma Priority) and applies to the task in which it appears. The argument must be of type String, and provides a name to be used for the task instance when the task is created. Note that this expression is not required to be static, and in particular, it can contain references to task discriminants. This facility can be used to provide different names for different tasks as they are created, as illustrated in the example below.

The task name is recorded internally in the run-time structures and is accessible to tools like the debugger. In addition the routine Ada.Task_Identification.Image will return this string, with a unique task address appended.

```
-- Example of the use of pragma Task_Name

with Ada.Task_Identification;
use Ada.Task_Identification;
with Text_IO; use Text_IO;
procedure t3 is

    type Astring is access String;

    task type Task\_Typ (Name : access String) is
        pragma Task\_Name (Name.all);
    end Task\_Typ;

    task body Task\_Typ is
        Nam : constant String := Image (Current_Task);
        begin
            Put\_Line ("-->", Nam (1 .. 14) & "<--");
        end Task\_Typ;

    type Ptr\_Task is access Task\_Typ;
    Task\_Var : Ptr\_Task;

    begin
```
Task_Var :=
    new Task_Typ (new String'("This is task 1"));
Task_Var :=
    new Task_Typ (new String'("This is task 2"));
end;

2.175 Pragma Task_Storage

Syntax:

pragma Task_Storage (  
    [Task_Type =>] LOCAL_NAME,  
    [Top_Guard =>] static_integer_EXPRESSION);  

This pragma specifies the length of the guard area for tasks. The guard area is an additional 
storage area allocated to a task. A value of zero means that either no guard area is created 
or a minimal guard area is created, depending on the target. This pragma can appear 
anywhere a Storage_Size attribute definition clause is allowed for a task type.

2.176 Pragma Test_Case

Syntax:

pragma Test_Case (  
    [Name =>] static_string_Expression  
    , [Mode =>] (Nominal | Robustness)  
    [, Requires => Boolean_Expression]  
    [, Ensures => Boolean_Expression]);

The Test_Case pragma allows defining fine-grain specifications for use by testing tools. 
The compiler checks the validity of the Test_Case pragma, but its presence does not lead 
to any modification of the code generated by the compiler.

Test_Case pragmas may only appear immediately following the (separate) declaration of a 
subprogram in a package declaration, inside a package spec unit. Only other pragmas may 
intervene (that is appear between the subprogram declaration and a test case).

The compiler checks that boolean expressions given in Requires and Ensures are valid, 
where the rules for Requires are the same as the rule for an expression in Precondition 
and the rules for Ensures are the same as the rule for an expression in Postcondition. In 
particular, attributes 'Old and 'Result can only be used within the Ensures expression. 
The following is an example of use within a package spec:

package Math_Functions is
...
    function Sqrt (Arg : Float) return Float;
    pragma Test_Case (Name => "Test 1",  
        Mode => Nominal,  
        Requires => Arg < 10000.0,  
        Ensures => Sqrt'Result < 10.0);  
...
end Math_Functions;
The meaning of a test case is that there is at least one context where \texttt{Requires} holds such that, if the associated subprogram is executed in that context, then \texttt{Ensures} holds when the subprogram returns. Mode \texttt{Nominal} indicates that the input context should also satisfy the precondition of the subprogram, and the output context should also satisfy its postcondition. Mode \texttt{Robustness} indicates that the precondition and postcondition of the subprogram should be ignored for this test case.

\subsection*{2.177 Pragma \texttt{Thread Local Storage}}

Syntax:
\begin{verbatim}
pragma Thread_Local_Storage ([Entity =>] LOCAL_NAME);
\end{verbatim}

This pragma specifies that the specified entity, which must be a variable declared in a library-level package, is to be marked as “Thread Local Storage” (TLS). On systems supporting this (which include Windows, Solaris, GNU/Linux, and VxWorks), this causes each thread (and hence each Ada task) to see a distinct copy of the variable. The variable must not have default initialization, and if there is an explicit initialization, it must be either \texttt{null} for an access variable, a static expression for a scalar variable, or a fully static aggregate for a composite type, that is to say, an aggregate all of whose components are static, and which does not include packed or discriminated components.

This provides a low-level mechanism similar to that provided by the \texttt{Ada.Task_Attributes} package, but much more efficient and is also useful in writing interface code that will interact with foreign threads.

If this pragma is used on a system where TLS is not supported, then an error message will be generated and the program will be rejected.

\subsection*{2.178 Pragma \texttt{Time Slice}}

Syntax:
\begin{verbatim}
pragma Time_Slice (static_duration_EXPRESSION);
\end{verbatim}

For implementations of GNAT on operating systems where it is possible to supply a time slice value, this pragma may be used for this purpose. It is ignored if it is used in a system that does not allow this control, or if it appears in other than the main program unit.

\subsection*{2.179 Pragma Title}

Syntax:
\begin{verbatim}
pragma Title (TITLING_OPTION [, TITLING OPTION]);
\end{verbatim}

\begin{verbatim}
TITLING_OPTION ::= [Title =>] STRING_LITERAL, | [Subtitle =>] STRING_LITERAL
\end{verbatim}

Syntax checked but otherwise ignored by GNAT. This is a listing control pragma used in DEC Ada 83 implementations to provide a title and/or subtitle for the program listing. The program listing generated by GNAT does not have titles or subtitles.

Unlike other pragmas, the full flexibility of named notation is allowed for this pragma, i.e., the parameters may be given in any order if named notation is used, and named and positional notation can be mixed following the normal rules for procedure calls in Ada.
2.180 Pragma Type_Invariant

Syntax:

```
pragma Type_Invariant
    ( [Entity =>] type_LOCAL_NAME,
    [Check =>] EXPRESSION);
```

The Type_Invariant pragma is intended to be an exact replacement for the language-defined Type_Invariant aspect, and shares its restrictions and semantics. It differs from the language defined Invariant pragma in that it does not permit a string parameter, and it is controlled by the assertion identifier Type_Invariant rather than Invariant.

2.181 Pragma Type_Invariant_Class

Syntax:

```
pragma Type_Invariant_Class
    ( [Entity =>] type_LOCAL_NAME,
    [Check =>] EXPRESSION);
```

The Type_Invariant_Class pragma is intended to be an exact replacement for the language-defined Type_Invariant’Class aspect, and shares its restrictions and semantics. Note: This pragma is called Type_Invariant_Class rather than Type_Invariant’Class because the latter would not be strictly conforming to the allowed syntax for pragmas. The motivation for providing pragmas equivalent to the aspects is to allow a program to be written using the pragmas, and then compiled if necessary using an Ada compiler that does not recognize the pragmas or aspects, but is prepared to ignore the pragmas. The assertion policy that controls this pragma is Type_Invariant’Class, not Type_Invariant_Class.

2.182 Pragma Unchecked_Union

Syntax:

```
pragma Unchecked_Union (first_subtype_LOCAL_NAME);
```

This pragma is used to specify a representation of a record type that is equivalent to a C union. It was introduced as a GNAT implementation defined pragma in the GNAT Ada 95 mode. Ada 2005 includes an extended version of this pragma, making it language defined, and GNAT fully implements this extended version in all language modes (Ada 83, Ada 95, and Ada 2005). For full details, consult the Ada 2012 Reference Manual, section B.3.3.

2.183 Pragma Unevaluated_Use_Of_Old

Syntax:

```
pragma Unevaluated_Use_Of_Old (Error | Warn | Allow);
```

This pragma controls the processing of attributes Old and Loop_Entry. If either of these attributes is used in a potentially unevaluated expression (e.g. the then or else parts of an if expression), then normally this usage is considered illegal if the prefix of the attribute is other than an entity name. The language requires this behavior for Old, and GNAT copies the same rule for Loop_Entry.
The reason for this rule is that otherwise, we can have a situation where we save the Old value, and this results in an exception, even though we might not evaluate the attribute. Consider this example:

```plaintext
class UnevalOld is
    K : Character;
    procedure U (A : String; C : Boolean) -- ERROR
        with Post => (if C then A(1)'Old = K else True);
    end;
end;
```

If procedure `U` is called with a string with a lower bound of 2, and `C` false, then an exception would be raised trying to evaluate `A(1)` on entry even though the value would not be actually used.

Although the rule guarantees against this possibility, it is sometimes too restrictive. For example if we know that the string has a lower bound of 1, then we will never raise an exception. The pragma `Unevaluated.Use_Of.Old` can be used to modify this behavior. If the argument is `Error` then an error is given (this is the default RM behavior). If the argument is `Warn` then the usage is allowed as legal but with a warning that an exception might be raised. If the argument is `Allow` then the usage is allowed as legal without generating a warning.

This pragma may appear as a configuration pragma, or in a declarative part or package specification. In the latter case it applies to uses up to the end of the corresponding statement sequence or sequence of package declarations.

### 2.184 Pragma Unimplemented_Unit

Syntax:

```plaintext
pragma Unimplemented_Unit;
```

If this pragma occurs in a unit that is processed by the compiler, GNAT aborts with the message `xxx not implemented`, where `xxx` is the name of the current compilation unit. This pragma is intended to allow the compiler to handle unimplemented library units in a clean manner.

The abort only happens if code is being generated. Thus you can use specs of unimplemented packages in syntax or semantic checking mode.

### 2.185 Pragma Universal_Aliasing

Syntax:

```plaintext
pragma Universal_Aliasing [([Entity =>] type_LOCAL_NAME)];
```

type_LOCAL_NAME must refer to a type declaration in the current declarative part. The effect is to inhibit strict type-based aliasing optimization for the given type. In other words, the effect is as though access types designating this type were subject to pragma `No.Strict.Aliasing`. For a detailed description of the strict aliasing optimization, and the situations in which it must be suppressed, see the section on `Optimization and Strict Aliasing` in the GNAT User's Guide.
2.186 Pragma Unmodified

Syntax:

\[ \text{pragma Unmodified (LOCAL\_NAME \{, LOCAL\_NAME\});} \]

This pragma signals that the assignable entities (variables, out parameters, in out parameters) whose names are listed are deliberately not assigned in the current source unit. This suppresses warnings about the entities being referenced but not assigned, and in addition a warning will be generated if one of these entities is in fact assigned in the same unit as the pragma (or in the corresponding body, or one of its subunits).

This is particularly useful for clearly signaling that a particular parameter is not modified, even though the spec suggests that it might be.

For the variable case, warnings are never given for unreferenced variables whose name contains one of the substrings DISCARD, DUMMY, IGNORE, JUNK, UNUSE, TMP, TEMP in any casing. Such names are typically to be used in cases where such warnings are expected. Thus it is never necessary to use \text{pragma Unmodified} for such variables, though it is harmless to do so.

2.187 Pragma Unreferenced

Syntax:

\[ \text{pragma Unreferenced (LOCAL\_NAME \{, LOCAL\_NAME\});} \]
\[ \text{pragma Unreferenced (library\_unit\_NAME \{, library\_unit\_NAME\});} \]

This pragma signals that the entities whose names are listed are deliberately not referenced in the current source unit after the occurrence of the pragma. This suppresses warnings about the entities being unreferenced, and in addition a warning will be generated if one of these entities is in fact subsequently referenced in the same unit as the pragma (or in the corresponding body, or one of its subunits).

This is particularly useful for clearly signaling that a particular parameter is not referenced in some particular subprogram implementation and that this is deliberate. It can also be useful in the case of objects declared only for their initialization or finalization side effects.

If \text{LOCAL\_NAME} identifies more than one matching homonym in the current scope, then the entity most recently declared is the one to which the pragma applies. Note that in the case of accept formals, the pragma Unreferenced may appear immediately after the keyword do which allows the indication of whether or not accept formals are referenced or not to be given individually for each accept statement.

The left hand side of an assignment does not count as a reference for the purpose of this pragma. Thus it is fine to assign to an entity for which pragma Unreferenced is given. However, use of an entity as an actual for an out parameter does count as a reference unless warnings for unread output parameters are enabled via \text{-gnatw.o}.

Note that if a warning is desired for all calls to a given subprogram, regardless of whether they occur in the same unit as the subprogram declaration, then this pragma should not be used (calls from another unit would not be flagged); pragma Obsolescent can be used instead for this purpose, see \cite{Pragma Obsolescent}, page 62.

The second form of pragma Unreferenced is used within a context clause. In this case the arguments must be unit names of units previously mentioned in \text{with} clauses (similar to the
usage of pragma Elaborate_All). The effect is to suppress warnings about unreferenced units and unreferenced entities within these units.

For the variable case, warnings are never given for unreferenced variables whose name contains one of the substrings DISCARD, DUMMY, IGNORE, JUNK, UNUSED in any casing. Such names are typically to be used in cases where such warnings are expected. Thus it is never necessary to use pragma Unreferenced for such variables, though it is harmless to do so.

### 2.188 Pragma Unreferenced_Objects

**Syntax:**

```plaintext
pragma Unreferenced_Objects (local_subtype_NAME {, local_subtype_NAME});
```

This pragma signals that for the types or subtypes whose names are listed, objects which are declared with one of these types or subtypes may not be referenced, and if no references appear, no warnings are given.

This is particularly useful for objects which are declared solely for their initialization and finalization effect. Such variables are sometimes referred to as RAII variables (Resource Acquisition Is Initialization). Using this pragma on the relevant type (most typically a limited controlled type), the compiler will automatically suppress unwanted warnings about these variables not being referenced.

### 2.189 Pragma Unreserve_All_Interrupts

**Syntax:**

```plaintext
pragma Unreserve_All_Interrupts;
```

Normally certain interrupts are reserved to the implementation. Any attempt to attach an interrupt causes Program_Error to be raised, as described in RM C.3.2(22). A typical example is the SIGINT interrupt used in many systems for a Ctrl-C interrupt. Normally this interrupt is reserved to the implementation, so that Ctrl-C can be used to interrupt execution.

If the pragma Unreserve_All_Interrupts appears anywhere in any unit in a program, then all such interrupts are unreserved. This allows the program to handle these interrupts, but disables their standard functions. For example, if this pragma is used, then pressing Ctrl-C will not automatically interrupt execution. However, a program can then handle the SIGINT interrupt as it chooses.

For a full list of the interrupts handled in a specific implementation, see the source code for the spec of Ada.Interrupts.Names in file a-intnam.ads. This is a target dependent file that contains the list of interrupts recognized for a given target. The documentation in this file also specifies what interrupts are affected by the use of the Unreserve_All_Interrupts pragma.

For a more general facility for controlling what interrupts can be handled, see pragma Interrupt_State, which subsumes the functionality of the Unreserve_All_Interrupts pragma.
2.190 Pragma Unsuppress

Syntax:

\[
\text{pragma Unsuppress (IDENTIFIER [, [On =>] NAME])};
\]

This pragma undoes the effect of a previous pragma \texttt{Suppress}. If there is no corresponding pragma \texttt{Suppress} in effect, it has no effect. The range of the effect is the same as for pragma \texttt{Suppress}. The meaning of the arguments is identical to that used in pragma \texttt{Suppress}.

One important application is to ensure that checks are on in cases where code depends on the checks for its correct functioning, so that the code will compile correctly even if the compiler switches are set to suppress checks. For example, in a program that depends on external names of tagged types and wants to ensure that the duplicated tag check occurs even if all run-time checks are suppressed by a compiler switch, the following configuration pragma will ensure this test is not suppressed:

\[
\text{pragma Unsuppress (Duplicated_Tag_Check)};
\]

This pragma is standard in Ada 2005. It is available in all earlier versions of Ada as an implementation-defined pragma.

Note that in addition to the checks defined in the Ada RM, GNAT recognizes a number of implementation-defined check names. See the description of pragma \texttt{Suppress} for full details.

2.191 Pragma Unused

Syntax:

\[
\text{pragma Unused (LOCAL_NAME {, LOCAL_NAME})};
\]

This pragma signals that the assignable entities (variables, \texttt{out} parameters, and \texttt{in out} parameters) whose names are listed deliberately do not get assigned or referenced in the current source unit after the occurrence of the pragma in the current source unit. This suppresses warnings about the entities that are unreferenced and/or not assigned, and, in addition, a warning will be generated if one of these entities gets assigned or subsequently referenced in the same unit as the pragma (in the corresponding body or one of its subunits).

This is particularly useful for clearly signaling that a particular parameter is not modified or referenced, even though the spec suggests that it might be.

For the variable case, warnings are never given for unreferenced variables whose name contains one of the substrings \texttt{DISCARD, DUMMY, IGNORE, JUNK, UNUSED} in any casing. Such names are typically to be used in cases where such warnings are expected. Thus it is never necessary to use \texttt{pragma Unused} for such variables, though it is harmless to do so.

2.192 Pragma Use_VADS_Size

Syntax:

\[
\text{pragma Use_VADS_Size};
\]

This is a configuration pragma. In a unit to which it applies, any use of the ‘Size attribute is automatically interpreted as a use of the ‘VADS_Size attribute. Note that this may result in incorrect semantic processing of valid Ada 95 or Ada 2005 programs. This is intended to aid in the handling of existing code which depends on the interpretation of Size as implemented in the VADS compiler. See description of the VADS_Size attribute for further details.
2.193 Pragma Validity_Checks

Syntax:

    pragma Validity_Checks (string_LITERAL | ALL_CHECKS | On | Off);

This pragma is used in conjunction with compiler switches to control the built-in validity checking provided by GNAT. The compiler switches, if set provide an initial setting for the switches, and this pragma may be used to modify these settings, or the settings may be provided entirely by the use of the pragma. This pragma can be used anywhere that a pragma is legal, including use as a configuration pragma (including use in the gnat.adc file).

The form with a string literal specifies which validity options are to be activated. The validity checks are first set to include only the default reference manual settings, and then a string of letters in the string specifies the exact set of options required. The form of this string is exactly as described for the `-gnatVx` compiler switch (see the GNAT User’s Guide for details). For example the following two methods can be used to enable validity checking for mode in and in out subprogram parameters:

*  
  
    pragma Validity_Checks ("im");

*  
    $ gcc -c -gnatVim ...

The form ALL_CHECKS activates all standard checks (its use is equivalent to the use of the gnatVa switch).

The forms with Off and On can be used to temporarily disable validity checks as shown in the following example:

    pragma Validity_Checks ("c"); -- validity checks for copies
    pragma Validity_Checks (Off); -- turn off validity checks
    A := B; -- B will not be validity checked
    pragma Validity_Checks (On); -- turn validity checks back on
    A := C; -- C will be validity checked

2.194 Pragma Volatile

Syntax:

    pragma Volatile (LOCAL_NAME);

This pragma is defined by the Ada Reference Manual, and the GNAT implementation is fully conformant with this definition. The reason it is mentioned in this section is that a pragma of the same name was supplied in some Ada 83 compilers, including DEC Ada 83. The Ada 95 / Ada 2005 implementation of pragma Volatile is upwards compatible with the implementation in DEC Ada 83.

2.195 Pragma Volatile_Full_Access

Syntax:

    pragma Volatile_Full_Access (LOCAL_NAME);
This is similar in effect to pragma Volatile, except that any reference to the object is guaranteed to be done only with instructions that read or write all the bits of the object. Furthermore, if the object is of a composite type, then any reference to a subcomponent of the object is guaranteed to read and/or write all the bits of the object.

The intention is that this be suitable for use with memory-mapped I/O devices on some machines. Note that there are two important respects in which this is different from pragma Atomic. First a reference to a Volatile_Full_Access object is not a sequential action in the RM 9.10 sense and, therefore, does not create a synchronization point. Second, in the case of pragma Atomic, there is no guarantee that all the bits will be accessed if the reference is not to the whole object; the compiler is allowed (and generally will) access only part of the object in this case.

2.196 Pragma Volatile_Function

Syntax:

```plaintext
pragma Volatile_Function [ (static_boolean_EXPRESSION) ];
```

For the semantics of this pragma, see the entry for aspect Volatile_Function in the SPARK 2014 Reference Manual, section 7.1.2.

2.197 Pragma Warning_As_Error

Syntax:

```plaintext
pragma Warning_As_Error (static_string_EXPRESSION);
```

This configuration pragma allows the programmer to specify a set of warnings that will be treated as errors. Any warning that matches the pattern given by the pragma argument will be treated as an error. This gives more precise control than -gnatwe, which treats warnings as errors.

This pragma can apply to regular warnings (messages enabled by -gnatw) and to style warnings (messages that start with “(style)”, enabled by -gnaty).

The pattern may contain asterisks, which match zero or more characters in the message. For example, you can use `pragma Warning_As_Error ("bits of*unused")` to treat the warning message `warning: 960 bits of "a" unused` as an error. All characters other than asterisk are treated as literal characters in the match. The match is case insensitive; for example XYZ matches xyz.

Note that the pattern matches if it occurs anywhere within the warning message string (it is not necessary to put an asterisk at the start and the end of the message, since this is implied).

Another possibility for the static_string_EXPRESSION which works whether or not error tags are enabled (`-gnatw.d`) is to use a single `-gnatw` tag string, enclosed in brackets, as shown in the example below, to treat one category of warnings as errors. Note that if you want to treat multiple categories of warnings as errors, you can use multiple pragma Warning_As_Error.

The above use of patterns to match the message applies only to warning messages generated by the front end. This pragma can also be applied to warnings provided by the back end and mentioned in [Pragma Warnings], page 103. By using a single full `-Wxxx` switch in the pragma, such warnings can also be treated as errors.
The pragma can appear either in a global configuration pragma file (e.g. gnat.adc), or at the start of a file. Given a global configuration pragma file containing:

```
pragma Warning_As_Error ("[-gnatwj]");
```

which will treat all obsolescent feature warnings as errors, the following program compiles as shown (compile options here are `-gnatwa.d -gnatl -gnatj55').

1. `pragma Warning_As_Error ("*never assigned*")`;
2. function Warnerr return String is
   3.   X : Integer;
       |
       >>> error: variable "X" is never read and never assigned [-gnatwv] [warning-as-error]
   4.   Y : Integer;
       |
       >>> warning: variable "Y" is assigned but never read [-gnatwu]
   5. begin
   6.   Y := 0;
   7.   return "%ABC%";
       |
       >>> error: use of "%" is an obsolescent feature (RM J.2(4)), use "" instead [-gnatwj] [warning-as-error]
   8. end;

8 lines: No errors, 3 warnings (2 treated as errors)

Note that this pragma does not affect the set of warnings issued in any way, it merely changes the effect of a matching warning if one is produced as a result of other warnings options. As shown in this example, if the pragma results in a warning being treated as an error, the tag is changed from “warning:“ to “error:“ and the string “[warning-as-error]” is appended to the end of the message.

### 2.198 Pragma Warnings

Syntax:

```
pragma Warnings ([TOOL_NAME,] DETAILS [, REASON]);
```

- `DETAILS ::= On | Off`
- `DETAILS ::= On | Off, local_NAME`
- `DETAILS ::= static_string_EXPRESSION`
- `DETAILS ::= On | Off, static_string_EXPRESSION`

- `TOOL_NAME ::= GNAT | GNATprove`
REASON ::= Reason => STRING_LITERAL {& STRING_LITERAL}

Note: in Ada 83 mode, a string literal may be used in place of a static string expression (which does not exist in Ada 83).

Note if the second argument of DETAILS is a local_NAME then the second form is always understood. If the intention is to use the fourth form, then you can write NAME & "" to force the interpretation as a 'static_string_EXPRESSION'.

Note: if the first argument is a valid TOOL_NAME, it will be interpreted that way. The use of the TOOL_NAME argument is relevant only to users of SPARK and GNATprove, see last part of this section for details.

Normally warnings are enabled, with the output being controlled by the command line switch. Warnings (Off) turns off generation of warnings until a Warnings (On) is encountered or the end of the current unit. If generation of warnings is turned off using this pragma, then some or all of the warning messages are suppressed, regardless of the setting of the command line switches.

The Reason parameter may optionally appear as the last argument in any of the forms of this pragma. It is intended purely for the purposes of documenting the reason for the Warnings pragma. The compiler will check that the argument is a static string but otherwise ignore this argument. Other tools may provide specialized processing for this string.

The form with a single argument (or two arguments if Reason present), where the first argument is ON or OFF may be used as a configuration pragma.

If the LOCAL_NAME parameter is present, warnings are suppressed for the specified entity. This suppression is effective from the point where it occurs till the end of the extended scope of the variable (similar to the scope of Suppress). This form cannot be used as a configuration pragma.

In the case where the first argument is other than ON or OFF, the third form with a single static_string_EXPRESSION argument (and possible reason) provides more precise control over which warnings are active. The string is a list of letters specifying which warnings are to be activated and which deactivated. The code for these letters is the same as the string used in the command line switch controlling warnings. For a brief summary, use the gnatmake command with no arguments, which will generate usage information containing the list of warnings switches supported. For full details see the section on Warning Message Control in the GNAT User’s Guide. This form can also be used as a configuration pragma.

The warnings controlled by the -gnatw switch are generated by the front end of the compiler. The GCC back end can provide additional warnings and they are controlled by the -W switch. Such warnings can be identified by the appearance of a string of the form [-W{xxx}] in the message which designates the -W‘xxx’ switch that controls the message. The form with a single 'static_string_EXPRESSION' argument also works for these warnings, but the string must be a single full -W‘xxx’ switch in this case. The above reference lists a few examples of these additional warnings.

The specified warnings will be in effect until the end of the program or another pragma Warnings is encountered. The effect of the pragma is cumulative. Initially the set of warnings is the standard default set as possibly modified by compiler switches. Then each pragma Warning modifies this set of warnings as specified. This form of the pragma may also be used as a configuration pragma.
The fourth form, with an `On|Off` parameter and a string, is used to control individual messages, based on their text. The string argument is a pattern that is used to match against the text of individual warning messages (not including the initial “warning: ” tag). The pattern may contain asterisks, which match zero or more characters in the message. For example, you can use `pragma Warnings (Off, "bits of unused")` to suppress the warning message `warning: 960 bits of "a" unused`. No other regular expression notations are permitted. All characters other than asterisk in these three specific cases are treated as literal characters in the match. The match is case insensitive, for example `XYZ` matches `xyz`.

Note that the pattern matches if it occurs anywhere within the warning message string (it is not necessary to put an asterisk at the start and the end of the message, since this is implied).

The above use of patterns to match the message applies only to warning messages generated by the front end. This form of the pragma with a string argument can also be used to control warnings provided by the back end and mentioned above. By using a single full `-W'xxx'` switch in the pragma, such warnings can be turned on and off.

There are two ways to use the pragma in this form. The OFF form can be used as a configuration pragma. The effect is to suppress all warnings (if any) that match the pattern string throughout the compilation (or match the `-W` switch in the back end case).

The second usage is to suppress a warning locally, and in this case, two pragmas must appear in sequence:

```plaintext
pragma Warnings (Off, Pattern);
... code where given warning is to be suppressed
pragma Warnings (On, Pattern);
```

In this usage, the pattern string must match in the Off and On pragmas, and (if `-gnatw.w` is given) at least one matching warning must be suppressed.

Note: if the ON form is not found, then the effect of the OFF form extends until the end of the file (pragma Warnings is purely textual, so its effect does not stop at the end of the enclosing scope).

Note: to write a string that will match any warning, use the string "***". It will not work to use a single asterisk or two asterisks since this looks like an operator name. This form with three asterisks is similar in effect to specifying `pragma Warnings (Off)` except (if `-gnatw.w` is given) that a matching `pragma Warnings (On, "***")` will be required. This can be helpful in avoiding forgetting to turn warnings back on.

Note: the debug flag `-gnatd.i` can be used to cause the compiler to entirely ignore all WARNINGS pragmas. This can be useful in checking whether obsolete pragmas in existing programs are hiding real problems.

Note: pragma Warnings does not affect the processing of style messages. See separate entry for pragma Style_Checks for control of style messages.

Users of the formal verification tool GNATprove for the SPARK subset of Ada may use the version of the pragma with a `TOOL_NAME` parameter.

If present, `TOOL_NAME` is the name of a tool, currently either `GNAT` for the compiler or `GNATprove` for the formal verification tool. A given tool only takes into account pragma Warnings that do not specify a tool name, or that specify the matching tool name. This
makes it possible to disable warnings selectively for each tool, and as a consequence to
detect useless pragma Warnings with switch -gnatw.w.

2.199 Pragma Weak_External

Syntax:

\[
\text{pragma Weak_External ([Entity =>] LOCAL_NAME);}\\
\]

LOCAL_NAME must refer to an object that is declared at the library level. This pragma
specifies that the given entity should be marked as a weak symbol for the linker. It is
equivalent to \text{__attribute__((weak))} in GNU C and causes LOCAL_NAME to be emitted as
a weak symbol instead of a regular symbol, that is to say a symbol that does not have to
be resolved by the linker if used in conjunction with a pragma Import.

When a weak symbol is not resolved by the linker, its address is set to zero. This is useful in
writing interfaces to external modules that may or may not be linked in the final executable,
for example depending on configuration settings.

If a program references at run time an entity to which this pragma has been applied, and
the corresponding symbol was not resolved at link time, then the execution of the program
is erroneous. It is not erroneous to take the Address of such an entity, for example to guard
potential references, as shown in the example below.

Some file formats do not support weak symbols so not all target machines support this
pragma.

\[\text{-- Example of the use of pragma Weak_External}\]

\[
\begin{verbatim}
package External_Module is
key : Integer;
pragma Import (C, key);
pragma Weak_External (key);
function Present return boolean;
end External_Module;
\end{verbatim}
\]

with System; use System;
package body External_Module is
function Present return boolean is
begin
return key'Address /= System.Null_Address;
end Present;
end External_Module;

2.200 Pragma Wide_Character_Encoding

Syntax:

\[
\text{pragma Wide_Character_Encoding (IDENTIFIER | CHARACTER_LITERAL);}\\
\]

This pragma specifies the wide character encoding to be used in program source text ap-
pearing subsequently. It is a configuration pragma, but may also be used at any point that
a pragma is allowed, and it is permissible to have more than one such pragma in a file,
allowing multiple encodings to appear within the same file.
However, note that the pragma cannot immediately precede the relevant wide character, because then the previous encoding will still be in effect, causing “illegal character” errors. The argument can be an identifier or a character literal. In the identifier case, it is one of HEX, UPPER, SHIFT_JIS, EUC, UTF8, or BRACKETS. In the character literal case it is correspondingly one of the characters h, u, s, e, 8, or b.

Note that when the pragma is used within a file, it affects only the encoding within that file, and does not affect withed units, specs, or subunits.
3 Implementation Defined Aspects

Ada defines (throughout the Ada 2012 reference manual, summarized in Annex K) a set of aspects that can be specified for certain entities. These language defined aspects are implemented in GNAT in Ada 2012 mode and work as described in the Ada 2012 Reference Manual.

In addition, Ada 2012 allows implementations to define additional aspects whose meaning is defined by the implementation. GNAT provides a number of these implementation-defined aspects which can be used to extend and enhance the functionality of the compiler. This section of the GNAT reference manual describes these additional aspects.

Note that any program using these aspects may not be portable to other compilers (although GNAT implements this set of aspects on all platforms). Therefore if portability to other compilers is an important consideration, you should minimize the use of these aspects.

Note that for many of these aspects, the effect is essentially similar to the use of a pragma or attribute specification with the same name applied to the entity. For example, if we write:

```ada
type R is range 1 .. 100
  with Value_Size => 10;
```

then the effect is the same as:

```ada
type R is range 1 .. 100;
for R'Value_Size use 10;
```

and if we write:

```ada
type R is new Integer
  with Shared => True;
```

then the effect is the same as:

```ada
type R is new Integer;
pragma Shared (R);
```

In the documentation below, such cases are simply marked as being boolean aspects equivalent to the corresponding pragma or attribute definition clause.

3.1 Aspect Abstract_State

This aspect is equivalent to [pragma Abstract_State], page 5.

3.2 Aspect Annotate

There are three forms of this aspect (where ID is an identifier, and ARG is a general expression), corresponding to [pragma Annotate], page 9.

‘Annotate => ID’
Equivalent to `pragma Annotate (ID, Entity => Name);`

‘Annotate => (ID)’
Equivalent to `pragma Annotate (ID, Entity => Name);`

‘Annotate => (ID ,ID {, ARG})’
Equivalent to `pragma Annotate (ID, ID {, ARG}, Entity => Name);`
3.3 Aspect Async_Readers
This boolean aspect is equivalent to [pragma Async_Readers], page 13.

3.4 Aspect Async_Writers
This boolean aspect is equivalent to [pragma Async_Writers], page 13.

3.5 Aspect Constant_After_Elaboration
This aspect is equivalent to [pragma Constant_After_Elaboration], page 20.

3.6 Aspect Contract_Cases
This aspect is equivalent to [pragma Contract_Cases], page 20, the sequence of clauses being enclosed in parentheses so that syntactically it is an aggregate.

3.7 Aspect Depends
This aspect is equivalent to [pragma Depends], page 25.

3.8 Aspect Default_Initial_Condition
This aspect is equivalent to [pragma Default_Initial_Condition], page 24.

3.9 Aspect Dimension
The Dimension aspect is used to specify the dimensions of a given subtype of a dimensioned numeric type. The aspect also specifies a symbol used when doing formatted output of dimensioned quantities. The syntax is:

```haskell
with Dimension =>
  ([Symbol =>] SYMBOL, DIMENSION_VALUE {, DIMENSION_Value})
```

```plaintext
SYMBOL ::= STRING_LITERAL | CHARACTER_LITERAL
DIMENSION_VALUE ::= RATIONAL
| others => RATIONAL
| DISCRETE_CHOICE_LIST => RATIONAL
```

```plaintext
RATIONAL ::= [-] NUMERIC_LITERAL /[ NUMERIC_LITERAL]
```

This aspect can only be applied to a subtype whose parent type has a Dimension_System aspect. The aspect must specify values for all dimensions of the system. The rational values are the powers of the corresponding dimensions that are used by the compiler to verify that physical (numeric) computations are dimensionally consistent. For example, the computation of a force must result in dimensions (L => 1, M => 1, T => -2). For further examples of the usage of this aspect, see package System.Dim.Mks. Note that when the dimensioned type is an integer type, then any dimension value must be an integer literal.
3.10 Aspect Dimension_System

The `Dimension_System` aspect is used to define a system of dimensions that will be used in subsequent subtype declarations with `Dimension` aspects that reference this system. The syntax is:

```ada
with Dimension_System => (DIMENSION {, DIMENSION});

DIMENSION ::= ([Unit_Name =>] IDENTIFIER,
               [Unit_Symbol =>] SYMBOL,
               [Dim_Symbol =>] SYMBOL)

SYMBOL ::= CHARACTER_LITERAL | STRING_LITERAL
```

This aspect is applied to a type, which must be a numeric derived type (typically a floating-point type), that will represent values within the dimension system. Each `DIMENSION` corresponds to one particular dimension. A maximum of 7 dimensions may be specified. `Unit_Name` is the name of the dimension (for example `Meter`). `Unit_Symbol` is the shorthand used for quantities of this dimension (for example `m` for `Meter`). `Dim_Symbol` gives the identification within the dimension system (typically this is a single letter, e.g. `L` standing for length for unit name `Meter`). The `Unit_Symbol` is used in formatted output of dimensioned quantities. The `Dim_Symbol` is used in error messages when numeric operations have inconsistent dimensions.

GNAT provides the standard definition of the International MKS system in the run-time package `System.Dim.Mks`. You can easily define similar packages for cgs units or British units, and define conversion factors between values in different systems. The MKS system is characterized by the following aspect:

```ada
type Mks_Type is new Long_Long_Float with
   Dimension_System => (
      (Unit_Name => Meter, Unit_Symbol => 'm', Dim_Symbol => 'L'),
      (Unit_Name => Kilogram, Unit_Symbol => "kg", Dim_Symbol => 'M'),
      (Unit_Name => Second, Unit_Symbol => 's', Dim_Symbol => 'T'),
      (Unit_Name => Ampere, Unit_Symbol => 'A', Dim_Symbol => 'I'),
      (Unit_Name => Kelvin, Unit_Symbol => 'K', Dim_Symbol => '@'),
      (Unit_Name => Mole, Unit_Symbol => "mol", Dim_Symbol => 'N'),
      (Unit_Name => Candela, Unit_Symbol => "cd", Dim_Symbol => 'J'));
```

Note that in the above type definition, we use the at symbol (@) to represent a theta character (avoiding the use of extended Latin-1 characters in this context).

See section ‘Performing Dimensionality Analysis in GNAT’ in the GNAT Users Guide for detailed examples of use of the dimension system.

3.11 Aspect Disable_Controlled

The aspect `Disable_Controlled` is defined for controlled record types. If active, this aspect causes suppression of all related calls to `Initialize`, `Adjust`, and `Finalize`. The intended use is for conditional compilation, where for example you might want a record to be controlled or not depending on whether some run-time check is enabled or suppressed.
3.12 Aspect Effective_Reads
This aspect is equivalent to [pragma Effective_Reads], page 26.

3.13 Aspect Effective_Writes
This aspect is equivalent to [pragma Effective_Writes], page 27.

3.14 Aspect Extensions_Visible
This aspect is equivalent to [pragma Extensions_Visible], page 37.

3.15 Aspect Favor_Top_Level
This boolean aspect is equivalent to [pragma Favor_Top_Level], page 39.

3.16 Aspect Ghost
This aspect is equivalent to [pragma Ghost], page 40.

3.17 Aspect Global
This aspect is equivalent to [pragma Global], page 40.

3.18 Aspect Initial_Condition
This aspect is equivalent to [pragma Initial_Condition], page 46.

3.19 Aspect Initializes
This aspect is equivalent to [pragma Initializes], page 48.

3.20 Aspect Inline_Always
This boolean aspect is equivalent to [pragma Inline_Always], page 48.

3.21 Aspect Invariant
This aspect is equivalent to [pragma Invariant], page 51. It is a synonym for the language defined aspect Type_Invariant except that it is separately controllable using pragma Assertion_Policy.

3.22 Aspect Invariant’Class
This aspect is equivalent to [pragma Type_Invariant_Class], page 96. It is a synonym for the language defined aspect Type_Invariant’Class except that it is separately controllable using pragma Assertion_Policy.
3.23 Aspect Iterable

This aspect provides a light-weight mechanism for loops and quantified expressions over container types, without the overhead imposed by the tampering checks of standard Ada 2012 iterators. The value of the aspect is an aggregate with six named components, of which the last three are optional: First, Next, Has_Element, Element, Last, and Previous. When only the first three components are specified, only the for .. in form of iteration over cursors is available. When Element is specified, both this form and the for .. of form of iteration over elements are available. If the last two components are specified, reverse iterations over the container can be specified (analogous to what can be done over predefined containers that support the Reverse_Iterator interface). The following is a typical example of use:

```ada
type List is private with
  Iterable => (First => First_Cursor,
               Next  => Advance,
               Has_Element => Cursor_Has_Element
               [,Element => Get_Element]
               [,Last  => Last_Cursor]
               [,Previous => Retreat]);
```

* The values of First and Last are primitive operations of the container type that return a Cursor, which must be a type declared in the container package or visible from it. For example:
  ```ada
  function First_Cursor (Cont : Container) return Cursor;
  function Last_Cursor (Cont : Container) return Cursor;
  ```

* The values of Next and Previous are primitive operations of the container type that take both a container and a cursor and yield a cursor. For example:
  ```ada
  function Advance (Cont : Container; Position : Cursor) return Cursor;
  function Retreat (Cont : Container; Position : Cursor) return Cursor;
  ```

* The value of Has_Element is a primitive operation of the container type that takes both a container and a cursor and yields a boolean. For example:
  ```ada
  function Cursor_Has_Element (Cont : Container; Position : Cursor) return Boolean;
  ```

* The value of Element is a primitive operation of the container type that takes both a container and a cursor and yields an Element_Type, which must be a type declared in the container package or visible from it. For example:
  ```ada
  function Get_Element (Cont : Container; Position : Cursor) return Element_Type;
  ```

This aspect is used in the GNAT-defined formal container packages.

3.24 Aspect Linker_Section

This aspect is equivalent to [pragma Linker_Section], page 54.

3.25 Aspect Lock_Free

This boolean aspect is equivalent to [pragma Lock_Free], page 55.
3.26 Aspect Max_Queue_Length
This aspect is equivalent to [pragma Max_Queue_Length], page 58.

3.27 Aspect No_Caching
This boolean aspect is equivalent to [pragma No_Caching], page 59.

3.28 Aspect No_Elaboration_Code_All
This aspect is equivalent to [pragma No_Elaboration_Code_All], page 59 for a program unit.

3.29 Aspect No_Inline
This boolean aspect is equivalent to [pragma No_Inline], page 60.

3.30 Aspect No_Tagged_Streams
This aspect is equivalent to [pragma No_Tagged_Streams], page 61 with an argument specifying a root tagged type (thus this aspect can only be applied to such a type).

3.31 Aspect No_Task_Parts
Applies to a type. If True, requires that the type and any descendants do not have any task parts. The rules for this aspect are the same as for the language-defined No_Controlled_Parts aspect (see RM-H.4.1), replacing “controlled” with “task”.
If No_Task_Parts is True for a type T, then the compiler can optimize away certain tasking-related code that would otherwise be needed for T'Class, because descendants of T might contain tasks.

3.32 Aspect Object_Size
This aspect is equivalent to [attribute Object_Size], page 124.

3.33 Aspect Obsolescent
This aspect is equivalent to [pragma Obsolescent], page 62. Note that the evaluation of this aspect happens at the point of occurrence, it is not delayed until the freeze point.

3.34 Aspect Part_Of
This aspect is equivalent to [pragma Part_Of], page 67.

3.35 Aspect Persistent_BSS
This boolean aspect is equivalent to [pragma Persistent_BSS], page 68.

3.36 Aspect Predicate
This aspect is equivalent to [pragma Predicate], page 72. It is thus similar to the language defined aspects Dynamic_Predicate and Static_Predicate except that whether the resulting predicate is static or dynamic is controlled by the form of the expression. It is also separately controllable using pragma Assertion_Policy.
3.37 Aspect Pure Function
This boolean aspect is equivalent to [pragma Pure Function], page 78.

3.38 Aspect Refined Depends
This aspect is equivalent to [pragma Refined Depends], page 79.

3.39 Aspect Refined Global
This aspect is equivalent to [pragma Refined Global], page 80.

3.40 Aspect Refined Post
This aspect is equivalent to [pragma Refined Post], page 80.

3.41 Aspect Refined State
This aspect is equivalent to [pragma Refined State], page 80.

3.42 Aspect Relaxed Initialization
For the syntax and semantics of this aspect, see the SPARK 2014 Reference Manual, section 6.10.

3.43 Aspect Remote Access Type
This aspect is equivalent to [pragma Remote Access Type], page 81.

3.44 Aspect Secondary Stack Size
This aspect is equivalent to [pragma Secondary Stack Size], page 83.

3.45 Aspect Scalar Storage Order
This aspect is equivalent to a [attribute Scalar Storage Order], page 127.

3.46 Aspect Shared
This boolean aspect is equivalent to [pragma Shared], page 84 and is thus a synonym for aspect Atomic.

3.47 Aspect Simple Storage Pool
This aspect is equivalent to [attribute Simple Storage Pool], page 130.

3.48 Aspect Simple Storage Pool Type
This boolean aspect is equivalent to [pragma Simple Storage Pool Type], page 84.

3.49 Aspect SPARK Mode
This aspect is equivalent to [pragma SPARK Mode], page 87 and may be specified for either or both of the specification and body of a subprogram or package.
3.50 Aspect Suppress_DEBUG_INFO
This boolean aspect is equivalent to [pragma Suppress_DEBUG_INFO], page 92.

3.51 Aspect Suppress_INITIALIZATION
This boolean aspect is equivalent to [pragma Suppress_INITIALIZATION], page 92.

3.52 Aspect TEST_CASE
This aspect is equivalent to [pragma TEST_CASE], page 94.

3.53 Aspect THREAD_LOCAL_STORAGE
This boolean aspect is equivalent to [pragma THREAD_LOCAL_STORAGE], page 95.

3.54 Aspect UNIVERSAL_ALIASING
This boolean aspect is equivalent to [pragma UNIVERSAL_ALIASING], page 97.

3.55 Aspect UNMODIFIED
This boolean aspect is equivalent to [pragma UNMODIFIED], page 97.

3.56 Aspect UNREFERENCED
This boolean aspect is equivalent to [pragma UNREFERENCED], page 98.
When using the -gnat2022 switch, this aspect is also supported on formal parameters, which is in particular the only form possible for expression functions.

3.57 Aspect UNREFERENCED_OBJECTS
This boolean aspect is equivalent to [pragma UNREFERENCED_OBJECTS], page 99.

3.58 Aspect VALUE_SIZE
This aspect is equivalent to [attribute VALUE_SIZE], page 138.

3.59 Aspect VOLATILE_FULL_ACCESS
This boolean aspect is equivalent to [pragma VOLATILE_FULL_ACCESS], page 101.

3.60 Aspect VOLATILE_FUNCTION
This boolean aspect is equivalent to [pragma VOLATILE_FUNCTION], page 102.

3.61 Aspect WARNINGS
This aspect is equivalent to the two argument form of [pragma WARNINGS], page 103, where the first argument is ON or OFF and the second argument is the entity.
4 Implementation Defined Attributes

Ada defines (throughout the Ada reference manual, summarized in Annex K), a set of attributes that provide useful additional functionality in all areas of the language. These language defined attributes are implemented in GNAT and work as described in the Ada Reference Manual.

In addition, Ada allows implementations to define additional attributes whose meaning is defined by the implementation. GNAT provides a number of these implementation-dependent attributes which can be used to extend and enhance the functionality of the compiler. This section of the GNAT reference manual describes these additional attributes. It also describes additional implementation-dependent features of standard language-defined attributes.

Note that any program using these attributes may not be portable to other compilers (although GNAT implements this set of attributes on all platforms). Therefore if portability to other compilers is an important consideration, you should minimize the use of these attributes.

4.1 Attribute Abort_Signal

Standard’Abort_Signal (Standard is the only allowed prefix) provides the entity for the special exception used to signal task abort or asynchronous transfer of control. Normally this attribute should only be used in the tasking runtime (it is highly peculiar, and completely outside the normal semantics of Ada, for a user program to intercept the abort exception).

4.2 Attribute Address_Size

Standard’Address_Size (Standard is the only allowed prefix) is a static constant giving the number of bits in an Address. It is the same value as System.Address’Size, but has the advantage of being static, while a direct reference to System.Address’Size is nonstatic because Address is a private type.

4.3 Attribute Asm_Input

The Asm_Input attribute denotes a function that takes two parameters. The first is a string, the second is an expression of the type designated by the prefix. The first (string) argument is required to be a static expression, and is the constraint for the parameter, (e.g., what kind of register is required). The second argument is the value to be used as the input argument. The possible values for the constant are the same as those used in the RTL, and are dependent on the configuration file used to built the GCC back end. [Machine Code Insertions], page 277

4.4 Attribute Asm_Output

The Asm_Output attribute denotes a function that takes two parameters. The first is a string, the second is the name of a variable of the type designated by the attribute prefix. The first (string) argument is required to be a static expression and designates the constraint for the parameter (e.g., what kind of register is required). The second argument is the variable to be updated with the result. The possible values for constraint are the same as
those used in the RTL, and are dependent on the configuration file used to build the GCC back end. If there are no output operands, then this argument may either be omitted, or explicitly given as `No_Output_Operands`. [Machine Code Insertions], page 277

### 4.5 Attribute `Atomic_Always_Lock_Free`

The prefix of the `Atomic_Always_Lock_Free` attribute is a type. The result is a Boolean value which is True if the type has discriminants, and False otherwise. The result indicates whether atomic operations are supported by the target for the given type.

### 4.6 Attribute `Bit`

`obj`'Bit`, where `obj` is any object, yields the bit offset within the storage unit (byte) that contains the first bit of storage allocated for the object. The value of this attribute is of the type `universal_integer` and is always a nonnegative number smaller than `System.Storage_Unit`.

For an object that is a variable or a constant allocated in a register, the value is zero. (The use of this attribute does not force the allocation of a variable to memory).

For an object that is a formal parameter, this attribute applies to either the matching actual parameter or to a copy of the matching actual parameter.

For an access object the value is zero. Note that `obj.all`'Bit` is subject to an `Access_Check` for the designated object. Similarly for a record component `X.C`'Bit` is subject to a discriminant check and `X(I).Bit` and `X(I1..I2)`'Bit` are subject to index checks.

This attribute is designed to be compatible with the DEC Ada 83 definition and implementation of the `Bit` attribute.

### 4.7 Attribute `Bit_Position`

`R.C`'Bit_Position`, where `R` is a record object and `C` is one of the fields of the record type, yields the bit offset within the record contains the first bit of storage allocated for the object. The value of this attribute is of the type `universal_integer`. The value depends only on the field `C` and is independent of the alignment of the containing record `R`.

### 4.8 Attribute `Code_Address`

The `'Address` attribute may be applied to subprograms in Ada 95 and Ada 2005, but the intended effect seems to be to provide an address value which can be used to call the subprogram by means of an address clause as in the following example:

```ada
procedure K is ...

procedure L;
   for L'Address use K'Address;
   pragma Import (Ada, L);
```

A call to `L` is then expected to result in a call to `K`. In Ada 83, where there were no access-to-subprogram values, this was a common work-around for getting the effect of an indirect call. GNAT implements the above use of `Address` and the technique illustrated by the example code works correctly.
However, for some purposes, it is useful to have the address of the start of the generated code for the subprogram. On some architectures, this is not necessarily the same as the \texttt{Address} value described above. For example, the \texttt{Address} value may reference a subprogram descriptor rather than the subprogram itself.

The \('\texttt{Code\_Address}'\) attribute, which can only be applied to subprogram entities, always returns the address of the start of the generated code of the specified subprogram, which may or may not be the same value as is returned by the corresponding \('\texttt{Address}'\) attribute.

### 4.9 Attribute Compiler\_Version

\texttt{Standard'Compiler\_Version} (\texttt{Standard} is the only allowed prefix) yields a static string identifying the version of the compiler being used to compile the unit containing the attribute reference.

### 4.10 Attribute Constrained

In addition to the usage of this attribute in the Ada RM, GNAT also permits the use of the \('\texttt{Constrained}'\) attribute in a generic template for any type, including types without discriminants. The value of this attribute in the generic instance when applied to a scalar type or a record type without discriminants is always \texttt{True}. This usage is compatible with older Ada compilers, including notably DEC Ada.

### 4.11 Attribute Default\_Bit\_Order

\texttt{Standard'Default\_Bit\_Order} (\texttt{Standard} is the only allowed prefix), provides the value \texttt{System.Default\_Bit\_Order} as a \texttt{Pos} value (0 for \texttt{High\_Order\_First}, 1 for \texttt{Low\_Order\_First}). This is used to construct the definition of \texttt{Default\_Bit\_Order} in package \texttt{System}.

### 4.12 Attribute Default\_Scalar\_Storage\_Order

\texttt{Standard'Default\_Scalar\_Storage\_Order} (\texttt{Standard} is the only allowed prefix), provides the current value of the default scalar storage order (as specified using pragma \texttt{Default\_Scalar\_Storage\_Order}, or equal to \texttt{Default\_Bit\_Order} if unspecified) as a \texttt{System.Bit\_Order} value. This is a static attribute.

### 4.13 Attribute Deref

The attribute \texttt{typ'Deref(expr)} where \texttt{expr} is of type \texttt{System.Address} yields the variable of type \texttt{typ} that is located at the given address. It is similar to \texttt{(totyp(expr).all)}, where \texttt{totyp} is an unchecked conversion from address to a named access-to-\texttt{typ} type, except that it yields a variable, so it can be used on the left side of an assignment.

### 4.14 Attribute Descriptor\_Size

Nonstatic attribute \texttt{Descriptor\_Size} returns the size in bits of the descriptor allocated for a type. The result is non-zero only for unconstrained array types and the returned value is of type universal integer. In GNAT, an array descriptor contains bounds information and is located immediately before the first element of the array.

```ada
type Unconstr_Array is array (Short_Short_Integer range <> ) of Positive;
```
Put_Line ("Descriptor size = " & Unconstr_Array'Descriptor_Size'Img);

The attribute takes into account any padding due to the alignment of the component type. In the example above, the descriptor contains two values of type Short_Short_Integer representing the low and high bound. But, since Positive has an alignment of 4, the size of the descriptor is \( 2 \times \text{Short}_\text{Short}_\text{Integer}'\text{Size} \) rounded up to the next multiple of 32, which yields a size of 32 bits, i.e. including 16 bits of padding.

4.15 Attribute Elaborated

The prefix of the 'Elaborated attribute must be a unit name. The value is a Boolean which indicates whether or not the given unit has been elaborated. This attribute is primarily intended for internal use by the generated code for dynamic elaboration checking, but it can also be used in user programs. The value will always be True once elaboration of all units has been completed. An exception is for units which need no elaboration, the value is always False for such units.

4.16 Attribute Elab_Body

This attribute can only be applied to a program unit name. It returns the entity for the corresponding elaboration procedure for elaborating the body of the referenced unit. This is used in the main generated elaboration procedure by the binder and is not normally used in any other context. However, there may be specialized situations in which it is useful to be able to call this elaboration procedure from Ada code, e.g., if it is necessary to do selective re-elaboration to fix some error.

4.17 Attribute Elab_Spec

This attribute can only be applied to a program unit name. It returns the entity for the corresponding elaboration procedure for elaborating the spec of the referenced unit. This is used in the main generated elaboration procedure by the binder and is not normally used in any other context. However, there may be specialized situations in which it is useful to be able to call this elaboration procedure from Ada code, e.g., if it is necessary to do selective re-elaboration to fix some error.

4.18 Attribute Elab_Subp_Body

This attribute can only be applied to a library level subprogram name and is only allowed in CodePeer mode. It returns the entity for the corresponding elaboration procedure for elaborating the body of the referenced subprogram unit. This is used in the main generated elaboration procedure by the binder in CodePeer mode only and is unrecognized otherwise.

4.19 Attribute Emax

The Emax attribute is provided for compatibility with Ada 83. See the Ada 83 reference manual for an exact description of the semantics of this attribute.

4.20 Attribute Enabled

The Enabled attribute allows an application program to check at compile time to see if the designated check is currently enabled. The prefix is a simple identifier, referencing any
predefined check name (other than All\_Checks) or a check name introduced by pragma Check\_Name. If no argument is given for the attribute, the check is for the general state of the check, if an argument is given, then it is an entity name, and the check indicates whether an Suppress or Unsuppress has been given naming the entity (if not, then the argument is ignored).

Note that instantiations inherit the check status at the point of the instantiation, so a useful idiom is to have a library package that introduces a check name with pragma Check\_Name, and then contains generic packages or subprograms which use the Enabled attribute to see if the check is enabled. A user of this package can then issue a pragma Suppress or pragma Unsuppress before instantiating the package or subprogram, controlling whether the check will be present.

4.21 Attribute Enum\_Rep

Note that this attribute is now standard in Ada 202x and is available as an implementation defined attribute for earlier Ada versions.

For every enumeration subtype $S$, $S$'Enum\_Rep denotes a function with the following spec:

$$\text{function } S'\text{Enum\_Rep (Arg : } S\text{'}Base) \text{ return } <\text{Universal}\_\text{Integer}>;$$

It is also allowable to apply Enum\_Rep directly to an object of an enumeration type or to a non-overloaded enumeration literal. In this case $S'\text{Enum\_Rep}$ is equivalent to typ'Enum\_Rep($S$) where typ is the type of the enumeration literal or object.

The function returns the representation value for the given enumeration value. This will be equal to value of the Pos attribute in the absence of an enumeration representation clause. This is a static attribute (i.e., the result is static if the argument is static).

$S'\text{Enum\_Rep}$ can also be used with integer types and objects, in which case it simply returns the integer value. The reason for this is to allow it to be used for (<> discrete formal arguments in a generic unit that can be instantiated with either enumeration types or integer types. Note that if Enum\_Rep is used on a modular type whose upper bound exceeds the upper bound of the largest signed integer type, and the argument is a variable, so that the universal integer calculation is done at run time, then the call to Enum\_Rep may raise Constraint\_Error.

4.22 Attribute Enum\_Val

Note that this attribute is now standard in Ada 202x and is available as an implementation defined attribute for earlier Ada versions.

For every enumeration subtype $S$, $S$'Enum\_Val denotes a function with the following spec:

$$\text{function } S'\text{Enum\_Val (Arg : } <\text{Universal}\_\text{Integer}>) \text{ return } S'\text{Base};$$

The function returns the enumeration value whose representation matches the argument, or raises Constraint\_Error if no enumeration literal of the type has the matching value. This will be equal to value of the Val attribute in the absence of an enumeration representation clause. This is a static attribute (i.e., the result is static if the argument is static).

4.23 Attribute Epsilon

The Epsilon attribute is provided for compatibility with Ada 83. See the Ada 83 reference manual for an exact description of the semantics of this attribute.
4.24 Attribute Fast.Math

\texttt{Standard’Fast.Math} (\texttt{Standard} is the only allowed prefix) yields a static Boolean value that is True if pragma \texttt{Fast.Math} is active, and False otherwise.

4.25 Attribute Finalization_Size

The prefix of attribute \texttt{Finalization_Size} must be an object or a non-class-wide type. This attribute returns the size of any hidden data reserved by the compiler to handle finalization-related actions. The type of the attribute is ‘universal_integer’.

\texttt{Finalization_Size} yields a value of zero for a type with no controlled parts, an object whose type has no controlled parts, or an object of a class-wide type whose tag denotes a type with no controlled parts.

Note that only heap-allocated objects contain finalization data.

4.26 Attribute Fixed_Value

For every fixed-point type \( S \), \( S’Fixed_Value \) denotes a function with the following specification:

\begin{verbatim}
function S’Fixed_Value (Arg : <Universal_Integer>) return S;
\end{verbatim}

The value returned is the fixed-point value \( V \) such that:

\[ V = \text{Arg} \ast S’\text{Small} \]

The effect is thus similar to first converting the argument to the integer type used to represent \( S \), and then doing an unchecked conversion to the fixed-point type. The difference is that there are full range checks, to ensure that the result is in range. This attribute is primarily intended for use in implementation of the input-output functions for fixed-point values.

4.27 Attribute From.Any

This internal attribute is used for the generation of remote subprogram stubs in the context of the Distributed Systems Annex.

4.28 Attribute Has_Access_Values

The prefix of the \texttt{Has_Access_Values} attribute is a type. The result is a Boolean value which is True if the is an access type, or is a composite type with a component (at any nesting depth) that is an access type, and is False otherwise. The intended use of this attribute is in conjunction with generic definitions. If the attribute is applied to a generic private type, it indicates whether or not the corresponding actual type has access values.

4.29 Attribute Has_Discriminants

The prefix of the \texttt{Has_Discriminants} attribute is a type. The result is a Boolean value which is True if the type has discriminants, and False otherwise. The intended use of this attribute is in conjunction with generic definitions. If the attribute is applied to a generic private type, it indicates whether or not the corresponding actual type has discriminants.
4.30 Attribute Has_Tagged_Values

The prefix of the `Has_Tagged_Values` attribute is a type. The result is a Boolean value which is True if the type is a composite type (array or record) that is either a tagged type or has a subcomponent that is tagged, and is False otherwise. The intended use of this attribute is in conjunction with generic definitions. If the attribute is applied to a generic private type, it indicates whether or not the corresponding actual type has access values.

4.31 Attribute Img

The `Img` attribute differs from `Image` in that, while both can be applied directly to an object, `Img` cannot be applied to types.

Example usage of the attribute:

```pascal
Put_Line ("X = " & X'Img);
```

which has the same meaning as the more verbose:

```pascal
Put_Line ("X = " & T'Image (X));
```

where T is the (sub)type of the object X.

Note that technically, in analogy to `Image`, `X'Img` returns a parameterless function that returns the appropriate string when called. This means that `X'Img` can be renamed as a function-returning-string, or used in an instantiation as a function parameter.

4.32 Attribute Initialized

For the syntax and semantics of this attribute, see the SPARK 2014 Reference Manual, section 6.10.

4.33 Attribute Integer_Value

For every integer type S, S'Integer_Value denotes a function with the following spec:

```pascal
function S'Integer_Value (Arg : <Universal_Fixed>) return S;
```

The value returned is the integer value V, such that:

```pascal
Arg = V * T'Small
```

where T is the type of Arg. The effect is thus similar to first doing an unchecked conversion from the fixed-point type to its corresponding implementation type, and then converting the result to the target integer type. The difference is that there are full range checks, to ensure that the result is in range. This attribute is primarily intended for use in implementation of the standard input-output functions for fixed-point values.

4.34 Attribute Invalid_Value

For every scalar type S, S'Invalid_Value returns an undefined value of the type. If possible this value is an invalid representation for the type. The value returned is identical to the value used to initialize an otherwise uninitialized value of the type if pragma Initialize_Scalars is used, including the ability to modify the value with the binder -Sxx flag and relevant environment variables at run time.
4.35 Attribute Iterable
Equivalent to Aspect Iterable.

4.36 Attribute Large
The Large attribute is provided for compatibility with Ada 83. See the Ada 83 reference manual for an exact description of the semantics of this attribute.

4.37 Attribute Library_Level
P’Library_Level, where P is an entity name, returns a Boolean value which is True if the entity is declared at the library level, and False otherwise. Note that within a generic instantiation, the name of the generic unit denotes the instance, which means that this attribute can be used to test if a generic is instantiated at the library level, as shown in this example:

```ada
generic
...
package Gen is
  pragma Compile_Time_Error
  (not Gen’Library_Level,
    "Gen can only be instantiated at library level");
...
end Gen;
```

4.38 Attribute Loop_Entry
Syntax:

X’Loop_Entry [(loop_name)]

The Loop_Entry attribute is used to refer to the value that an expression had upon entry to a given loop in much the same way that the Old attribute in a subprogram postcondition can be used to refer to the value an expression had upon entry to the subprogram. The relevant loop is either identified by the given loop name, or it is the innermost enclosing loop when no loop name is given.

A Loop_Entry attribute can only occur within an Assert, Assert_And_Cut, Assume, Loop_Variant or Loop_Invariant pragma. In addition, such a pragma must be one of the items in the sequence of statements of a loop body, or nested inside block statements that appear in the sequence of statements of a loop body. A common use of Loop_Entry is to compare the current value of objects with their initial value at loop entry, in a Loop_Invariant pragma.

The effect of using X’Loop_Entry is the same as declaring a constant initialized with the initial value of X at loop entry. This copy is not performed if the loop is not entered, or if the corresponding pragmas are ignored or disabled.

4.39 Attribute Machine_Size
This attribute is identical to the Object_Size attribute. It is provided for compatibility with the DEC Ada 83 attribute of this name.
4.40 Attribute Mantissa

The Mantissa attribute is provided for compatibility with Ada 83. See the Ada 83 reference manual for an exact description of the semantics of this attribute.

4.41 Attribute Maximum_Alignment

Standard'Maximum_Alignment (Standard is the only allowed prefix) provides the maximum useful alignment value for the target. This is a static value that can be used to specify the alignment for an object, guaranteeing that it is properly aligned in all cases.

4.42 Attribute Max_Integer_Size

Standard'Max_Integer_Size (Standard is the only allowed prefix) provides the size of the largest supported integer type for the target. The result is a static constant.

4.43 Attribute Mechanism_Code

func'Mechanism_Code yields an integer code for the mechanism used for the result of function func, and subprog'Mechanism_Code (n) yields the mechanism used for formal parameter number ‘n’ (a static integer value, with 1 meaning the first parameter) of subprogram subprog. The code returned is:

‘1’ by copy (value)

‘2’ by reference

4.44 Attribute Null_Parameter

A reference T’Null_Parameter denotes an imaginary object of type or subtype T allocated at machine address zero. The attribute is allowed only as the default expression of a formal parameter, or as an actual expression of a subprogram call. In either case, the subprogram must be imported.

The identity of the object is represented by the address zero in the argument list, independent of the passing mechanism (explicit or default).

This capability is needed to specify that a zero address should be passed for a record or other composite object passed by reference. There is no way of indicating this without the Null_Parameter attribute.

4.45 Attribute Object_Size

The size of an object is not necessarily the same as the size of the type of an object. This is because by default object sizes are increased to be a multiple of the alignment of the object. For example, Natural'Size is 31, but by default objects of type Natural will have a size of 32 bits. Similarly, a record containing an integer and a character:

```ada
type Rec is record
    I : Integer;
    C : Character;
```
end record;

will have a size of 40 (that is \texttt{Rec’S\texttt{Size}} will be 40). The alignment will be 4, because of the integer field, and so the default size of record objects for this type will be 64 (8 bytes).

If the alignment of the above record is specified to be 1, then the object size will be 40 (5 bytes). This is true by default, and also an object size of 40 can be explicitly specified in this case.

A consequence of this capability is that different object sizes can be given to subtypes that would otherwise be considered in Ada to be statically matching. But it makes no sense to consider such subtypes as statically matching. Consequently, GNAT adds a rule to the static matching rules that requires object sizes to match. Consider this example:

1. procedure BadAVConvert is
2. type R is new Integer;
3. subtype R1 is R range 1 .. 10;
4. subtype R2 is R range 1 .. 10;
5. for R1’Object_Size use 8;
6. for R2’Object_Size use 16;
7. type R1P is access all R1;
8. type R2P is access all R2;
9. R1PV : R1P := new R1’(4);
10. R2PV : R2P;
11. begin
12. R2PV := R2P (R1PV);
13. end;

In the absence of lines 5 and 6, types \texttt{R1} and \texttt{R2} statically match and hence the conversion on line 12 is legal. But since lines 5 and 6 cause the object sizes to differ, GNAT considers that types \texttt{R1} and \texttt{R2} are not statically matching, and line 12 generates the diagnostic shown above.

Similar additional checks are performed in other contexts requiring statically matching subtypes.

4.46 Attribute Old

In addition to the usage of \texttt{Old} defined in the Ada 2012 RM (usage within \texttt{Post} aspect), GNAT also permits the use of this attribute in implementation defined pragmas \texttt{Postcondition, Contract_Cases} and \texttt{Test_Case}. Also usages of \texttt{Old} which would be illegal according to the Ada 2012 RM definition are allowed under control of implementation defined pragma \texttt{Unevaluated_Use_Of_Old}.

4.47 Attribute Passed_By_Reference

\texttt{typ’Passed_By_Reference} for any subtype \texttt{typ} returns a value of type \texttt{Boolean} value that is \texttt{True} if the type is normally passed by reference and \texttt{False} if the type is normally passed
by copy in calls. For scalar types, the result is always False and is static. For non-scalar types, the result is nonstatic.

### 4.48 Attribute Pool_Address

$X'\text{Pool\_Address}$ for any object $X$ returns the address of $X$ within its storage pool. This is the same as $X'\text{Address}$, except that for an unconstrained array whose bounds are allocated just before the first component, $X'\text{Pool\_Address}$ returns the address of those bounds, whereas $X'\text{Address}$ returns the address of the first component.

Here, we are interpreting ‘storage pool’ broadly to mean wherever the object is allocated, which could be a user-defined storage pool, the global heap, on the stack, or in a static memory area. For an object created by new, $\text{Ptr.all}'\text{Pool\_Address}$ is what is passed to Allocate and returned from Deallocate.

### 4.49 Attribute Range_Length

$\text{typ}'\text{Range\_Length}$ for any discrete type $\text{typ}$ yields the number of values represented by the subtype (zero for a null range). The result is static for static subtypes. $\text{Range\_Length}$ applied to the index subtype of a one dimensional array always gives the same result as $\text{Length}$ applied to the array itself.

### 4.50 Attribute Restriction_Set

This attribute allows compile time testing of restrictions that are currently in effect. It is primarily intended for specializing code in the run-time based on restrictions that are active (e.g. don’t need to save fpt registers if restriction No_Floating_Point is known to be in effect), but can be used anywhere.

There are two forms:

- $\text{System}'\text{Restriction\_Set} (\text{partition\_boolean\_restriction\_NAME})$
- $\text{System}'\text{Restriction\_Set} (\text{No\_Dependence} => \text{library\_unit\_NAME})$

In the case of the first form, the only restriction names allowed are parameterless restrictions that are checked for consistency at bind time. For a complete list see the subtype $\text{System.Rident.Partition\_Boolean\_Restrictions}$.

The result returned is True if the restriction is known to be in effect, and False if the restriction is known not to be in effect. An important guarantee is that the value of a Restriction_Set attribute is known to be consistent throughout all the code of a partition. This is trivially achieved if the entire partition is compiled with a consistent set of restriction pragmas. However, the compilation model does not require this. It is possible to compile one set of units with one set of pragmas, and another set of units with another set of pragmas. It is even possible to compile a spec with one set of pragmas, and then WITH the same spec with a different set of pragmas. Inconsistencies in the actual use of the restriction are checked at bind time.

In order to achieve the guarantee of consistency for the Restriction_Set pragma, we consider that a use of the pragma that yields False is equivalent to a violation of the restriction.

So for example if you write
if System'Restriction_Set (No_Floating_Point) then
  ...
else
  ...
end if;

And the result is False, so that the else branch is executed, you can assume that this restriction is not set for any unit in the partition. This is checked by considering this use of the restriction pragma to be a violation of the restriction No_Floating_Point. This means that no other unit can attempt to set this restriction (if some unit does attempt to set it, the binder will refuse to bind the partition).

Technical note: The restriction name and the unit name are interpreted entirely syntactically, as in the corresponding Restrictions pragma, they are not analyzed semantically, so they do not have a type.

### 4.51 Attribute Result

`function'Result` can only be used with in a Postcondition pragma for a function. The prefix must be the name of the corresponding function. This is used to refer to the result of the function in the postcondition expression. For a further discussion of the use of this attribute and examples of its use, see the description of pragma Postcondition.

### 4.52 Attribute Safe_Emax

The `Safe_Emax` attribute is provided for compatibility with Ada 83. See the Ada 83 reference manual for an exact description of the semantics of this attribute.

### 4.53 Attribute Safe_Large

The `Safe_Large` attribute is provided for compatibility with Ada 83. See the Ada 83 reference manual for an exact description of the semantics of this attribute.

### 4.54 Attribute Safe_Small

The `Safe_Small` attribute is provided for compatibility with Ada 83. See the Ada 83 reference manual for an exact description of the semantics of this attribute.

### 4.55 Attribute Scalar_Storage_Order

For every array or record type `S`, the representation attribute `Scalar_Storage_Order` denotes the order in which storage elements that make up scalar components are ordered within `S`. The value given must be a static expression of type `System.Bit_Order`. The following is an example of the use of this feature:

```ada
-- Component type definitions
subtype Yr_Type is Natural range 0 .. 127;
subtype Mo_Type is Natural range 1 .. 12;
subtype Da_Type is Natural range 1 .. 31;
```
-- Record declaration

type Date is record
   Years_Since_1980 : Yr_Type;
   Month            : Mo_Type;
   Day_Of_Month    : Da_Type;
end record;

-- Record representation clause

for Date use record
   Years_Since_1980 at 0 range 0 .. 6;
   Month            at 0 range 7 .. 10;
   Day_Of_Month    at 0 range 11 .. 15;
end record;

-- Attribute definition clauses

for Date'Bit_Order use System.High_Order_First;
for Date'Scalar_Storage_Order use System.High_Order_First;
-- If Scalar_Storage_Order is specified, it must be consistent with
-- Bit_Order, so it's best to always define the latter explicitly if
-- the former is used.

Other properties are as for the standard representation attribute Bit_Order defined by Ada
RM 13.5.3(4). The default is System.Default_Bit_Order.

For a record type T, if T'Scalar_Storage_Order is specified explicitly, it shall be equal
to T'Bit_Order. Note: this means that if a Scalar_Storage_Order attribute definition
clause is not confirming, then the type’s Bit_Order shall be specified explicitly and set to
the same value.

Derived types inherit an explicitly set scalar storage order from their parent types. This
may be overridden for the derived type by giving an explicit scalar storage order for it.
However, for a record extension, the derived type must have the same scalar storage order
as the parent type.

A component of a record type that is itself a record or an array and that does not start and
end on a byte boundary must have have the same scalar storage order as the record type.
A component of a bit-packed array type that is itself a record or an array must have the
same scalar storage order as the array type.

No component of a type that has an explicit Scalar_Storage_Order attribute definition
may be aliased.

A confirming Scalar_Storage_Order attribute definition clause (i.e. with a value equal to
System.Default_Bit_Order) has no effect.

If the opposite storage order is specified, then whenever the value of a scalar component of
an object of type S is read, the storage elements of the enclosing machine scalar are first
reversed (before retrieving the component value, possibly applying some shift and mask
operatings on the enclosing machine scalar), and the opposite operation is done for writes.
In that case, the restrictions set forth in 13.5.1(10.3/2) for scalar components are relaxed. Instead, the following rules apply:

* the underlying storage elements are those at positions \( (\text{position} + \frac{\text{first\_bit}}{\text{storage\_element\_size}}) \ldots (\text{position} + \frac{(\text{last\_bit} + \text{storage\_element\_size} - 1)}{\text{storage\_element\_size}}) \)

* the sequence of underlying storage elements shall have a size no greater than the largest machine scalar

* the enclosing machine scalar is defined as the smallest machine scalar starting at a position no greater than \( \text{position} + \frac{\text{first\_bit}}{\text{storage\_element\_size}} \) and covering storage elements at least up to \( \text{position} + \frac{(\text{last\_bit} + \text{storage\_element\_size} - 1)}{\text{storage\_element\_size}} \)

* the position of the component is interpreted relative to that machine scalar.

If no scalar storage order is specified for a type (either directly, or by inheritance in the case of a derived type), then the default is normally the native ordering of the target, but this default can be overridden using pragma `Default_Scalar_Storage_Order`.

If a component of `T` is itself of a record or array type, the specified `Scalar_Storage_Order` does `not` apply to that nested type: an explicit attribute definition clause must be provided for the component type as well if desired.

Representation changes that explicitly or implicitly toggle the scalar storage order are not supported and may result in erroneous execution of the program, except when performed by means of an instance of `Ada.Unchecked_Conversion`.

In particular, overlays are not supported and a warning is given for them:

```ada
type Rec_LE is record
  I : Integer;
end record;

for Rec_LE use record
  I at 0 range 0 .. 31;
end record;

for Rec_LE’Bit_Order use System.Low_Order_First;
for Rec_LE’Scalar_Storage_Order use System.Low_Order_First;

type Rec_BE is record
  I : Integer;
end record;

for Rec_BE use record
  I at 0 range 0 .. 31;
end record;

for Rec_BE’Bit_Order use System.High_Order_First;
for Rec_BE’Scalar_Storage_Order use System.High_Order_First;

R_LE : Rec_LE;
```
R_BE : Rec_BE;
for R_BE'Address use R_LE'Address;

warning: overlay changes scalar storage order [enabled by default]
In most cases, such representation changes ought to be replaced by an instantiation of a
function or procedure provided by GNAT.Byte_Swapping.
Note that the scalar storage order only affects the in-memory data representation. It has
no effect on the representation used by stream attributes.
Note that debuggers may be unable to display the correct value of scalar components of a
type for which the opposite storage order is specified.

4.56 Attribute Simple_Storage_Pool

For every nonformal, nonderived access-to-object type Acc, the representation attribute
Simple_Storage_Pool may be specified via an attribute_definition_clause (or by specifying
the equivalent aspect):

    My_Pool : My_Simple_Storage_Pool_Type;

type Acc is access My_Data_Type;

    for Acc'Simple_Storage_Pool use My_Pool;

The name given in an attribute_definition_clause for the Simple_Storage_Pool
attribute shall denote a variable of a 'simple storage pool type' (see pragma
Simple_Storage_Pool_Type).

The use of this attribute is only allowed for a prefix denoting a type for which it has been
specified. The type of the attribute is the type of the variable specified as the simple storage
pool of the access type, and the attribute denotes that variable.

It is illegal to specify both Storage_Pool and Simple_Storage_Pool for the same access
type.

If the Simple_Storage_Pool attribute has been specified for an access type, then applying
the Storage_Pool attribute to the type is flagged with a warning and its evaluation raises
the exception Program_Error.

If the Simple_Storage_Pool attribute has been specified for an access type S, then the
evaluation of the attribute S'Storage_Size returns the result of calling Storage_Size
(S'Simple_Storage_Pool), which is intended to indicate the number of storage elements
reserved for the simple storage pool. If the Storage_Size function has not been defined for
the simple storage pool type, then this attribute returns zero.

If an access type S has a specified simple storage pool of type SSP, then the evaluation of an
allocator for that access type calls the primitive Allocate procedure for type SSP, passing
S'Simple_Storage_Pool as the pool parameter. The detailed semantics of such allocators
is the same as those defined for allocators in section 13.11 of the Ada Reference Manual,
with the term 'simple storage pool' substituted for 'storage pool'.

If an access type S has a specified simple storage pool of type SSP, then a call to an instance
of the Ada.Unchecked_Deallocation for that access type invokes the primitive Deallocate
procedure for type SSP, passing \texttt{S’Simple\_Storage\_Pool} as the pool parameter. The detailed semantics of such unchecked deallocations is the same as defined in section 13.11.2 of the Ada Reference Manual, except that the term ‘simple storage pool’ is substituted for ‘storage pool’.

\section*{4.57 Attribute Small}

The \texttt{Small} attribute is defined in Ada 95 (and Ada 2005) only for fixed-point types. GNAT also allows this attribute to be applied to floating-point types for compatibility with Ada 83. See the Ada 83 reference manual for an exact description of the semantics of this attribute when applied to floating-point types.

\section*{4.58 Attribute Small\_Denominator}

\texttt{typ’S\_Small\_Denominator} for any fixed-point subtype \texttt{typ} yields the denominator in the representation of \texttt{typ’S\_Small} as a rational number with coprime factors (i.e. as an irreducible fraction).

\section*{4.59 Attribute Small\_Numerator}

\texttt{typ’S\_Small\_Numerator} for any fixed-point subtype \texttt{typ} yields the numerator in the representation of \texttt{typ’S\_Small} as a rational number with coprime factors (i.e. as an irreducible fraction).

\section*{4.60 Attribute Storage\_Unit}

\texttt{Standard’S\_Storage\_Unit} (\texttt{Standard} is the only allowed prefix) provides the same value as \texttt{System.Storage\_Unit}.

\section*{4.61 Attribute Stub\_Type}

The GNAT implementation of remote access-to-classwide types is organized as described in AARM section E.4 (20.t): a value of an RACW type (designating a remote object) is represented as a normal access value, pointing to a “stub” object which in turn contains the necessary information to contact the designated remote object. A call on any dispatching operation of such a stub object does the remote call, if necessary, using the information in the stub object to locate the target partition, etc.

For a prefix \texttt{T} that denotes a remote access-to-classwide type, \texttt{T’S\_Stub\_Type} denotes the type of the corresponding stub objects.

By construction, the layout of \texttt{T’S\_Stub\_Type} is identical to that of type \texttt{RACW\_Stub\_Type} declared in the internal implementation-defined unit \texttt{System\_Partition\_Interface}. Use of this attribute will create an implicit dependency on this unit.

\section*{4.62 Attribute System\_Allocator\_Alignment}

\texttt{Standard’S\_System\_Allocator\_Alignment} (\texttt{Standard} is the only allowed prefix) provides the observable guaranteed to be honored by the system allocator (\texttt{malloc}). This is a static value that can be used in user storage pools based on \texttt{malloc} either to reject allocation with alignment too large or to enable a realignment circuitry if the alignment request is larger than this value.
4.63 Attribute Target_Name

Standard’Target_Name (Standard is the only allowed prefix) provides a static string value that identifies the target for the current compilation. For GCC implementations, this is the standard gcc target name without the terminating slash (for example, GNAT 5.0 on windows yields “i586-pc-mingw32msv”).

4.64 Attribute To_Address

The System’To_Address (System is the only allowed prefix) denotes a function identical to System.Storage_Elements.To_Address except that it is a static attribute. This means that if its argument is a static expression, then the result of the attribute is a static expression. This means that such an expression can be used in contexts (e.g., preelaborable packages) which require a static expression and where the function call could not be used (since the function call is always nonstatic, even if its argument is static). The argument must be in the range \(-(2^{*}(m-1)) \ldots 2^{*}m-1\), where \(m\) is the memory size (typically 32 or 64). Negative values are interpreted in a modular manner (e.g., -1 means the same as \(16\#FFFF_FFFF\#\) on a 32 bits machine).

4.65 Attribute To.Any

This internal attribute is used for the generation of remote subprogram stubs in the context of the Distributed Systems Annex.

4.66 Attribute Type.Class

\texttt{typ’Type.Class} for any type or subtype \texttt{typ} yields the value of the type class for the full type of \texttt{typ}. If \texttt{typ} is a generic formal type, the value is the value for the corresponding actual subtype. The value of this attribute is of type System.Aux_DEC.Type.Class, which has the following definition:

\begin{verbatim}
   type Type.Class is
      (Type.Class_Enumeration,
       Type.Class_Integer,
       Type.Class_Fixed_Point,
       Type.Class_Floating_Point,
       Type.Class_Array,
       Type.Class_Record,
       Type.Class_Access,
       Type.Class_Task,
       Type.Class_Address);
\end{verbatim}

Protected types yield the value \texttt{Type.Class_Task}, which thus applies to all concurrent types. This attribute is designed to be compatible with the DEC Ada 83 attribute of the same name.

4.67 Attribute Type.Key

The \texttt{Type.Key} attribute is applicable to a type or subtype and yields a value of type Standard.String containing encoded information about the type or subtype. This provides improved compatibility with other implementations that support this attribute.
4.68 Attribute TypeCode

This internal attribute is used for the generation of remote subprogram stubs in the context of the Distributed Systems Annex.

4.69 Attribute Unconstrained_Array

The Unconstrained_Array attribute can be used with a prefix that denotes any type or subtype. It is a static attribute that yields True if the prefix designates an unconstrained array, and False otherwise. In a generic instance, the result is still static, and yields the result of applying this test to the generic actual.

4.70 Attribute Universal_Literal_String

The prefix of Universal_Literal_String must be a named number. The static result is the string consisting of the characters of the number as defined in the original source. This allows the user program to access the actual text of named numbers without intermediate conversions and without the need to enclose the strings in quotes (which would preclude their use as numbers).

For example, the following program prints the first 50 digits of pi:

```ada
with Text_IO; use Text_IO;
with Ada.Numerics;
procedure Pi is
begin
  Put (Ada.Numerics.Pi’Universal_Literal_String);
end;
```

4.71 Attribute Unrestricted_Access

The Unrestricted_Access attribute is similar to Access except that all accessibility and aliased view checks are omitted. This is a user-beware attribute.

For objects, it is similar to Address, for which it is a desirable replacement where the value desired is an access type. In other words, its effect is similar to first applying the Address attribute and then doing an unchecked conversion to a desired access type.

For subprograms, P’Unrestricted_Access may be used where P’Access would be illegal, to construct a value of a less-nested named access type that designates a more-nested subprogram. This value may be used in indirect calls, so long as the more-nested subprogram still exists; once the subprogram containing it has returned, such calls are erroneous. For example:

```ada
package body P is

  type Less_Nested is not null access procedure;
  Global : Less_Nested;

  procedure P1 is
  begin
    Global.all;
  end P1;
```
procedure P2 is
  Local_Var : Integer;

  procedure More_Nested is
  begin
    ... Local_Var ...
  end More_Nested;
begin
  Global := More_Nested'Unrestricted_Access;
  P1;
end P2;
end P;

When P1 is called from P2, the call via Global is OK, but if P1 were called after P2 returns, it would be an erroneous use of a dangling pointer.

For objects, it is possible to use \texttt{Unrestricted\_Access} for any type. However, if the result is of an access-to-unconstrained array subtype, then the resulting pointer has the same scope as the context of the attribute, and must not be returned to some enclosing scope. For instance, if a function uses \texttt{Unrestricted\_Access} to create an access-to-unconstrained-array and returns that value to the caller, the result will involve dangling pointers. In addition, it is only valid to create pointers to unconstrained arrays using this attribute if the pointer has the normal default ‘fat’ representation where a pointer has two components, one points to the array and one points to the bounds. If a size clause is used to force ‘thin’ representation for a pointer to unconstrained where there is only space for a single pointer, then the resulting pointer is not usable.

In the simple case where a direct use of \texttt{Unrestricted\_Access} attempts to make a thin pointer for a non-aliased object, the compiler will reject the use as illegal, as shown in the following example:

\begin{verbatim}
with System; use System;
procedure SliceUA2 is
  type A is access all String;
  for A'Size use Standard'Address_Size;

  procedure P (Arg : A) is
  begin
    null;
  end P;

  X : String := "hello world!";
  X2 : aliased String := "hello world!";

  AV : A := X'Unrestricted_Access;  -- ERROR
  | >>> illegal use of Unrestricted\_Access attribute
  >>> attempt to generate thin pointer to unaliased object
\end{verbatim}
begin
  P (X'Unrestricted_Access); -- ERROR
  | >>> illegal use of Unrestricted_Access attribute
  >>> attempt to generate thin pointer to unaliased object
  P (X(7 .. 12)'Unrestricted_Access); -- ERROR
  | >>> illegal use of Unrestricted_Access attribute
  >>> attempt to generate thin pointer to unaliased object
  P (X2'Unrestricted_Access); -- OK
end;

but other cases cannot be detected by the compiler, and are considered to be erroneous. Consider the following example:

with System; use System;
with System; use System;
procedure SliceUA is
  type AF is access all String;
  type A is access all String;
  for A'Size use Standard'Address_Size;

  procedure P (Arg : A) is
    begin
      if Arg'Length /= 6 then
        raise Program_Error;
      end if;
    end P;

  X : String := "hello world!";
  Y : AF := X (7 .. 12)'Unrestricted_Access;

begin
  P (A (Y));
end;

A normal unconstrained array value or a constrained array object marked as aliased has the bounds in memory just before the array, so a thin pointer can retrieve both the data and the bounds. But in this case, the non-aliased object X does not have the bounds before the string. If the size clause for type A were not present, then the pointer would be a fat pointer, where one component is a pointer to the bounds, and all would be well. But with the size clause present, the conversion from fat pointer to thin pointer in the call loses the bounds, and so this is erroneous, and the program likely raises a Program_Error exception.

In general, it is advisable to completely avoid mixing the use of thin pointers and the use of Unrestricted_Access where the designated type is an unconstrained array. The use of
thin pointers should be restricted to cases of porting legacy code that implicitly assumes
the size of pointers, and such code should not in any case be using this attribute.

Another erroneous situation arises if the attribute is applied to a constant. The resulting
pointer can be used to access the constant, but the effect of trying to modify a constant in
this manner is not well-defined. Consider this example:

\[
\begin{align*}
P & : \text{constant Integer} := 4; \\
type R & : \text{access all Integer}; \\
RV & : R := P'Unrestricted\_Access; \\
& \ldots \\
RV.all & := 3;
\end{align*}
\]

Here we attempt to modify the constant \( P \) from 4 to 3, but the compiler may or may not
notice this attempt, and subsequent references to \( P \) may yield either the value 3 or the value
4 or the assignment may blow up if the compiler decides to put \( P \) in read-only memory.

One particular case where \textit{Unrestricted\_Access} can be used in this way is to modify the
value of an \textit{in} parameter:

\[
\begin{align*}
\text{procedure K (S : in String) is} \\
& \quad \text{type R is access all Character;} \\
& \quad RV : R := S (3)'Unrestricted\_Access; \\
& \quad \text{begin} \\
& \quad \quad RV.all := 'a'; \\
& \quad \text{end;}
\end{align*}
\]

In general this is a risky approach. It may appear to “work” but such uses of \textit{Unrestricted\_Access} are potentially non-portable, even from one version of GNAT to another, so are best
avoided if possible.

4.72 Attribute Update

The \textit{Update} attribute creates a copy of an array or record value with one or more modified
components. The syntax is:

\[
\begin{align*}
\text{PREFIX'}\text{Update ( RECORD\_COMPONENT\_ASSOCIATION\_LIST )} \\
\text{PREFIX'}\text{Update ( ARRAY\_COMPONENT\_ASSOCIATION {, ARRAY\_COMPONENT\_ASSOCIATION } )} \\
\text{PREFIX'}\text{Update ( MULTIDIMENSIONAL\_ARRAY\_COMPONENT\_ASSOCIATION} \\
\quad \{-, MULTIDIMENSIONAL\_ARRAY\_COMPONENT\_ASSOCIATION \})
\end{align*}
\]

\[
\begin{align*}
\text{MULTIDIMENSIONAL\_ARRAY\_COMPONENT\_ASSOCIATION} & ::= \text{INDEX\_EXPRESSION\_LIST\_LIST} \Rightarrow \text{EXPRESSION} \\
\text{INDEX\_EXPRESSION\_LIST\_LIST} & ::= \text{INDEX\_EXPRESSION\_LIST} \{ | \text{INDEX\_EXPRESSION\_LIST} \}
\end{align*}
\]

where \textit{PREFIX} is the name of an array or record object, the association list in parentheses
does not contain an \textit{others} choice and the box symbol <> may not appear in any expression.
The effect is to yield a copy of the array or record value which is unchanged apart from
the components mentioned in the association list, which are changed to the indicated value.
The original value of the array or record value is not affected. For example:

\[
\begin{align*}
\text{type Arr is Array (1 .. 5) of Integer;} \\
& \ldots \\
Avar1 : Arr := (1,2,3,4,5);
\end{align*}
\]
Avar2 : Arr := Avar1'Update (2 => 10, 3 .. 4 => 20);
yields a value for Avar2 of 1,10,20,20,5 with Avar1 begin unmodified. Similarly:

```haskell
type Rec is A, B, C : Integer;
...
Rvar1 : Rec := (A => 1, B => 2, C => 3);
Rvar2 : Rec := Rvar1'Update (B => 20);
yields a value for Rvar2 of (A = 1, B = 20, C = 3), with Rvar1 being unmodified. Note
that the value of the attribute reference is computed completely before it is used. This
means that if you write:

Avar1 := Avar1'Update (1 => 10, 2 => Function_Call);
then the value of Avar1 is not modified if Function_Call raises an exception, unlike the
effect of a series of direct assignments to elements of Avar1. In general this requires that
two extra complete copies of the object are required, which should be kept in mind when
considering efficiency.

The Update attribute cannot be applied to prefixes of a limited type, and cannot reference
discriminants in the case of a record type. The accessibility level of an Update attribute
result object is defined as for an aggregate.

In the record case, no component can be mentioned more than once. In the array case, two
overlapping ranges can appear in the association list, in which case the modifications are
processed left to right.

Multi-dimensional arrays can be modified, as shown by this example:

```haskell
A : array (1 .. 10, 1 .. 10) of Integer;
...
A := A'Update ((1, 2) => 20, (3, 4) => 30);
```
which changes element (1,2) to 20 and (3,4) to 30.

### 4.73 Attribute Valid_Value

The 'Valid_Value attribute is defined for enumeration types other than those in package
Standard. This attribute is a function that takes a String, and returns Boolean. T'Valid_
Value (S) returns True if and only if T'Value (S) would not raise Constraint_Error.

### 4.74 Attribute Valid_Scalars

The 'Valid_Scalars attribute is intended to make it easier to check the validity of scalar
subcomponents of composite objects. The attribute is defined for any prefix P which denotes
an object. Prefix P can be any type except for tagged private or Unchecked_Union types.
The value of the attribute is of type Boolean.

P'Valid_Scalars yields True if and only if the evaluation of C'Valid yields True for every
scalar subcomponent C of P, or if P has no scalar subcomponents. Attribute 'Valid_Scalars
is equivalent to attribute 'Valid for scalar types.

It is not specified in what order the subcomponents are checked, nor whether any more are
checked after any one of them is determined to be invalid. If the prefix P is of a class-wide
type T'Class (where T is the associated specific type), or if the prefix P is of a specific
tagged type \( T \), then only the subcomponents of \( T \) are checked; in other words, components of extensions of \( T \) are not checked even if \( T'\text{Class (P)'}\text{Tag} /= T'\text{Tag} \).

The compiler will issue a warning if it can be determined at compile time that the prefix of the attribute has no scalar subcomponents.

Note: \texttt{Valid_Scalars} can generate a lot of code, especially in the case of a large variant record. If the attribute is called in many places in the same program applied to objects of the same type, it can reduce program size to write a function with a single use of the attribute, and then call that function from multiple places.

\section*{4.75 Attribute VADS\_Size}

The 'VADS\_Size attribute is intended to make it easier to port legacy code which relies on the semantics of 'Size as implemented by the VADS Ada 83 compiler. GNAT makes a best effort at duplicating the same semantic interpretation. In particular, 'VADS\_Size applied to a predefined or other primitive type with no Size clause yields the Object\_Size (for example, \texttt{Natural'Size} is 32 rather than 31 on typical machines). In addition 'VADS\_Size applied to an object gives the result that would be obtained by applying the attribute to the corresponding type.

\section*{4.76 Attribute Value\_Size}

\texttt{type'Value\_Size} is the number of bits required to represent a value of the given subtype. It is the same as \texttt{type'Size}, but, unlike \texttt{Size}, may be set for non-first subtypes.

\section*{4.77 Attribute Wchar\_T\_Size}

\texttt{Standard'Wchar\_T\_Size} (\texttt{Standard} is the only allowed prefix) provides the size in bits of the \texttt{C wchar_t} type primarily for constructing the definition of this type in package \texttt{Interfaces.C}. The result is a static constant.

\section*{4.78 Attribute Word\_Size}

\texttt{Standard'Word\_Size} (\texttt{Standard} is the only allowed prefix) provides the value \texttt{System.Word\_Size}. The result is a static constant.
5 Standard and Implementation Defined Restrictions

All Ada Reference Manual-defined Restriction identifiers are implemented:

* language-defined restrictions (see 13.12.1)
* tasking restrictions (see D.7)
* high integrity restrictions (see H.4)

GNAT implements additional restriction identifiers. All restrictions, whether language defined or GNAT-specific, are listed in the following.

5.1 Partition-Wide Restrictions

There are two separate lists of restriction identifiers. The first set requires consistency throughout a partition (in other words, if the restriction identifier is used for any compilation unit in the partition, then all compilation units in the partition must obey the restriction).

5.1.1 Immediate_Reclamation

[RM H.4] This restriction ensures that, except for storage occupied by objects created by allocators and not deallocated via unchecked deallocation, any storage reserved at run time for an object is immediately reclaimed when the object no longer exists.

5.1.2 Max_Asynchronous_Select_Nesting

[RM D.7] Specifies the maximum dynamic nesting level of asynchronous selects. Violations of this restriction with a value of zero are detected at compile time. Violations of this restriction with values other than zero cause Storage_Error to be raised.

5.1.3 Max_EntryQueueLength

[RM D.7] This restriction is a declaration that any protected entry compiled in the scope of the restriction has at most the specified number of tasks waiting on the entry at any one time, and so no queue is required. Note that this restriction is checked at run time. Violation of this restriction results in the raising of Program_Error exception at the point of the call.

The restriction Max_EntryQueueDepth is recognized as a synonym for Max_EntryQueueLength. This is retained for historical compatibility purposes (and a warning will be generated for its use if warnings on obsolescent features are activated).

5.1.4 Max_Protected_Entries

[RM D.7] Specifies the maximum number of entries per protected type. The bounds of every entry family of a protected unit shall be static, or shall be defined by a discriminant of a subtype whose corresponding bound is static.

5.1.5 Max_Select_Alternatives

[RM D.7] Specifies the maximum number of alternatives in a selective accept.
5.1.6 Max_Storage_At_Blocking
[RM D.7] Specifies the maximum portion (in storage elements) of a task’s Storage_Size that can be retained by a blocked task. A violation of this restriction causes Storage_Error to be raised.

5.1.7 Max_Task_Entrees
[RM D.7] Specifies the maximum number of entries per task. The bounds of every entry family of a task unit shall be static, or shall be defined by a discriminant of a subtype whose corresponding bound is static.

5.1.8 Max_Tasks
[RM D.7] Specifies the maximum number of task that may be created, not counting the creation of the environment task. Violations of this restriction with a value of zero are detected at compile time. Violations of this restriction with values other than zero cause Storage_Error to be raised.

5.1.9 No_Abort_Statements
[RM D.7] There are no abort_statements, and there are no calls to Task_Identification.Abort_Task.

5.1.10 No_Access_Parameter_Allocators
[RM H.4] This restriction ensures at compile time that there are no occurrences of an allocator as the actual parameter to an access parameter.

5.1.11 No_Access_Subprograms
[RM H.4] This restriction ensures at compile time that there are no declarations of access-to-subprogram types.

5.1.12 No_Allocators
[RM H.4] This restriction ensures at compile time that there are no occurrences of an allocator.

5.1.13 No_Anonymous_Allocators
[RM H.4] This restriction ensures at compile time that there are no occurrences of an allocator of anonymous access type.

5.1.14 No_Asynchronous_Control
[RM J.13] This restriction ensures at compile time that there are no semantic dependences on the predefined package Asynchronous_Task_Control.

5.1.15 No_Calendar
[GNAT] This restriction ensures at compile time that there are no semantic dependences on package Calendar.

5.1.16 No_Coextensions
[RM H.4] This restriction ensures at compile time that there are no coextensions. See 3.10.2.
5.1.17 **No_DefaultInitialization**

[GNAT] This restriction prohibits any instance of default initialization of variables. The binder implements a consistency rule which prevents any unit compiled without the restriction from with'ing a unit with the restriction (this allows the generation of initialization procedures to be skipped, since you can be sure that no call is ever generated to an initialization procedure in a unit with the restriction active). If used in conjunction with Initialize_Scalars or Normalize_Scalars, the effect is to prohibit all cases of variables declared without a specific initializer (including the case of OUT scalar parameters).

5.1.18 **No_Delay**

[RM H.4] This restriction ensures at compile time that there are no delay statements and no semantic dependences on package Calendar.

5.1.19 **No_Dependence**

[RM 13.12.1] This restriction ensures at compile time that there are no dependences on a library unit. For GNAT, this includes implicit implementation dependences on units of the runtime library that are created by the compiler to support specific constructs of the language.

5.1.20 **No_Direct_Boolen_Operators**

[GNAT] This restriction ensures that no logical operators (and/or/xor) are used on operands of type Boolean (or any type derived from Boolean). This is intended for use in safety critical programs where the certification protocol requires the use of short-circuit (and then, or else) forms for all composite boolean operations.

5.1.21 **No_Dispach**

[RM H.4] This restriction ensures at compile time that there are no occurrences of T’Class, for any (tagged) subtype T.

5.1.22 **No_Dispaching_Calls**

[GNAT] This restriction ensures at compile time that the code generated by the compiler involves no dispatching calls. The use of this restriction allows the safe use of record extensions, classwide membership tests and other classwide features not involving implicit dispatching. This restriction ensures that the code contains no indirect calls through a dispatching mechanism. Note that this includes internally-generated calls created by the compiler, for example in the implementation of class-wide objects assignments. The membership test is allowed in the presence of this restriction, because its implementation requires no dispatching. This restriction is comparable to the official Ada restriction No_Dispatch except that it is a bit less restrictive in that it allows all classwide constructs that do not imply dispatching. The following example indicates constructs that violate this restriction.

```ada
package Pkg is
   type T is tagged record
      Data : Natural;
   end record;
   procedure P (X : T);
```

```
type DT is new T with record
  More_Data : Natural;
end record;
procedure Q (X : DT);
end Pkg;

with Pkg; use Pkg;
procedure Example is
  procedure Test (O : T'Class) is
    N : Natural := O'Size; -- Error: Dispatching call
    C : T'Class := O; -- Error: implicit Dispatching Call
  begin
    if O in DT'Class then -- OK : Membership test
      Q (DT (O)); -- OK : Type conversion plus direct call
    else
      P (O); -- Error: Dispatching call
    end if;
  end Test;

  Obj : DT;
  begin
    P (Obj); -- OK : Direct call
    P (T (Obj)); -- OK : Type conversion plus direct call
    P (T'Class (Obj)); -- Error: Dispatching call
    Test (Obj); -- OK : Type conversion
    if Obj in T'Class then -- OK : Membership test
      null;
    end if;
  end Example;

5.1.23 No_Dynamic_Attachment
[RM D.7] This restriction ensures that there is no call to any of the operations defined in package Ada.Interrupts (Is_Reserved, Is_Attached, Current_Handler, Attach_Handler, Exchange_Handler, Detach_Handler, and Reference).
The restriction No_Dynamic_Interrupts is recognized as a synonym for No_Dynamic_Attachment. This is retained for historical compatibility purposes (and a warning will be generated for its use if warnings on obsolescent features are activated).

5.1.24 No_Dynamic_Priorities
[RM D.7] There are no semantic dependencies on the package Dynamic_Priorities.

5.1.25 No_Entry_Calls_In_Elaboration_Code
[GNAT] This restriction ensures at compile time that no task or protected entry calls are made during elaboration code. As a result of the use of this restriction, the compiler can
assume that no code past an accept statement in a task can be executed at elaboration
time.

5.1.26 No Enumeration Maps

[GNAT] This restriction ensures at compile time that no operations requiring enumeration
maps are used (that is Image and Value attributes applied to enumeration types).

5.1.27 No Exception Handlers

[GNAT] This restriction ensures at compile time that there are no explicit exception han-
dlers. It also indicates that no exception propagation will be provided. In this mode,
extceptions may be raised but will result in an immediate call to the last chance handler, a
routine that the user must define with the following profile:

```lang-pl
procedure Last_Chance_Handler
  (Source_Location : System.Address; Line : Integer);
pragma Export (C, Last_Chance_Handler,
       "__gnat_last_chance_handler");
```

The parameter is a C null-terminated string representing a message to be associated with the
exception (typically the source location of the raise statement generated by the compiler).
The Line parameter when nonzero represents the line number in the source program where
the raise occurs.

5.1.28 No Exception Propagation

[GNAT] This restriction guarantees that exceptions are never propagated to an outer sub-
program scope. The only case in which an exception may be raised is when the handler
is statically in the same subprogram, so that the effect of a raise is essentially like a goto
statement. Any other raise statement (implicit or explicit) will be considered unhandled.
Exception handlers are allowed, but may not contain an exception occurrence identifier (ex-
ception choice). In addition, use of the package GNAT.Current_Exception is not permitted,
and reraise statements (raise with no operand) are not permitted.

5.1.29 No Exception Registration

[GNAT] This restriction ensures at compile time that no stream operations for types Excep-
tion_Id or Exception_Occurrence are used. This also makes it impossible to pass exceptions
to or from a partition with this restriction in a distributed environment. If this restric-
tion is active, the generated code is simplified by omitting the otherwise-required global
registration of exceptions when they are declared.

5.1.30 No Exceptions

[RM H.4] This restriction ensures at compile time that there are no raise statements and no
exception handlers and also suppresses the generation of language-defined run-time checks.

5.1.31 No Finalization

[GNAT] This restriction disables the language features described in chapter 7.6 of the Ada
2005 RM as well as all form of code generation performed by the compiler to support these
features. The following types are no longer considered controlled when this restriction is in
effect:
* Ada.Finalization.Controlled
* Ada.Finalization.Limited_Controlled
* Derivations from Controlled or Limited_Controlled
* Class-wide types
* Protected types
* Task types
* Array and record types with controlled components

The compiler no longer generates code to initialize, finalize or adjust an object or a nested component, either declared on the stack or on the heap. The deallocation of a controlled object no longer finalizes its contents.

5.1.32 No_Fixed_Point
[RM H.4] This restriction ensures at compile time that there are no occurrences of fixed point types and operations.

5.1.33 No_Floating_Point
[RM H.4] This restriction ensures at compile time that there are no occurrences of floating point types and operations.

5.1.34 No_Implicit_Conditionals
[GNAT] This restriction ensures that the generated code does not contain any implicit conditionals, either by modifying the generated code where possible, or by rejecting any construct that would otherwise generate an implicit conditional. Note that this check does not include run time constraint checks, which on some targets may generate implicit conditionals as well. To control the latter, constraint checks can be suppressed in the normal manner. Constructs generating implicit conditionals include comparisons of composite objects and the Max/Min attributes.

5.1.35 No_Implicit_Dynamic_Code
[GNAT] This restriction prevents the compiler from building ‘trampolines’. This is a structure that is built on the stack and contains dynamic code to be executed at run time. On some targets, a trampoline is built for the following features: Access, Unrestricted_Access, or Address of a nested subprogram; nested task bodies; primitive operations of nested tagged types. Trampolines do not work on machines that prevent execution of stack data. For example, on windows systems, enabling DEP (data execution protection) will cause trampolines to raise an exception. Trampolines are also quite slow at run time.

On many targets, trampolines have been largely eliminated. Look at the version of system.ads for your target — if it has Always_Compatible_Rep equal to False, then trampolines are largely eliminated. In particular, a trampoline is built for the following features: Address of a nested subprogram; Access or Unrestricted_Access of a nested subprogram, but only if pragma Favor_Top_Level applies, or the access type has a foreign-language convention; primitive operations of nested tagged types.

5.1.36 No_Implicit_Heap_Allocations
[RM D.7] No constructs are allowed to cause implicit heap allocation.
5.1.37 No_Implicit_Protected_Object_Allocations

[GNAT] No constructs are allowed to cause implicit heap allocation of a protected object.

5.1.38 No_Implicit_Task_Allocations

[GNAT] No constructs are allowed to cause implicit heap allocation of a task.

5.1.39 No_Initialize_Scalars

[GNAT] This restriction ensures that no unit in the partition is compiled with pragma Initialize_Scalars. This allows the generation of more efficient code, and in particular eliminates dummy null initialization routines that are otherwise generated for some record and array types.

5.1.40 No_IO

[RM H.4] This restriction ensures at compile time that there are no dependences on any of the library units Sequential_IO, Direct_IO, Text_IO, Wide_Text_IO, Wide_Wide_Text_IO, or Stream_IO.

5.1.41 No_Local_Allocators

[RM H.4] This restriction ensures at compile time that there are no occurrences of an allocator in subprograms, generic subprograms, tasks, and entry bodies.

5.1.42 No_Local_Protected_Objects

[RM D.7] This restriction ensures at compile time that protected objects are only declared at the library level.

5.1.43 No_Local_Tagged_TYPES

[GNAT] This restriction ensures at compile time that tagged types are only declared at the library level.

5.1.44 No_Local_Timing_Events


5.1.45 No_Long_Long_Integers

[GNAT] This partition-wide restriction forbids any explicit reference to type Standard.Long_Long_Integer, and also forbids declaring range types whose implicit base type is Long_Long_Integer, and modular types whose size exceeds Long_Integer’Size.

5.1.46 No_Multiple_Elaboration

[GNAT] When this restriction is active and the static elaboration model is used, and -fpreserve-control-flow is not used, the compiler is allowed to suppress the elaboration counter normally associated with the unit, even if the unit has elaboration code. This counter is typically used to check for access before elaboration and to control multiple elaboration attempts. If the restriction is used, then the situations in which multiple elaboration is possible, including non-Ada main programs and Stand Alone libraries, are not permitted and will be diagnosed by the Ada binder.
5.1.47 No_Nested_Finalization
[RM D.7] All objects requiring finalization are declared at the library level.

5.1.48 NoProtected_Type_Allocators
[RM D.7] This restriction ensures at compile time that there are no allocator expressions that attempt to allocate protected objects.

5.1.49 No_Protected_Types
[RM H.4] This restriction ensures at compile time that there are no declarations of protected types or protected objects.

5.1.50 No_Recursion
[RM H.4] A program execution is erroneous if a subprogram is invoked as part of its execution.

5.1.51 No_Reentrancy
[RM H.4] A program execution is erroneous if a subprogram is executed by two tasks at the same time.

5.1.52 No_Relative_Delay
[RM D.7] This restriction ensures at compile time that there are no delay relative statements and prevents expressions such as delay 1.23; from appearing in source code.

5.1.53 No_Requeue_Statements
[RM D.7] This restriction ensures at compile time that no requeue statements are permitted and prevents keyword requeue from being used in source code.

The restriction No_Requeue is recognized as a synonym for No_Requeue_Statements. This is retained for historical compatibility purposes (and a warning will be generated for its use if warnings on obsolescent features are activated).

5.1.54 No_Secondary_Stack
[GNAT] This restriction ensures at compile time that the generated code does not contain any reference to the secondary stack. The secondary stack is used to implement functions returning unconstrained objects (arrays or records) on some targets. Suppresses the allocation of secondary stacks for tasks (excluding the environment task) at run time.

5.1.55 No_Select_Statements
[RM D.7] This restriction ensures at compile time no select statements of any kind are permitted, that is the keyword select may not appear.

5.1.56 No_Specific_Termination_Handler
[RM D.7] There are no calls to Ada.Task_Termination.Set_Specific_Handler or to Ada.Task_Termination.Specific_Handler.
5.1.57 NoSpecification_of_Aspect

[RM 13.12.1] This restriction checks at compile time that no aspect specification, attribute definition clause, or pragma is given for a given aspect.

5.1.58 NoStandard_Allocators_After_Elaboration

[RM D.7] Specifies that an allocator using a standard storage pool should never be evaluated at run time after the elaboration of the library items of the partition has completed. Otherwise, Storage_Error is raised.

5.1.59 NoStandard_Storage_Pools

[GNAT] This restriction ensures at compile time that no access types use the standard default storage pool. Any access type declared must have an explicit Storage_Pool attribute defined specifying a user-defined storage pool.

5.1.60 No_Stream_Optimizations

[GNAT] This restriction affects the performance of stream operations on types String, Wide_String and Wide_Wide_String. By default, the compiler uses block reads and writes when manipulating String objects due to their superior performance. When this restriction is in effect, the compiler performs all IO operations on a per-character basis.

5.1.61 No_Streams

[GNAT] This restriction ensures at compile/bind time that there are no stream objects created and no use of stream attributes. This restriction does not forbid dependences on the package Ada.Streams. So it is permissible to with AdaStreams (or another package that does so itself) as long as no actual stream objects are created and no stream attributes are used.

Note that the use of restriction allows optimization of tagged types, since they do not need to worry about dispatching stream operations. To take maximum advantage of this space-saving optimization, any unit declaring a tagged type should be compiled with the restriction, though this is not required.

5.1.62 No_Tagged_Type_Registration

[GNAT] If this restriction is active, then class-wide streaming attributes are not supported. In addition, the subprograms in Ada.Tags are not supported. If this restriction is active, the generated code is simplified by omitting the otherwise-required global registration of tagged types when they are declared. This restriction may be necessary in order to also apply the No_Elaboration_Code restriction.

5.1.63 No_Task_Allocators

[RM D.7] There are no allocators for task types or types containing task subcomponents.

5.1.64 No_Task_At_Interrupt_Priority

[GNAT] This restriction ensures at compile time that there is no Interrupt_Priority aspect or pragma for a task or a task type. As a consequence, the tasks are always created with a priority below that an interrupt priority.
5.1.65  **No_Task_Attributes_Package**

[GNAT] This restriction ensures at compile time that there are no implicit or explicit dependencies on the package `Ada.Task_Attributes`.

The restriction `No_Task_Attributes` is recognized as a synonym for `No_Task_Attributes_Package`. This is retained for historical compatibility purposes (and a warning will be generated for its use if warnings on obsolescent features are activated).

5.1.66  **No_Task_Hierarchy**

[RM D.7] All (non-environment) tasks depend directly on the environment task of the partition.

5.1.67  **No_Task_Termination**

[RM D.7] Tasks that terminate are erroneous.

5.1.68  **No_Tasking**

[GNAT] This restriction prevents the declaration of tasks or task types throughout the partition. It is similar in effect to the use of `Max_Tasks => 0` except that violations are caught at compile time and cause an error message to be output either by the compiler or binder.

5.1.69  **No_Terminate_Alternatives**

[RM D.7] There are no selective accepts with terminate alternatives.

5.1.70  **No_Unchecked_Access**

[RM H.4] This restriction ensures at compile time that there are no occurrences of the `Unchecked_Access` attribute.

5.1.71  **No_Unchecked_Conversion**

[RM J.13] This restriction ensures at compile time that there are no semantic dependences on the predefined generic function `Unchecked_Conversion`.

5.1.72  **No_Unchecked_Deallocation**

[RM J.13] This restriction ensures at compile time that there are no semantic dependences on the predefined generic procedure `Unchecked_Deallocation`.

5.1.73  **No_Use_Of_Entity**

[GNAT] This restriction ensures at compile time that there are no references to the entity given in the form

\[
\text{No_Use_Of_Entity} => \text{Name}
\]

where `Name` is the fully qualified entity, for example

\[
\text{No_Use_Of_Entity} => \text{Ada.Text_IO.Put_Line}
\]

5.1.74  **Pure_Barrers**

[GNAT] This restriction ensures at compile time that protected entry barriers are restricted to:
* components of the protected object (excluding selection from dereferences),
* constant declarations,
* named numbers,
* enumeration literals,
* integer literals,
* real literals,
* character literals,
* implicitly defined comparison operators,
* uses of the Standard."not" operator,
* short-circuit operator,
* the Count attribute

This restriction is a relaxation of the Simple_Barriers restriction, but still ensures absence of side effects, exceptions, and recursion during the evaluation of the barriers.

### 5.1.75 Simple_Barriers

[RM D.7] This restriction ensures at compile time that barriers in entry declarations for protected types are restricted to either static boolean expressions or references to simple boolean variables defined in the private part of the protected type. No other form of entry barriers is permitted.

The restriction Boolean_Entry_Barriers is recognized as a synonym for Simple_Barriers. This is retained for historical compatibility purposes (and a warning will be generated for its use if warnings on obsolescent features are activated).

### 5.1.76 Static_Priorities

[GNAT] This restriction ensures at compile time that all priority expressions are static, and that there are no dependences on the package Ada.Dynamic_Priorities.

### 5.1.77 Static_Storage_Size

[GNAT] This restriction ensures at compile time that any expression appearing in a Storage_Size pragma or attribute definition clause is static.

### 5.2 Program Unit Level Restrictions

The second set of restriction identifiers does not require partition-wide consistency. The restriction may be enforced for a single compilation unit without any effect on any of the other compilation units in the partition.

#### 5.2.1 No_Elaboration_Code

[GNAT] This restriction ensures at compile time that no elaboration code is generated. Note that this is not the same condition as is enforced by pragma Preelaborate. There are cases in which pragma Preelaborate still permits code to be generated (e.g., code to initialize a large array to all zeroes), and there are cases of units which do not meet the requirements for pragma Preelaborate, but for which no elaboration code is generated. Generally, it is the case that preelaborable units will meet the restrictions, with the exception of large
aggregates initialized with an others clause, and exception declarations (which generate calls to a run-time registry procedure). This restriction is enforced on a unit by unit basis, it need not be obeyed consistently throughout a partition.

In the case of aggregates with others, if the aggregate has a dynamic size, there is no way to eliminate the elaboration code (such dynamic bounds would be incompatible with Preelaborate in any case). If the bounds are static, then use of this restriction actually modifies the code choice of the compiler to avoid generating a loop, and instead generate the aggregate statically if possible, no matter how many times the data for the others clause must be repeatedly generated.

It is not possible to precisely document the constructs which are compatible with this restriction, since, unlike most other restrictions, this is not a restriction on the source code, but a restriction on the generated object code. For example, if the source contains a declaration:

\[
\text{Val : constant Integer := X;}
\]

where \( X \) is not a static constant, it may be possible, depending on complex optimization circuitry, for the compiler to figure out the value of \( X \) at compile time, in which case this initialization can be done by the loader, and requires no initialization code. It is not possible to document the precise conditions under which the optimizer can figure this out.

Note that this the implementation of this restriction requires full code generation. If it is used in conjunction with “semantics only” checking, then some cases of violations may be missed.

When this restriction is active, we are not requesting control-flow preservation with -fpreserve-control-flow, and the static elaboration model is used, the compiler is allowed to suppress the elaboration counter normally associated with the unit. This counter is typically used to check for access before elaboration and to control multiple elaboration attempts.

5.2.2 No Dynamic Accessibility Checks

[GNAT] No dynamic accessibility checks are generated when this restriction is in effect. Instead, dangling references are prevented via more conservative compile-time checking. More specifically, existing compile-time checks are enforced but with more conservative assumptions about the accessibility levels of the relevant entities. These conservative assumptions eliminate the need for dynamic accessibility checks.

These new rules for computing (at compile-time) the accessibility level of an anonymous access type \( T \) are as follows:

* If \( T \) is a function result type then, from the caller’s perspective, its level is that of the innermost master enclosing the function call. From the callee’s perspective, the level of parameters and local variables of the callee is statically deeper than the level of \( T \).

For any other accessibility level \( L \) such that the level of parameters and local variables of the callee is statically deeper than \( L \), the level of \( T \) (from the callee’s perspective) is also statically deeper than \( L \).

* If \( T \) is the type of a formal parameter then, from the caller’s perspective, its level is at least as deep as that of the type of the corresponding actual parameter (whatever that actual parameter might be). From the callee’s perspective, the level of parameters and local variables of the callee is statically deeper than the level of \( T \).
* If T is the type of a discriminant then its level is that of the discriminated type.
* If T is the type of a stand-alone object then its level is the level of the object.
* In all other cases, the level of T is as defined by the existing rules of Ada.

5.2.3 No_Dynamic_Sized_Objects

[GNAT] This restriction disallows certain constructs that might lead to the creation of dynamic-sized composite objects (or array or discriminated type). An array subtype indication is illegal if the bounds are not static or references to discriminants of an enclosing type. A discriminated subtype indication is illegal if the type has discriminant-dependent array components or a variant part, and the discriminants are not static. In addition, array and record aggregates are illegal in corresponding cases. Note that this restriction does not forbid access discriminants. It is often a good idea to combine this restriction with No_Secondary_Stack.

5.2.4 No_Entry_Queue

[GNAT] This restriction is a declaration that any protected entry compiled in the scope of the restriction has at most one task waiting on the entry at any one time, and so no queue is required. This restriction is not checked at compile time. A program execution is erroneous if an attempt is made to queue a second task on such an entry.

5.2.5 No_Implementation_Aspect_Specifications

[RM 13.12.1] This restriction checks at compile time that no GNAT-defined aspects are present. With this restriction, the only aspects that can be used are those defined in the Ada Reference Manual.

5.2.6 No_Implementation_Attributes

[RM 13.12.1] This restriction checks at compile time that no GNAT-defined attributes are present. With this restriction, the only attributes that can be used are those defined in the Ada Reference Manual.

5.2.7 No_Implementation_Identifiers

[RM 13.12.1] This restriction checks at compile time that no implementation-defined identifiers (marked with pragma Implementation_DEFINED) occur within language-defined packages.

5.2.8 No_Implementation_Pragmas

[RM 13.12.1] This restriction checks at compile time that no GNAT-defined pragmas are present. With this restriction, the only pragmas that can be used are those defined in the Ada Reference Manual.

5.2.9 No_Implementation_Restrictions

[GNAT] This restriction checks at compile time that no GNAT-defined restriction identifiers (other than No_Implementation_Restrictions itself) are present. With this restriction, the only other restriction identifiers that can be used are those defined in the Ada Reference Manual.
5.2.10 No_Implementation_Units

[RM 13.12.1] This restriction checks at compile time that there is no mention in the context clause of any implementation-defined descendants of packages Ada, Interfaces, or System.

5.2.11 No_Implicit_Aliasing

[GNAT] This restriction, which is not required to be partition-wide consistent, requires an explicit aliased keyword for an object to which ‘Access, ‘Unchecked_Access, or ‘Address is applied, and forbids entirely the use of the ‘Unrestricted_Access attribute for objects. Note: the reason that Unrestricted_Access is forbidden is that it would require the prefix to be aliased, and in such cases, it can always be replaced by the standard attribute Unchecked_Access which is preferable.

5.2.12 No_Implicit_Loops

[GNAT] This restriction ensures that the generated code of the unit marked with this restriction does not contain any implicit for loops, either by modifying the generated code where possible, or by rejecting any construct that would otherwise generate an implicit for loop. If this restriction is active, it is possible to build large array aggregates with all static components without generating an intermediate temporary, and without generating a loop to initialize individual components. Otherwise, a loop is created for arrays larger than about 5000 scalar components. Note that if this restriction is set in the spec of a package, it will not apply to its body.

5.2.13 No_Obsolescent_Features

[RM 13.12.1] This restriction checks at compile time that no obsolescent features are used, as defined in Annex J of the Ada Reference Manual.

5.2.14 No_Wide_Characters

[GNAT] This restriction ensures at compile time that no uses of the types Wide_Character or Wide_String or corresponding wide wide types appear, and that no wide or wide wide string or character literals appear in the program (that is literals representing characters not in type Character).

5.2.15 Static_Dispacth_Tables

[GNAT] This restriction checks at compile time that all the artifacts associated with dispatch tables can be placed in read-only memory.

5.2.16 SPARK_05

[GNAT] This restriction no longer has any effect and is superseded by SPARK 2014, whose restrictions are checked by the tool GNATprove. To check that a codebase respects SPARK 2014 restrictions, mark the code with pragma or aspect SPARK_Mode, and run the tool GNATprove at Stone assurance level, as follows:

```
    gnatprove -P project.gpr --mode=stone
```

or equivalently:

```
    gnatprove -P project.gpr --mode=check_all
```
6 Implementation Advice

The main text of the Ada Reference Manual describes the required behavior of all Ada compilers, and the GNAT compiler conforms to these requirements.

In addition, there are sections throughout the Ada Reference Manual headed by the phrase 'Implementation advice'. These sections are not normative, i.e., they do not specify requirements that all compilers must follow. Rather they provide advice on generally desirable behavior. They are not requirements, because they describe behavior that cannot be provided on all systems, or may be undesirable on some systems.

As far as practical, GNAT follows the implementation advice in the Ada Reference Manual. Each such RM section corresponds to a section in this chapter whose title specifies the RM section number and paragraph number and the subject of the advice. The contents of each section consists of the RM text within quotation marks, followed by the GNAT interpretation of the advice. Most often, this simply says 'followed', which means that GNAT follows the advice. However, in a number of cases, GNAT deliberately deviates from this advice, in which case the text describes what GNAT does and why.

6.1 RM 1.1.3(20): Error Detection

“If an implementation detects the use of an unsupported Specialized Needs Annex feature at run time, it should raise Program_Error if feasible.”

Not relevant. All specialized needs annex features are either supported, or diagnosed at compile time.

6.2 RM 1.1.3(31): Child Units

“If an implementation wishes to provide implementation-defined extensions to the functionality of a language-defined library unit, it should normally do so by adding children to the library unit.”

Followed.

6.3 RM 1.1.5(12): Bounded Errors

“If an implementation detects a bounded error or erroneous execution, it should raise Program_Error.”

Followed in all cases in which the implementation detects a bounded error or erroneous execution. Not all such situations are detected at runtime.

6.4 RM 2.8(16): Pragmas

“Normally, implementation-defined pragmas should have no semantic effect for error-free programs; that is, if the implementation-defined pragmas are removed from a working program, the program should still be legal, and should still have the same semantics.”

The following implementation defined pragmas are exceptions to this rule:
<table>
<thead>
<tr>
<th>Pragma</th>
<th>Explanation</th>
</tr>
</thead>
<tbody>
<tr>
<td>‘Abort_Defer’</td>
<td>Affects semantics</td>
</tr>
<tr>
<td>‘Ada_83’</td>
<td>Affects legality</td>
</tr>
<tr>
<td>‘Assert’</td>
<td>Affects semantics</td>
</tr>
<tr>
<td>‘CPP_Class’</td>
<td>Affects semantics</td>
</tr>
<tr>
<td>‘CPP_Constructor’</td>
<td>Affects semantics</td>
</tr>
<tr>
<td>‘Debug’</td>
<td>Affects semantics</td>
</tr>
<tr>
<td>‘Interface_Name’</td>
<td>Affects semantics</td>
</tr>
<tr>
<td>‘Machine_Attribute’</td>
<td>Affects semantics</td>
</tr>
<tr>
<td>‘Unimplemented_Unit’</td>
<td>Affects legality</td>
</tr>
<tr>
<td>‘Unchecked_Union’</td>
<td>Affects semantics</td>
</tr>
</tbody>
</table>

In each of the above cases, it is essential to the purpose of the pragma that this advice not be followed. For details see [Implementation Defined Pragmas], page 4.

6.5 RM 2.8(17-19): Pragmas

“Normally, an implementation should not define pragmas that can make an illegal program legal, except as follows:

* A pragma used to complete a declaration, such as a pragma Import;
* A pragma used to configure the environment by adding, removing, or replacing library_items.”

See [RM 2.8(16); Pragmas], page 153.

6.6 RM 3.5.2(5): Alternative Character Sets

“If an implementation supports a mode with alternative interpretations for Character and Wide_Character, the set of graphic characters of Character should nevertheless remain a proper subset of the set of graphic characters of Wide_Character. Any character set ‘localizations’ should be reflected in the results of the subprograms defined in the language-defined package Characters.Handling (see A.3) available in such a mode. In a mode with an alternative interpretation of Character, the implementation should also support a corresponding change in what is a legal identifier_letter.”

Not all wide character modes follow this advice, in particular the JIS and IEC modes reflect standard usage in Japan, and in these encoding, the upper half of the Latin-1 set is not
part of the wide-character subset, since the most significant bit is used for wide character encoding. However, this only applies to the external forms. Internally there is no such restriction.

6.7 RM 3.5.4(28): Integer Types

“An implementation should support Long_Integer in addition to Integer if the target machine supports 32-bit (or longer) arithmetic. No other named integer subtypes are recommended for package Standard. Instead, appropriate named integer subtypes should be provided in the library package Interfaces (see B.2).”

Long_Integer is supported. Other standard integer types are supported so this advice is not fully followed. These types are supported for convenient interface to C, and so that all hardware types of the machine are easily available.

6.8 RM 3.5.4(29): Integer Types

“An implementation for a two’s complement machine should support modular types with a binary modulus up to System.Max_Int*2+2. An implementation should support a non-binary modules up to Integer’Last.”

Followed.

6.9 RM 3.5.5(8): Enumeration Values

“For the evaluation of a call on S’Pos for an enumeration subtype, if the value of the operand does not correspond to the internal code for any enumeration literal of its type (perhaps due to an un-initialized variable), then the implementation should raise Program_Error. This is particularly important for enumeration types with noncontiguous internal codes specified by an enumeration_representation_clause.”

Followed.

6.10 RM 3.5.7(17): Float Types

“An implementation should support Long_Float in addition to Float if the target machine supports 11 or more digits of precision. No other named floating point subtypes are recommended for package Standard. Instead, appropriate named floating point subtypes should be provided in the library package Interfaces (see B.2).”

Short_Float and Long_Long_Float are also provided. The former provides improved compatibility with other implementations supporting this type. The latter corresponds to the highest precision floating-point type supported by the hardware. On most machines, this will be the same as Long_Float, but on some machines, it will correspond to the IEEE extended form. The notable case is all x86 implementations, where Long_Long_Float corresponds to the 80-bit extended precision format supported in hardware on this processor. Note that the 128-bit format on SPARC is not supported, since this is a software rather than a hardware format.
6.11 RM 3.6.2(11): Multidimensional Arrays

“An implementation should normally represent multidimensional arrays in row-major order, consistent with the notation used for multidimensional array aggregates (see 4.3.3). However, if a pragma `Convention (Fortran, . . .)` applies to a multidimensional array type, then column-major order should be used instead (see B.5, ‘Interfacing with Fortran’).”

Followed.

6.12 RM 9.6(30-31): Duration’Small

“Whenever possible in an implementation, the value of `Duration’Small` should be no greater than 100 microseconds.”

Followed. (`Duration’Small = 10**(-9)).

“The time base for `delay_relative_statements` should be monotonic; it need not be the same time base as used for `Calendar.Clock`.”

Followed.

6.13 RM 10.2.1(12): Consistent Representation

“In an implementation, a type declared in a pre-elaborated package should have the same representation in every elaboration of a given version of the package, whether the elaborations occur in distinct executions of the same program, or in executions of distinct programs or partitions that include the given version.”

Followed, except in the case of tagged types. Tagged types involve implicit pointers to a local copy of a dispatch table, and these pointers have representations which thus depend on a particular elaboration of the package. It is not easy to see how it would be possible to follow this advice without severely impacting efficiency of execution.

6.14 RM 11.4.1(19): Exception Information

“`Exception_Message` by default and `Exception_Information` should produce information useful for debugging. `Exception_Message` should be short, about one line. `Exception_Information` can be long. `Exception_Message` should not include the `Exception_Name`. `Exception_Information` should include both the `Exception_Name` and the `Exception_Message`.”

Followed. For each exception that doesn’t have a specified `Exception_Message`, the compiler generates one containing the location of the raise statement. This location has the form ‘file_name:line’, where file_name is the short file name (without path information) and line is the line number in the file. Note that in the case of the Zero Cost Exception mechanism, these messages become redundant with the `Exception_Information` that contains a full backtrace of the calling sequence, so they are disabled. To disable explicitly the generation of the source location message, use the Pragma `Discard_Names`.
6.15 RM 11.5(28): Suppression of Checks

“The implementation should minimize the code executed for checks that have been suppressed.”

Followed.

6.16 RM 13.1 (21-24): Representation Clauses

“The recommended level of support for all representation items is qualified as follows:

An implementation need not support representation items containing nonstatic expressions, except that an implementation should support a representation item for a given entity if each nonstatic expression in the representation item is a name that statically denotes a constant declared before the entity.”

Followed. In fact, GNAT goes beyond the recommended level of support by allowing non-static expressions in some representation clauses even without the need to declare constants initialized with the values of such expressions. For example:

```
X : Integer;
Y : Float;
for Y'Address use X'Address;>>
```

"An implementation need not support a specification for the ‘‘Size’’ for a given composite subtype, nor the size or storage place for an object (including a component) of a given composite subtype, unless the constraints on the subtype and its composite subcomponents (if any) are all static constraints."

Followed. Size Clauses are not permitted on nonstatic components, as described above.

“An aliased component, or a component whose type is by-reference, should always be allocated at an addressable location.”

Followed.

6.17 RM 13.2(6-8): Packed Types

“If a type is packed, then the implementation should try to minimize storage allocated to objects of the type, possibly at the expense of speed of accessing components, subject to reasonable complexity in addressing calculations.

The recommended level of support pragma Pack is:

For a packed record type, the components should be packed as tightly as possible subject to the Sizes of the component subtypes, and subject to any ‘record_representation_clause’ that applies to the type; the implementation may, but need not, reorder components or cross aligned word boundaries to improve the packing. A component whose Size is greater than the word size may be allocated an integral number of words.”
Followed. Tight packing of arrays is supported for all component sizes up to 64-bits. If the array component size is 1 (that is to say, if the component is a boolean type or an enumeration type with two values) then values of the type are implicitly initialized to zero. This happens both for objects of the packed type, and for objects that have a subcomponent of the packed type.

6.18 RM 13.3(14-19): Address Clauses

“For an array $X$, $X'\text{Address}$ should point at the first component of the array, and not at the array bounds.”

Followed.

“The recommended level of support for the Address attribute is:

$X'\text{Address}$ should produce a useful result if $X$ is an object that is aliased or of a by-reference type, or is an entity whose Address has been specified.”

Followed. A valid address will be produced even if none of those conditions have been met. If necessary, the object is forced into memory to ensure the address is valid.

“An implementation should support Address clauses for imported subprograms.”

Followed.

“Objects (including subcomponents) that are aliased or of a by-reference type should be allocated on storage element boundaries.”

Followed.

“If the Address of an object is specified, or it is imported or exported, then the implementation should not perform optimizations based on assumptions of no aliases.”

Followed.

6.19 RM 13.3(29-35): Alignment Clauses

“The recommended level of support for the Alignment attribute for subtypes is:

An implementation should support specified Alignments that are factors and multiples of the number of storage elements per word, subject to the following:”

Followed.

“An implementation need not support specified Alignments for combinations of Sizes and Alignments that cannot be easily loaded and stored by available machine instructions.”

Followed.

“An implementation need not support specified Alignments that are greater than the maximum Alignment the implementation ever returns by default.”

Followed.
“The recommended level of support for the Alignment attribute for objects is:

Same as above, for subtypes, but in addition:”

Followed.

“For stand-alone library-level objects of statically constrained subtypes, the implementation should support all alignments supported by the target linker. For example, page alignment is likely to be supported for such objects, but not for subtypes.”

Followed.

6.20 RM 13.3(42-43): Size Clauses

“The recommended level of support for the Size attribute of objects is:

A Size clause should be supported for an object if the specified Size is at least as large as its subtype’s Size, and corresponds to a size in storage elements that is a multiple of the object’s Alignment (if the Alignment is nonzero).”

Followed.

6.21 RM 13.3(50-56): Size Clauses

“If the Size of a subtype is specified, and allows for efficient independent addressability (see 9.10) on the target architecture, then the Size of the following objects of the subtype should equal the Size of the subtype:

Aliased objects (including components).”

Followed.

“Size clause on a composite subtype should not affect the internal layout of components.”

Followed. But note that this can be overridden by use of the implementation pragma Implicit_Packing in the case of packed arrays.

“The recommended level of support for the Size attribute of subtypes is:

The Size (if not specified) of a static discrete or fixed point subtype should be the number of bits needed to represent each value belonging to the subtype using an unbiased representation, leaving space for a sign bit only if the subtype contains negative values. If such a subtype is a first subtype, then an implementation should support a specified Size for it that reflects this representation.”

Followed.

“For a subtype implemented with levels of indirection, the Size should include the size of the pointers, but not the size of what they point at.”

Followed.
6.22 RM 13.3(71-73): Component Size Clauses

“The recommended level of support for the Component_Size attribute is:
An implementation need not support specified Component_Sizes that are less than the Size of the component subtype.”

Followed.

“An implementation should support specified Component_Sizes that are factors and multiples of the word size. For such Component_Sizes, the array should contain no gaps between components. For other Component_Sizes (if supported), the array should contain no gaps between components when packing is also specified; the implementation should forbid this combination in cases where it cannot support a no-gaps representation.”

Followed.

6.23 RM 13.4(9-10): Enumeration Representation Clauses

“The recommended level of support for enumeration representation clauses is:
An implementation need not support enumeration representation clauses for boolean types, but should at minimum support the internal codes in the range System.Min_Int .. System.Max_Int.”

Followed.

6.24 RM 13.5.1(17-22): Record Representation Clauses

“The recommended level of support for ‘record_representation_clause’s is:
An implementation should support storage places that can be extracted with a load, mask, shift sequence of machine code, and set with a load, shift, mask, store sequence, given the available machine instructions and run-time model.”

Followed.

“A storage place should be supported if its size is equal to the Size of the component subtype, and it starts and ends on a boundary that obeys the Alignment of the component subtype.”

Followed.

“If the default bit ordering applies to the declaration of a given type, then for a component whose subtype’s Size is less than the word size, any storage place that does not cross an aligned word boundary should be supported.”

Followed.

“An implementation may reserve a storage place for the tag field of a tagged type, and disallow other components from overlapping that place.”

Followed. The storage place for the tag field is the beginning of the tagged record, and its size is Address’Size. GNAT will reject an explicit component clause for the tag field.
“An implementation need not support a ‘component_clause’ for a component of an extension part if the storage place is not after the storage places of all components of the parent type, whether or not those storage places had been specified.”

Followed. The above advice on record representation clauses is followed, and all mentioned features are implemented.

### 6.25 RM 13.5.2(5): Storage Place Attributes

“If a component is represented using some form of pointer (such as an offset) to the actual data of the component, and this data is contiguous with the rest of the object, then the storage place attributes should reflect the place of the actual data, not the pointer. If a component is allocated discontinuously from the rest of the object, then a warning should be generated upon reference to one of its storage place attributes.”

Followed. There are no such components in GNAT.

### 6.26 RM 13.5.3(7-8): Bit Ordering

“The recommended level of support for the non-default bit ordering is: If Word_Size = Storage_Unit, then the implementation should support the non-default bit ordering in addition to the default bit ordering.”

Followed. Word size does not equal storage size in this implementation. Thus non-default bit ordering is not supported.

### 6.27 RM 13.7(37): Address as Private

“Address should be of a private type.”

Followed.

### 6.28 RM 13.7.1(16): Address Operations

“Operations in System and its children should reflect the target environment semantics as closely as is reasonable. For example, on most machines, it makes sense for address arithmetic to ‘wrap around’. Operations that do not make sense should raise Program_Error.”

Followed. Address arithmetic is modular arithmetic that wraps around. No operation raises Program_Error, since all operations make sense.

### 6.29 RM 13.9(14-17): Unchecked Conversion

“The Size of an array object should not include its bounds; hence, the bounds should not be part of the converted data.”

Followed.

“The implementation should not generate unnecessary run-time checks to ensure that the representation of S is a representation of the target type. It should take advantage of the permission to return by reference
when possible. Restrictions on unchecked conversions should be avoided unless required by the target environment.”

Followed. There are no restrictions on unchecked conversion. A warning is generated if the source and target types do not have the same size since the semantics in this case may be target dependent.

“The recommended level of support for unchecked conversions is:
Unchecked conversions should be supported and should be reversible in the cases where this clause defines the result. To enable meaningful use of unchecked conversion, a contiguous representation should be used for elementary subtypes, for statically constrained array subtypes whose component subtype is one of the subtypes described in this paragraph, and for record subtypes without discriminants whose component subtypes are described in this paragraph.”

Followed.

6.30 RM 13.11(23-25): Implicit Heap Usage

“An implementation should document any cases in which it dynamically allocates heap storage for a purpose other than the evaluation of an allocator.”

Followed, the only other points at which heap storage is dynamically allocated are as follows:

* At initial elaboration time, to allocate dynamically sized global objects.
* To allocate space for a task when a task is created.
* To extend the secondary stack dynamically when needed. The secondary stack is used for returning variable length results.

“A default (implementation-provided) storage pool for an access-to-constant type should not have overhead to support deallocation of individual objects.”

Followed.

“A storage pool for an anonymous access type should be created at the point of an allocator for the type, and be reclaimed when the designated object becomes inaccessible.”

Followed.

6.31 RM 13.11.2(17): Unchecked Deallocation

“For a standard storage pool, Free should actually reclaim the storage.”

Followed.

6.32 RM 13.13.2(1.6): Stream Oriented Attributes

“If not specified, the value of Stream_Size for an elementary type should be the number of bits that corresponds to the minimum number of stream elements required by the first subtype of the type, rounded up to the
nearest factor or multiple of the word size that is also a multiple of the stream element size.”

Followed, except that the number of stream elements is 1, 2, 3, 4 or 8. The Stream_Size may be used to override the default choice.

The default implementation is based on direct binary representations and is therefore target- and endianness-dependent. To address this issue, GNAT also supplies an alternate implementation of the stream attributes Read and Write, which uses the target-independent XDR standard representation for scalar types. This XDR alternative can be enabled via the binder switch -xdr.

6.33 RM A.1(52): Names of Predefined Numeric Types

“If an implementation provides additional named predefined integer types, then the names should end with Integer as in Long_Integer. If an implementation provides additional named predefined floating point types, then the names should end with Float as in Long_Float.”

Followed.

6.34 RM A.3.2(49): Ada.Characters.Handling

“If an implementation provides a localized definition of Character or Wide_Character, then the effects of the subprograms in Characters.Handling should reflect the localizations. See also 3.5.2.”

Followed. GNAT provides no such localized definitions.


“Bounded string objects should not be implemented by implicit pointers and dynamic allocation.”

Followed. No implicit pointers or dynamic allocation are used.

6.36 RM A.5.2(46-47): Random Number Generation

“Any storage associated with an object of type Generator should be reclaimed on exit from the scope of the object.”

Followed.

“If the generator period is sufficiently long in relation to the number of distinct initiator values, then each possible value of Initiator passed to Reset should initiate a sequence of random numbers that does not, in a practical sense, overlap the sequence initiated by any other value. If this is not possible, then the mapping between initiator values and generator states should be a rapidly varying function of the initiator value.”

Followed. The generator period is sufficiently long for the first condition here to hold true.
6.37 RM A.10.7(23): Get_Immediate

“The Get_Immediate procedures should be implemented with unbuffered input. For a device such as a keyboard, input should be available if a key has already been typed, whereas for a disk file, input should always be available except at end of file. For a file associated with a keyboard-like device, any line-editing features of the underlying operating system should be disabled during the execution of Get_Immediate.”

Followed on all targets except VxWorks. For VxWorks, there is no way to provide this functionality that does not result in the input buffer being flushed before the Get_Immediate call. A special unit Interfaces.Vxworks.IO is provided that contains routines to enable this functionality.

6.38 RM A.18: Containers

All implementation advice pertaining to Ada.Containers and its child units (that is, all implementation advice occurring within section A.18 and its subsections) is followed except for A.18.24(17):

“Bounded ordered set objects should be implemented without implicit pointers or dynamic allocation. “

The implementations of the two Reference_Preserving_Key functions of the generic package Ada.Containers.Bounded_Ordered_Sets each currently make use of dynamic allocation; other operations on bounded ordered set objects follow the implementation advice.

6.39 RM B.1(39-41): Pragma Export

“If an implementation supports pragma Export to a given language, then it should also allow the main subprogram to be written in that language. It should support some mechanism for invoking the elaboration of the Ada library units included in the system, and for invoking the finalization of the environment task. On typical systems, the recommended mechanism is to provide two subprograms whose link names are adainit and adafinal. adainit should contain the elaboration code for library units. adafinal should contain the finalization code. These subprograms should have no effect the second and subsequent time they are called.”

Followed.

“Automatic elaboration of pre-elaborated packages should be provided when pragma Export is supported.”

Followed when the main program is in Ada. If the main program is in a foreign language, then adainit must be called to elaborate pre-elaborated packages.

“For each supported convention ‘L’ other than Intrinsic, an implementation should support Import and Export pragmas for objects of ‘L’-compatible types and for subprograms, and pragma Convention for ‘L’-eligible types and for subprograms, presuming the other language has corresponding features. Pragma Convention need not be supported for scalar types.”

Followed.
**6.40 RM B.2(12-13): Package Interfaces**

“For each implementation-defined convention identifier, there should be a child package of package Interfaces with the corresponding name. This package should contain any declarations that would be useful for interfacing to the language (implementation) represented by the convention. Any declarations useful for interfacing to any language on the given hardware architecture should be provided directly in Interfaces.”

Followed.

“An implementation supporting an interface to C, COBOL, or Fortran should provide the corresponding package or packages described in the following clauses.”

Followed. GNAT provides all the packages described in this section.

**6.41 RM B.3(63-71): Interfacing with C**

“An implementation should support the following interface correspondences between Ada and C.”

Followed.

“An Ada procedure corresponds to a void-returning C function.”

Followed.

“An Ada function corresponds to a non-void C function.”

Followed.

“An Ada in scalar parameter is passed as a scalar argument to a C function.”

Followed.

“An Ada in parameter of an access-to-object type with designated type T is passed as a t* argument to a C function, where t is the C type corresponding to the Ada type T.”

Followed.

“An Ada access T parameter, or an Ada out or in out parameter of an elementary type T, is passed as a t* argument to a C function, where t is the C type corresponding to the Ada type T. In the case of an elementary out or in out parameter, a pointer to a temporary copy is used to preserve by-copy semantics.”

Followed.

“An Ada parameter of a record type T, of any mode, is passed as a t* argument to a C function, where t is the C structure corresponding to the Ada type T.”

Followed. This convention may be overridden by the use of the C_Pass_By_Copy pragma, or Convention, or by explicitly specifying the mechanism for a given call using an extended import or export pragma.

“An Ada parameter of an array type with component type T, of any mode, is passed as a t* argument to a C function, where t is the C type corresponding to the Ada type T.”
Followed.

“An Ada parameter of an access-to-subprogram type is passed as a pointer to a C function whose prototype corresponds to the designated subprogram’s specification.”

Followed.

**6.42 RM B.4(95-98): Interfacing with COBOL**

“An Ada implementation should support the following interface correspondences between Ada and COBOL.”

Followed.

“An Ada access T parameter is passed as a BY REFERENCE data item of the COBOL type corresponding to T.”

Followed.

“An Ada in scalar parameter is passed as a BY CONTENT data item of the corresponding COBOL type.”

Followed.

“Any other Ada parameter is passed as a BY REFERENCE data item of the COBOL type corresponding to the Ada parameter type; for scalars, a local copy is used if necessary to ensure by-copy semantics.”

Followed.

**6.43 RM B.5(22-26): Interfacing with Fortran**

“An Ada implementation should support the following interface correspondences between Ada and Fortran:”

Followed.

“An Ada procedure corresponds to a Fortran subroutine.”

Followed.

“An Ada function corresponds to a Fortran function.”

Followed.

“An Ada parameter of an elementary, array, or record type T is passed as a T argument to a Fortran procedure, where T is the Fortran type corresponding to the Ada type T, and where the INTENT attribute of the corresponding dummy argument matches the Ada formal parameter mode; the Fortran implementation’s parameter passing conventions are used. For elementary types, a local copy is used if necessary to ensure by-copy semantics.”

Followed.

“An Ada parameter of an access-to-subprogram type is passed as a reference to a Fortran procedure whose interface corresponds to the designated subprogram’s specification.”

Followed.

“The machine code or intrinsic support should allow access to all operations normally available to assembly language programmers for the target environment, including privileged instructions, if any.”

Followed.

“The interfacing pragmas (see Annex B) should support interface to assembler; the default assembler should be associated with the convention identifier Assembler.”

Followed.

“If an entity is exported to assembly language, then the implementation should allocate it at an addressable location, and should ensure that it is retained by the linking process, even if not otherwise referenced from the Ada code. The implementation should assume that any call to a machine code or assembler subprogram is allowed to read or update every object that is specified as exported.”

Followed.


“The implementation should ensure that little or no overhead is associated with calling intrinsic and machine-code subprograms.”

Followed for both intrinsics and machine-code subprograms.

“It is recommended that intrinsic subprograms be provided for convenient access to any machine operations that provide special capabilities or efficiency and that are not otherwise available through the language constructs.”

Followed. A full set of machine operation intrinsic subprograms is provided.

“Atomic read-modify-write operations—e.g., test and set, compare and swap, decrement and test, enqueue/dequeue.”

Followed on any target supporting such operations.

“Standard numeric functions—e.g., sin, log.”

Followed on any target supporting such operations.

“String manipulation operations—e.g., translate and test.”

Followed on any target supporting such operations.

“Vector operations—e.g., compare vector against thresholds.”

Followed on any target supporting such operations.

“Direct operations on I/O ports.”

Followed on any target supporting such operations.
6.46 RM C.3(28): Interrupt Support

“If the Ceiling_Locking policy is not in effect, the implementation should provide means for the application to specify which interrupts are to be blocked during protected actions, if the underlying system allows for a finer-grain control of interrupt blocking.”

Followed. The underlying system does not allow for finer-grain control of interrupt blocking.

6.47 RM C.3.1(20-21): Protected Procedure Handlers

“Whenever possible, the implementation should allow interrupt handlers to be called directly by the hardware.”

Followed on any target where the underlying operating system permits such direct calls.

“Whenever practical, violations of any implementation-defined restrictions should be detected before run time.”

Followed. Compile time warnings are given when possible.

6.48 RM C.3.2(25): Package Interrupts

“If implementation-defined forms of interrupt handler procedures are supported, such as protected procedures with parameters, then for each such form of a handler, a type analogous to Parameterless_Handler should be specified in a child package of Interrupts, with the same operations as in the predefined package Interrupts.”

Followed.

6.49 RM C.4(14): Pre-elaboration Requirements

“It is recommended that pre-elaborated packages be implemented in such a way that there should be little or no code executed at run time for the elaboration of entities not already covered by the Implementation Requirements.”

Followed. Executable code is generated in some cases, e.g., loops to initialize large arrays.

6.50 RM C.5(8): Pragma Discard_Names

“If the pragma applies to an entity, then the implementation should reduce the amount of storage used for storing names associated with that entity.”

Followed.

6.51 RM C.7.2(30): The Package Task_Attributes

“Some implementations are targeted to domains in which memory use at run time must be completely deterministic. For such implementations, it is recommended that the storage for task attributes will be pre-allocated statically and not from the heap. This can be accomplished by either placing restrictions on the number and the size of the task’s attributes,
or by using the pre-allocated storage for the first \( N \) attribute objects, and
the heap for the others. In the latter case, \( N \) should be documented."
Not followed. This implementation is not targeted to such a domain.

6.52 RM D.3(17): Locking Policies

“The implementation should use names that end with \_Locking for locking policies defined by the implementation.”
Followed. Two implementation-defined locking policies are defined, whose names (\texttt{Inheritance\_Locking} and \texttt{Concurrent\_Readers\_Locking}) follow this suggestion.

6.53 RM D.4(16): Entry Queuing Policies

“Names that end with \_Queuing should be used for all implementation-defined queuing policies.”
Followed. No such implementation-defined queuing policies exist.

6.54 RM D.6(9-10): Preemptive Abort

“Even though the ‘abort\_statement’ is included in the list of potentially blocking operations (see 9.5.1), it is recommended that this statement be implemented in a way that never requires the task executing the ‘abort\_statement’ to block.”
Followed.

“On a multi-processor, the delay associated with aborting a task on another processor should be bounded; the implementation should use periodic polling, if necessary, to achieve this.”
Followed.

6.55 RM D.7(21): Tasking Restrictions

“When feasible, the implementation should take advantage of the specified restrictions to produce a more efficient implementation.”
GNAT currently takes advantage of these restrictions by providing an optimized run time when the Ravenscar profile and the GNAT restricted run time set of restrictions are specified. See pragma \texttt{Profile (Ravenscar)} and pragma \texttt{Profile (Restricted)} for more details.

6.56 RM D.8(47-49): Monotonic Time

“When appropriate, implementations should provide configuration mechanisms to change the value of \texttt{Tick}.”
Such configuration mechanisms are not appropriate to this implementation and are thus not supported.

“It is recommended that \texttt{Calendar\_Clock} and \texttt{Real\_Time\_Clock} be implemented as transformations of the same time base.”
Followed.
“It is recommended that the best time base which exists in the underlying system be available to the application through Clock. *Best* may mean highest accuracy or largest range.”

Followed.

**6.57 RM E.5(28-29): Partition Communication Subsystem**

“Whenever possible, the PCS on the called partition should allow for multiple tasks to call the RPC-receiver with different messages and should allow them to block until the corresponding subprogram body returns.”

Followed by GLADE, a separately supplied PCS that can be used with GNAT.

“The *Write* operation on a stream of type *Params_Stream_Type* should raise *Storage_Error* if it runs out of space trying to write the *Item* into the stream.”

Followed by GLADE, a separately supplied PCS that can be used with GNAT.

**6.58 RM F(7): COBOL Support**

“If COBOL (respectively, C) is widely supported in the target environment, implementations supporting the Information Systems Annex should provide the child package *Interfaces.COBOL* (respectively, *Interfaces.C*) specified in Annex B and should support a *convention_identifier* of COBOL (respectively, C) in the interfacing pragmas (see Annex B), thus allowing Ada programs to interface with programs written in that language.”

Followed.

**6.59 RM F.1(2): Decimal Radix Support**

“Packed decimal should be used as the internal representation for objects of subtype *S* when *S’Machine_Radix = 10*.”

Not followed. GNAT ignores *S’Machine_Radix* and always uses binary representations.

**6.60 RM G: Numerics**

“If Fortran (respectively, C) is widely supported in the target environment, implementations supporting the Numerics Annex should provide the child package *Interfaces.Fortran* (respectively, *Interfaces.C*) specified in Annex B and should support a *convention_identifier* of Fortran (respectively, C) in the interfacing pragmas (see Annex B), thus allowing Ada programs to interface with programs written in that language.”

Followed.
6.61 RM G.1.1(56-58): Complex Types

“Because the usual mathematical meaning of multiplication of a complex operand and a real operand is that of the scaling of both components of the former by the latter, an implementation should not perform this operation by first promoting the real operand to complex type and then performing a full complex multiplication. In systems that, in the future, support an Ada binding to IEC 559:1989, the latter technique will not generate the required result when one of the components of the complex operand is infinite. (Explicit multiplication of the infinite component by the zero component obtained during promotion yields a NaN that propagates into the final result.) Analogous advice applies in the case of multiplication of a complex operand and a pure-imaginary operand, and in the case of division of a complex operand by a real or pure-imaginary operand.”

Not followed.

“Similarly, because the usual mathematical meaning of addition of a complex operand and a real operand is that the imaginary operand remains unchanged, an implementation should not perform this operation by first promoting the real operand to complex type and then performing a full complex addition. In implementations in which the Signed_Zeros attribute of the component type is True (and which therefore conform to IEC 559:1989 in regard to the handling of the sign of zero in predefined arithmetic operations), the latter technique will not generate the required result when the imaginary component of the complex operand is a negatively signed zero. (Explicit addition of the negative zero to the zero obtained during promotion yields a positive zero.) Analogous advice applies in the case of addition of a complex operand and a pure-imaginary operand, and in the case of subtraction of a complex operand and a real or pure-imaginary operand.”

Not followed.

“Implementations in which Real’Signed_Zeros is True should attempt to provide a rational treatment of the signs of zero results and result components. As one example, the result of the Argument function should have the sign of the imaginary component of the parameter X when the point represented by that parameter lies on the positive real axis; as another, the sign of the imaginary component of the Compose_From_Polar function should be the same as (respectively, the opposite of) that of the Argument parameter when that parameter has a value of zero and the Modulus parameter has a nonnegative (respectively, negative) value.”

Followed.

6.62 RM G.1.2(49): Complex Elementary Functions

“Implementations in which Complex_Types.Real’Signed_Zeros is True should attempt to provide a rational treatment of the signs of zero results and result components. For example, many of the complex elementary
functions have components that are odd functions of one of the parameter components; in these cases, the result component should have the sign of the parameter component at the origin. Other complex elementary functions have zero components whose sign is opposite that of a parameter component at the origin, or is always positive or always negative.”

Followed.

6.63 RM G.2.4(19): Accuracy Requirements

“The versions of the forward trigonometric functions without a `Cycle` parameter should not be implemented by calling the corresponding version with a `Cycle` parameter of `2.0*Numerics.Pi`, since this will not provide the required accuracy in some portions of the domain. For the same reason, the version of `Log` without a `Base` parameter should not be implemented by calling the corresponding version with a `Base` parameter of `Numerics.e`.”

Followed.

6.64 RM G.2.6(15): Complex Arithmetic Accuracy

“The version of the `Compose_From_Polar` function without a `Cycle` parameter should not be implemented by calling the corresponding version with a `Cycle` parameter of `2.0*Numerics.Pi`, since this will not provide the required accuracy in some portions of the domain.”

Followed.

6.65 RM H.6(15/2): Pragma Partition_Elaboration_Policy

“If the partition elaboration policy is `Sequential` and the Environment task becomes permanently blocked during elaboration then the partition is deadlocked and it is recommended that the partition be immediately terminated.”

Not followed.
7 Implementation Defined Characteristics

In addition to the implementation dependent pragmas and attributes, and the implementation advice, there are a number of other Ada features that are potentially implementation dependent and are designated as implementation-defined. These are mentioned throughout the Ada Reference Manual, and are summarized in Annex M.

A requirement for conforming Ada compilers is that they provide documentation describing how the implementation deals with each of these issues. In this chapter you will find each point in Annex M listed, followed by a description of how GNAT handles the implementation dependence.

You can use this chapter as a guide to minimizing implementation dependent features in your programs if portability to other compilers and other operating systems is an important consideration. The numbers in each entry below correspond to the paragraph numbers in the Ada Reference Manual.

* "Whether or not each recommendation given in Implementation Advice is followed. See 1.1.2(37)."

See [Implementation Advice], page 152.

* "Capacity limitations of the implementation. See 1.1.3(3)."

The complexity of programs that can be processed is limited only by the total amount of available virtual memory, and disk space for the generated object files.

* "Variations from the standard that are impractical to avoid given the implementation’s execution environment. See 1.1.3(6)."

There are no variations from the standard.

* "Which code statements cause external interactions. See 1.1.3(10)."

Any ‘code statements’ can potentially cause external interactions.

* "The coded representation for the text of an Ada program. See 2.1(4)."

See separate section on source representation.

* "The semantics of an Ada program whose text is not in Normalization Form C. See 2.1(4)."

See separate section on source representation.

* "The representation for an end of line. See 2.2(2)."

See separate section on source representation.

* "Maximum supported line length and lexical element length. See 2.2(15)."

The maximum line length is 255 characters and the maximum length of a lexical element is also 255 characters. This is the default setting if not overridden by the use of compiler switch ‘-gnaty’ (which sets the maximum to 79) or ‘-gnatyMnn’ which allows the maximum line length to be specified to be any value up to 32767. The maximum length of a lexical element is the same as the maximum line length.

* "Implementation defined pragmas. See 2.8(14)."

See [Implementation Defined Pragmas], page 4.
Pragma `Optimize`, if given with a `Time` or `Space` parameter, checks that the optimization flag is set, and aborts if it is not.

In the case of a `Dynamic_Predicate` aspect, the string is “Dynamic_Predicate failed at <source position>”, where “<source position>” might be something like “foo.adb:123”. The `Static_Predicate` case is handled analogously.

* “The predefined integer types declared in `Standard`. See 3.5.4(25).”

<table>
<thead>
<tr>
<th>Type</th>
<th>Representation</th>
</tr>
</thead>
<tbody>
<tr>
<td>‘Short_Short_Integer’</td>
<td>8-bit signed</td>
</tr>
<tr>
<td>‘Short_Integer’</td>
<td>16-bit signed</td>
</tr>
<tr>
<td>‘Integer’</td>
<td>32-bit signed</td>
</tr>
<tr>
<td>‘Long_Integer’</td>
<td>64-bit signed (on most 64-bit targets, depending on the C definition of long) 32-bit signed (on all other targets)</td>
</tr>
<tr>
<td>‘Long_Long_Integer’</td>
<td>64-bit signed</td>
</tr>
<tr>
<td>‘Long_Long_Long_Integer’</td>
<td>128-bit signed (on 64-bit targets) 64-bit signed (on 32-bit targets)</td>
</tr>
</tbody>
</table>

* “Any nonstandard integer types and the operators defined for them. See 3.5.4(26).”

There are no nonstandard integer types.

* “Any nonstandard real types and the operators defined for them. See 3.5.6(8).”

There are no nonstandard real types.

* “What combinations of requested decimal precision and range are supported for floating point types. See 3.5.7(7).”

The precision and range are defined by the IEEE Standard for Floating-Point Arithmetic (IEEE 754-2019).

* “The predefined floating point types declared in `Standard`. See 3.5.7(16).”

<table>
<thead>
<tr>
<th>Type</th>
<th>Representation</th>
</tr>
</thead>
<tbody>
<tr>
<td>‘Short_Float’</td>
<td>IEEE Binary32 (Single)</td>
</tr>
<tr>
<td>‘Float’</td>
<td>IEEE Binary32 (Single)</td>
</tr>
</tbody>
</table>
Chapter 7: Implementation Defined Characteristics

'Long_Float' IEEE Binary64 (Double)

'Long_Long_Float' IEEE Binary64 (Double) on non-x86 architectures IEEE 80-bit Extended on x86 architecture

The default rounding mode specified by the IEEE 754 Standard is assumed both for static and dynamic computations (that is, round to nearest, ties to even). The input routines yield correctly rounded values for Short_Float, Float, and Long_Float at least. The output routines can compute up to twice as many exact digits as the value of T'Digits for any type, for example 30 digits for Long_Float; if more digits are requested, zeros are printed.

* "The small of an ordinary fixed point type. See 3.5.9(8)."

The small is the largest power of two that does not exceed the delta.

* "What combinations of small, range, and digits are supported for fixed point types. See 3.5.9(10)."

For an ordinary fixed point type, on 32-bit platforms, the small must lie in 2.0**(−80) .. 2.0**80 and the range in -9.0E+36 .. 9.0E+36; any combination is permitted that does not result in a mantissa larger than 63 bits.

On 64-bit platforms, the small must lie in 2.0**(−127) .. 2.0**127 and the range in -1.0E+76 .. 1.0E+76; any combination is permitted that does not result in a mantissa larger than 63 bits, and any combination is permitted that results in a mantissa between 64 and 127 bits if the small is the ratio of two integers that lie in 1 .. 2.0**127.

If the small is the ratio of two integers with 64-bit magnitude on 32-bit platforms and 128-bit magnitude on 64-bit platforms, which is the case if no small clause is provided, then the operations of the fixed point type are entirely implemented by means of integer instructions. In the other cases, some operations, in particular input and output, may be implemented by means of floating-point instructions and may be affected by accuracy issues on architectures other than x86.

For a decimal fixed point type, on 32-bit platforms, the small must lie in 1.0E−18 .. 1.0E+18 and the digits in 1 .. 18. On 64-bit platforms, the small must lie in 1.0E−38 .. 1.0E+38 and the digits in 1 .. 38.

* "The result of Tags.Expanded_Name for types declared within an unnamed 'block_statement'. See 3.9(10)."

Block numbers of the form Bnnn, where ‘nnn’ is a decimal integer are allocated.

* "The sequence of characters of the value returned by Tags.Expanded_Name (respectively, Tags.Wide_Expanded_Name) when some of the graphic characters of Tags.Wide_Wide_Expanded_Name are not defined in Character (respectively, Wide_Character). See 3.9(10.1)."

This is handled in the same way as the implementation-defined behavior referenced in A.4.12(34).

* "Implementation-defined attributes. See 4.1.4(12)."

See [Implementation Defined Attributes], page 115.

* "The value of the parameter to Empty for some container aggregates. See 4.3.5(40)."
As per the suggestion given in the Annotated Ada RM, the default value of the formal parameter is used if one exists and zero is used otherwise.

* “The maximum number of chunks for a parallel reduction expression without a chunk_specification. See 4.5.10(21).”

Feature unimplemented.

* “Round to even is used in all such cases.

* “The maximum number of chunks for a parallel generalized iterator without a chunk_specification. See 5.5.2(10).”

Feature unimplemented.

* “The number of chunks for an array component iterator. See 5.5.2(11).”

Feature unimplemented.

* “Any extensions of the Global aspect. See 6.1.2(43).”

Feature unimplemented.

* “The circumstances the implementation passes in the null value for a view conversion of an access type used as an out parameter. See 6.4.1(19).”

Difficult to characterize.

* “Any extensions of the Default_Initial.Condition aspect. See 7.3.3(11).”

SPARK allows specifying ‘null’ as the Default_Initial.Condition aspect of a type. See the SPARK reference manual for further details.

* “Any implementation-defined time types. See 9.6(6).”

There are no implementation-defined time types.

* “The time base associated with relative delays. See 9.6(20).”

See 9.6(20). The time base used is that provided by the C library function gettimeofday.

* “The time base of the type Calendar.Time. See 9.6(23).”

The time base used is that provided by the C library function gettimeofday.

* “The time zone used for package Calendar operations. See 9.6(24).”

The time zone used by package Calendar is the current system time zone setting for local time, as accessed by the C library function localtime.

* “Any limit on ‘delay_until_statements’ of ‘select_statements’. See 9.6(29).”

There are no such limits.

* “The result of Calendar.Formatting.Image if its argument represents more than 100 hours. See 9.6.1(86).”

Calendar.Time_Error is raised.

* “Implementation-defined conflict check policies. See 9.10.1(5).”

There are no implementation-defined conflict check policies.
Chapter 7: Implementation Defined Characteristics

* “The representation for a compilation. See 10.1(2).”
A compilation is represented by a sequence of files presented to the compiler in a single invocation of the ‘gcc’ command.

* “Any restrictions on compilations that contain multiple compilation units. See 10.1(4).”
No single file can contain more than one compilation unit, but any sequence of files can be presented to the compiler as a single compilation.

* “The mechanisms for creating an environment and for adding and replacing compilation units. See 10.1.4(3).”
See separate section on compilation model.

* “The manner of explicitly assigning library units to a partition. See 10.2(2).”
If a unit contains an Ada main program, then the Ada units for the partition are determined by recursive application of the rules in the Ada Reference Manual section 10.2(2-6). In other words, the Ada units will be those that are needed by the main program, and then this definition of need is applied recursively to those units, and the partition contains the transitive closure determined by this relationship. In short, all the necessary units are included, with no need to explicitly specify the list. If additional units are required, e.g., by foreign language units, then all units must be mentioned in the context clause of one of the needed Ada units.

If the partition contains no main program, or if the main program is in a language other than Ada, then GNAT provides the binder options ‘-z’ and ‘-n’ respectively, and in this case a list of units can be explicitly supplied to the binder for inclusion in the partition (all units needed by these units will also be included automatically). For full details on the use of these options, refer to ‘GNAT Make Program gnatmake’ in the GNAT User’s Guide.

* “The implementation-defined means, if any, of specifying which compilation units are needed by a given compilation unit. See 10.2(2).”
The units needed by a given compilation unit are as defined in the Ada Reference Manual section 10.2(2-6). There are no implementation-defined pragmas or other implementation-defined means for specifying needed units.

* “The manner of designating the main subprogram of a partition. See 10.2(7).”
The main program is designated by providing the name of the corresponding ALI file as the input parameter to the binder.

* “The order of elaboration of ‘library_items’. See 10.2(18).”
The first constraint on ordering is that it meets the requirements of Chapter 10 of the Ada Reference Manual. This still leaves some implementation-dependent choices, which are resolved by analyzing the elaboration code of each unit and identifying implicit elaboration-order dependencies.

* “Parameter passing and function return for the main subprogram. See 10.2(21).”
The main program has no parameters. It may be a procedure, or a function returning an integer type. In the latter case, the returned integer value is the return code of the program (overriding any value that may have been set by a call to Ada.Command_Line.Set_Exit_Status).
* “The mechanisms for building and running partitions. See 10.2(24).”

GNAT itself supports programs with only a single partition. The GNATDIST tool provided with the GLADE package (which also includes an implementation of the PCS) provides a completely flexible method for building and running programs consisting of multiple partitions. See the separate GLADE manual for details.

* “The details of program execution, including program termination. See 10.2(25).”

See separate section on compilation model.

* “The semantics of any non-active partitions supported by the implementation. See 10.2(28).”

Passive partitions are supported on targets where shared memory is provided by the operating system. See the GLADE reference manual for further details.

* “The information returned by Exception_Message. See 11.4.1(10).”

Exception message returns the null string unless a specific message has been passed by the program.

* “The result of Exceptions.Exception_Name for types declared within an unnamed ‘block_statement’. See 11.4.1(12).”

Blocks have implementation defined names of the form Bnnn where ‘nnn’ is an integer.

* “The information returned by Exception_Information. See 11.4.1(13).”

Exception_Information returns a string in the following format:

```plaintext
*Exception_Name:* nnnnn
*Message:* mmmmm
*PID:* ppp
*Load address:* 0xhhhh
*Call stack traceback locations:* 0xhhhh 0xhhhh 0xhhhh ... 0xhhh
```

where

* nnnn is the fully qualified name of the exception in all upper case letters. This line is always present.
* mmmmm is the message (this line present only if message is non-null)
* ppp is the Process Id value as a decimal integer (this line is present only if the Process Id is nonzero). Currently we are not making use of this field.
* The Load address line, the Call stack traceback locations line and the following values are present only if at least one traceback location was recorded. The Load address indicates the address at which the main executable was loaded; this line may not be present if operating system hasn’t relocated the main executable. The values are given in C style format, with lower case letters for a-f, and only as many digits present as are necessary. The line terminator sequence at the end of each line, including the last line is a single LF character (16#0A#).
* “The sequence of characters of the value returned by Exceptions.Exception_Name (respectively, Exceptions.Wide_Exception_Name) when some of the graphic characters of Exceptions.Wide_Exception_Name are not defined in Character (respectively, Wide_Character). See 11.4.1(12.1).”

This is handled in the same way as the implementation-defined behavior referenced in A.4.12(34).

* “The information returned by Exception_Information. See 11.4.1(13).”

The exception name and the source location at which the exception was raised are included.

* “Implementation-defined policy_identifiers and assertion_aspect_marks allowed in a pragma Assertion_Policy. See 11.4.2(9).”

Implementation-defined assertion_aspect_marks include Assert_And_Cut, Assume, Contract_Cases, Debug, Ghost, Initial_Condition, Loop_Invariant, Loop_Variant, Postcondition, Precondition, Predicate, Refined_Post, Statement Assertions, and Subprogram_Variant. Implementation-defined policy_identifiers include Ignore and Suppressible.

* “The default assertion policy. See 11.4.2(10).”

The default assertion policy is Ignore, although this can be overridden via compiler switches such as “-gnata”.

* “Implementation-defined check names. See 11.5(27).”

The implementation defined check names include Alignment_Check, Atomic_Synchronization, Duplicated_Tag_Check, Container_Checks, Tampering_Check, Predicate_Check, and Validity_Check. In addition, a user program can add implementation-defined check names by means of the pragma Check_Name. See the description of pragma Suppress for full details.

* “Existence and meaning of second parameter of pragma Unsuppress. See 11.5(27.1).”

The legality rules for and semantics of the second parameter of pragma Unsuppress match those for the second argument of pragma Suppress.

* “The cases that cause conflicts between the representation of the ancestors of a type_declaration. See 13.1(13.1).”

No such cases exist.

* “The interpretation of each representation aspect. See 13.1(20).”

See separate section on data representations.

* “Any restrictions placed upon the specification of representation aspects. See 13.1(20).”

See separate section on data representations.

* “Implementation-defined aspects, including the syntax for specifying such aspects and the legality rules for such aspects. See 13.1.1(38).”

See [Implementation Defined Aspects], page 107.

* “The set of machine scalars. See 13.3(8.1).”

See separate section on data representations.

* “The meaning of Size for indefinite subtypes. See 13.3(48).”
The Size attribute of an indefinite subtype is not less than the Size attribute of any object of that type.

* "The meaning of Object_Size for indefinite subtypes. See 13.3(58)."

The Object_Size attribute of an indefinite subtype is not less than the Object_Size attribute of any object of that type.

* "The default external representation for a type tag. See 13.3(75)."

The default external representation for a type tag is the fully expanded name of the type in upper case letters.

* "What determines whether a compilation unit is the same in two different partitions. See 13.3(76)."

A compilation unit is the same in two different partitions if and only if it derives from the same source file.

* "Implementation-defined components. See 13.5.1(15)."

The only implementation defined component is the tag for a tagged type, which contains a pointer to the dispatching table.

* "If Word_Size = Storage_Unit, the default bit ordering. See 13.5.3(5)."

Word_Size (32) is not the same as Storage_Unit (8) for this implementation, so no non-default bit ordering is supported. The default bit ordering corresponds to the natural endianness of the target architecture.

* "The contents of the visible part of package System. See 13.7(2)."

See the definition of package System in system.ads. Note that two declarations are added to package System.

```
Max_Priority : constant Positive := Priority'Last;
Max_Interrupt_Priority : constant Positive := Interrupt_Priority'Last;
```

* "The range of Storage_Elements.Storage_Offset, the modulus of Storage_Elements.Storage_Element, and the declaration of Storage_Elements.Integer_Address. See 13.7.1(11)."

See the definition of package System.Storage_Elements in s-stoele.ads.

* "The contents of the visible part of package System.Machine_Code, and the meaning of 'code_statements'. See 13.8(7)."

See the definition and documentation in file s-maccod.ads.

* "The result of unchecked conversion for instances with scalar result types whose result is not defined by the language. See 13.9(11)."

Unchecked conversion between types of the same size results in an uninterpreted transmission of the bits from one type to the other. If the types are of unequal sizes, then in the case of discrete types, a shorter source is first zero or sign extended as necessary, and a shorter target is simply truncated on the left. For all non-discrete types, the source is first copied if necessary to ensure that the alignment requirements of the target are met, then a pointer is constructed to the source value, and the result is obtained by dereferencing this pointer after converting it to be a pointer to the target type. Unchecked conversions where the target subtype is an unconstrained array are not permitted. If the target alignment
is greater than the source alignment, then a copy of the result is made with appropriate alignment.
* "The result of unchecked conversion for instances with nonscalar result types whose result is not defined by the language. See 13.9(11)."

See preceding definition for the scalar result case.

* "Whether or not the implementation provides user-accessible names for the standard pool type(s). See 13.11(17)."

There are 3 different standard pools used by the compiler when Storage_Pool is not specified depending whether the type is local to a subprogram or defined at the library level and whether Storage_Size is specified or not. See documentation in the runtime library units  ‘System.Pool_Global, System.Pool_Size and System.Pool_Local’ in files s-poosiz.ads, s-pooglo.ads and s-pooloc.ads for full details on the default pools used. All these pools are accessible by means of withing these units.

* "The meaning of Storage_Size when neither the Storage_Size nor the Storage_Pool is specified for an access type. See 13.11(18)."

Storage_Size is measured in storage units, and refers to the total space available for an access type collection, or to the primary stack space for a task.

* "The effect of specifying aspect Default_Storage_Pool on an instance of a language-defined generic unit. See 13.11.3(5)."

Instances of language-defined generic units are treated the same as other instances with respect to the Default_Storage_Pool aspect.

* "Implementation-defined restrictions allowed in a pragma Restrictions. See 13.12(8.7)."

See [Standard and Implementation Defined Restrictions], page 138.

* "The consequences of violating limitations on Restrictions pragmas. See 13.12(9)."

Restrictions that can be checked at compile time are enforced at compile time; violations are illegal. For other restrictions, any violation during program execution results in erroneous execution.

* "Implementation-defined usage profiles allowed in a pragma Profile. See 13.12(15)."

See [Implementation Defined Pragmas], page 4.

* "The contents of the stream elements read and written by the Read and Write attributes of elementary types. See 13.13.2(9)."

The representation is the in-memory representation of the base type of the type, using the number of bits corresponding to the type’s Size value, and the natural ordering of the machine.

* "The names and characteristics of the numeric subtypes declared in the visible part of package Standard. See A.1(3)."

See items describing the integer and floating-point types supported.

* "The values returned by Strings.Hash. See A.4.9(3)."

This hash function has predictable collisions and is subject to equivalent substring attacks. It is not suitable for construction of a hash table keyed on possibly malicious user input.
* “The value returned by a call to a Text_Buffer Get procedure if any character in the returned sequence is not defined in Character. See A.4.12(34).”

The contents of a buffer is represented internally as a UTF_8 string. The value return by Text_Buffer.Get is the result of passing that UTF_8 string to UTF_Encoding.Strings.Decode.

* “The value returned by a call to a Text_Buffer Wide_Get procedure if any character in the returned sequence is not defined in Wide_Character. See A.4.12(34).”

The contents of a buffer is represented internally as a UTF_8 string. The value return by Text_Buffer.Wide_Get is the result of passing that UTF_8 string to UTF_Encoding.Wide_Strings.Decode.

* “The accuracy actually achieved by the elementary functions. See A.5.1(1).”

The elementary functions correspond to the functions available in the C library. Only fast math mode is implemented.

* “The sign of a zero result from some of the operators or functions in Numerics.Generic_Elementary_Functions, when Float_Type’Signed_Zeros is True. See A.5.1(46).”

The sign of zeroes follows the requirements of the IEEE 754 standard on floating-point.


Maximum image width is 6864, see library file s-rannum.ads.


Maximum image width is 6864, see library file s-rannum.ads.

* “The string representation of a random number generator’s state. See A.5.2(38).”

The value returned by the Image function is the concatenation of the fixed-width decimal representations of the 624 32-bit integers of the state vector.

* “The values of the Model_Mantissa, Model_Emin, Model_Epsilon, Model, Safe_First, and Safe_Last attributes, if the Numerics Annex is not supported. See A.5.3(72).”

Running the compiler with ‘-gnatS’ to produce a listing of package Standard displays the values of these attributes.

* “The value of Buffer_Size in Storage_IO. See A.9(10).”

All type representations are contiguous, and the Buffer_Size is the value of type’Size rounded up to the next storage unit boundary.

* “External files for standard input, standard output, and standard error See A.10(5).”

These files are mapped onto the files provided by the C streams libraries. See source file i-cstrea.ads for further details.

* “The accuracy of the value produced by Put. See A.10.9(36).”

If more digits are requested in the output than are represented by the precision of the value, zeroes are output in the corresponding least significant digit positions.

* “Current size for a stream file for which positioning is not supported. See A.12.1(1.1).”

Positioning is supported.
* “The meaning of Argument_Count, Argument, and Command_Name. See A.15(1).”

These are mapped onto the argv and argc parameters of the main program in the natural manner.

* “The interpretation of file names and directory names. See A.16(46).”

These names are interpreted consistently with the underlying file system.

* “The maximum value for a file size in Directories. See A.16(87).”

Directories.File_Size’Last is equal to Long_Long_Integer’Last.

* “The result for Directories.Size for a directory or special file. See A.16(93).”

Name_Error is raised.

* “The result for Directories.Modification_Time for a directory or special file. See A.16(93).”

Name_Error is raised.

* “The interpretation of a nonnull search pattern in Directories. See A.16(104).”

When the Pattern parameter is not the null string, it is interpreted according to the syntax of regular expressions as defined in the GNAT.Regexp package.

See [GNAT.Regexp (g-regexp.ads)], page 265.

* “The results of a Directories search if the contents of the directory are altered while a search is in progress. See A.16(110).”

The effect of a call to Get_Next_Entry is determined by the current state of the directory.

* “The definition and meaning of an environment variable. See A.17(1).”

This definition is determined by the underlying operating system.

* “The circumstances where an environment variable cannot be defined. See A.17(16).”

There are no such implementation-defined circumstances.

* “Environment names for which Set has the effect of Clear. See A.17(17).”

There are no such names.


Containers.Hash_Type’Modulus is 2**32. Containers.Count_Type’Last is 2**31 - 1.

* “Implementation-defined convention names. See B.1(11).”

The following convention names are supported

<table>
<thead>
<tr>
<th>Convention Name</th>
<th>Interpretation</th>
</tr>
</thead>
<tbody>
<tr>
<td>‘Ada’</td>
<td>Ada</td>
</tr>
<tr>
<td>‘Ada_Pass_By_Copy’</td>
<td>Allowed for any types except by-reference types such as limited records. Convention Ada, but causes any parameters with this convention to be passed by copy.</td>
</tr>
</tbody>
</table>
‘Ada_Pass_By_Reference’ Allowed for any types except by-copy types such as scalars. Compatible with Ada, but causes any parameters with this convention to be passed by reference.

‘Assembler’ Assembly language

‘Asm’ Synonym for Assembler

‘Assembly’ Synonym for Assembler

‘C’ C

‘C_Pass_By_Copy’ Allowed only for record types, like C, but also notes that record is to be passed by copy rather than reference.

‘COBOL’ COBOL

‘C_Plus_Plus (or CPP)’ C++

‘Default’ Treated the same as C

‘External’ Treated the same as C

‘Fortran’ Fortran

‘Intrinsic’ For support of pragma Import with convention Intrinsic, see separate section on Intrinsic Subprograms.

‘Stdcall’ Stdcall (used for Windows implementations only). This convention corresponds to the (previously called Pascal convention) C/C++ convention under Windows. A routine with this convention cleans the stack before exit. This pragma cannot be applied to a dispatching call.

‘DLL’ Synonym for Stdcall

‘Win32’ Synonym for Stdcall

‘Stubbed’ Stubbed is a special convention used to indicate that the body of the subprogram will be entirely ignored. Any call to the subprogram is converted into a raise of the Program_Error exception. If a pragma Import specifies convention stubbed then no body need be present at all. This convention is useful during development for the inclusion of subprograms whose body has not yet been written. In addition, all otherwise unrecognized convention names are also treated as being synonymous with convention C. In all implementations, use of such other names results in a warning.

* “The meaning of link names. See B.1(36).”

Link names are the actual names used by the linker.

* “The manner of choosing link names when neither the link name nor the address of an imported or exported entity is specified. See B.1(36).”
The default linker name is that which would be assigned by the relevant external language, interpreting the Ada name as being in all lower case letters.

* “The effect of pragma Linker_Options. See B.1(37).”

The string passed to Linker_Options is presented uninterpreted as an argument to the link command, unless it contains ASCII.NUL characters. NUL characters if they appear act as argument separators, so for example

```ada
pragma Linker_Options ('-labc' & ASCII.NUL & '-ldef');
```

causes two separate arguments `-labc` and `-ldef` to be passed to the linker. The order of linker options is preserved for a given unit. The final list of options passed to the linker is in reverse order of the elaboration order. For example, linker options for a body always appear before the options from the corresponding package spec.

* “The contents of the visible part of package Interfaces and its language-defined descendants. See B.2(1).”

See files with prefix `i-` in the distributed library.

* “Implementation-defined children of package Interfaces. The contents of the visible part of package Interfaces. See B.2(11).”

See files with prefix `i-` in the distributed library.

* “The definitions of certain types and constants in Interfaces.C. See B.3(41).”

See source file `i-c.ads`.

* “The types Floating, Long_Floating, Binary, Long_Binary, Decimal_Element, and COBOL_Character; and the initialization of the variables Ada_To_COBOL and COBOL_To_Ada, in Interfaces.COBOL. See B.4(50).”

<table>
<thead>
<tr>
<th>COBOL</th>
<th>Ada</th>
</tr>
</thead>
<tbody>
<tr>
<td>'Floating'</td>
<td>Float</td>
</tr>
<tr>
<td>'Long_Floating'</td>
<td>(Floating) Long_Float</td>
</tr>
<tr>
<td>'Binary'</td>
<td>Integer</td>
</tr>
<tr>
<td>'Long_Binary'</td>
<td>Long_Long_Integer</td>
</tr>
<tr>
<td>'Decimal_Element'</td>
<td>Character</td>
</tr>
<tr>
<td>'COBOL_Character'</td>
<td>Character</td>
</tr>
</tbody>
</table>

For initialization, see the file `i-cobol.ads` in the distributed library.

* “The types Fortran_Integer, Real, Double_Precision, and Character_Set in Interfaces.Fortran. See B.5(17).”

See source file `i-fortra.ads`. These types are derived, respectively, from Integer, Float, Long_Float, and Character.
* “Implementation-defined intrinsic subprograms. See C.1(1).”

See separate section on Intrinsic Subprograms.

* “Any restrictions on a protected procedure or its containing type when an aspect Attach_Handler or Interrupt_Handler is specified. See C.3.1(17).”

There are no such restrictions.

* “Any other forms of interrupt handler supported by the Attach_Handler and Interrupt_Handler aspects. See C.3.1(19).”

There are no such forms.

* “The semantics of some attributes and functions of an entity for which aspect Discard_Names is True. See C.5(7).”

If Discard_Names is True for an enumeration type, the Image attribute provides the image of the Pos of the literal, and Value accepts Pos values.

If both of the aspects “Discard_Names” and No_Tagged_Streams are true for a tagged type, its Expanded_Name and External_Tag values are empty strings. This is useful to avoid exposing entity names at binary level.

* “The modulus and size of Test_and_Set_Flag. See C.6.3(8).”

The modulus is $2^{**8}$. The size is 8.

* “The value used to represent the set value for Atomic_Test_and_Set. See C.6.3(10).”

The value is 1.

* “The result of the Task_Identification.Image attribute. See C.7.1(7).”

The result of this attribute is a string that identifies the object or component that denotes a given task. If a variable Var has a task type, the image for this task will have the form Var_XXXXXXXXX, where the suffix ‘XXXXXXXX’ is the hexadecimal representation of the virtual address of the corresponding task control block. If the variable is an array of tasks, the image of each task will have the form of an indexed component indicating the position of a given task in the array, e.g., Group(5)_XXXXXXXX. If the task is a component of a record, the image of the task will have the form of a selected component. These rules are fully recursive, so that the image of a task that is a subcomponent of a composite object corresponds to the expression that designates this task.

If a task is created by an allocator, its image depends on the context. If the allocator is part of an object declaration, the rules described above are used to construct its image, and this image is not affected by subsequent assignments. If the allocator appears within an expression, the image includes only the name of the task type.

If the configuration pragma Discard_Names is present, or if the restriction No_Implicit_Heap_Allocation is in effect, the image reduces to the numeric suffix, that is to say the hexadecimal representation of the virtual address of the control block of the task.

* “The value of Current_Task when in a protected entry or interrupt handler. See C.7.1(17).”

Protected entries or interrupt handlers can be executed by any convenient thread, so the value of Current_Task is undefined.

* “Granularity of locking for Task_Attributes. See C.7.2(16).”
No locking is needed if the formal type Attribute has the size and alignment of either Integer or System.Address and the bit representation of Initial_Value is all zeroes. Otherwise, locking is performed.

* “The declarations of Any_Priority and Priority. See D.1(11).”

See declarations in file system.ads.

* “Implementation-defined execution resources. See D.1(15).”

There are no implementation-defined execution resources.

* “Whether, on a multiprocessor, a task that is waiting for access to a protected object keeps its processor busy. See D.2.1(3).”

On a multi-processor, a task that is waiting for access to a protected object does not keep its processor busy.

* “The affect of implementation defined execution resources on task dispatching. See D.2.1(9).”

Tasks map to threads in the threads package used by GNAT. Where possible and appropriate, these threads correspond to native threads of the underlying operating system.

* “Implementation-defined task dispatching policies. See D.2.2(3).”

There are no implementation-defined task dispatching policies.

* “The value of Default_Quantum in Dispatching.Round_Robin. See D.2.5(4).”

The value is 10 milliseconds.

* “Implementation-defined ‘policy_identifiers’ allowed in a pragma Locking_Policy. See D.3(4).”

The two implementation defined policies permitted in GNAT are Inheritance_Locking and Concurrent_Readers_Locking. On targets that support the Inheritance_Locking policy, locking is implemented by inheritance, i.e., the task owning the lock operates at a priority equal to the highest priority of any task currently requesting the lock. On targets that support the Concurrent_Readers_Locking policy, locking is implemented with a read/write lock allowing multiple protected object functions to enter concurrently.

* “Default ceiling priorities. See D.3(10).”

The ceiling priority of protected objects of the type System.Interrupt_Priority’Last as described in the Ada Reference Manual D.3(10),

* “The ceiling of any protected object used internally by the implementation. See D.3(16).”

The ceiling priority of internal protected objects is System.Priority’Last.

* “Implementation-defined queuing policies. See D.4(1).”

There are no implementation-defined queuing policies.

* “Implementation-defined admission policies. See D.4.1(1).”

There are no implementation-defined admission policies.

* “Any operations that implicitly require heap storage allocation. See D.7(8).”

The only operation that implicitly requires heap storage allocation is task creation.
* “When restriction \texttt{No\_Dynamic\_CPU\_Assignment} applies to a partition, the processor on which a task with a CPU value of a \texttt{Not\_A\_Specific\_CPU} will execute. See \texttt{D.7(10)}.”

Unknown.
* “When restriction \texttt{No\_Task\_Termination} applies to a partition, what happens when a task terminates. See \texttt{D.7(15.1)}.”

Execution is erroneous in that case.
* “The behavior when restriction \texttt{Max\_Storage\_At\_Blocking} is violated. See \texttt{D.7(17)}.”

Execution is erroneous in that case.
* “The behavior when restriction \texttt{Max\_Asynchronous\_Select\_Nesting} is violated. See \texttt{D.7(18)}.”

Execution is erroneous in that case.
* “The behavior when restriction \texttt{Max\_Tasks} is violated. See \texttt{D.7(19)}.”

Execution is erroneous in that case.
* “Whether the use of pragma \texttt{Restrictions} results in a reduction in program code or data size or execution time. See \texttt{D.7(20)}.”

Yes it can, but the precise circumstances and properties of such reductions are difficult to characterize.
* “The value of \texttt{Barrier\_Limit’Last} in \texttt{Synchronous\_Barriers}. See \texttt{D.10.1(4)}.”

\texttt{Synchronous\_Barriers.Barrier\_Limit’Last} is \texttt{Integer’Last}.
* “When an aborted task that is waiting on a \texttt{Synchronous\_Barrier} is aborted. See \texttt{D.10.1(13)}.”

Difficult to characterize.
* “The value of \texttt{Min\_Handler\_Ceiling} in \texttt{Execution\_Time\_Group\_Budgets}. See \texttt{D.14.2(7)}.”

See source file \texttt{a-etgrbu.ads}.
* “The value of \texttt{CPU\_Range’Last} in \texttt{System\_Multiprocessors}. See \texttt{D.16(4)}.”

See source file \texttt{s-multip.ads}.
* “The processor on which the environment task executes in the absence of a value for the aspect \texttt{CPU}. See \texttt{D.16(13)}.”

Unknown.
* “The means for creating and executing distributed programs. See \texttt{E(5)}.”

The \texttt{GLADE} package provides a utility \texttt{GNATDIST} for creating and executing distributed programs. See the \texttt{GLADE} reference manual for further details.
* “Any events that can result in a partition becoming inaccessible. See \texttt{E.1(7)}.”

See the \texttt{GLADE} reference manual for full details on such events.
* “The scheduling policies, treatment of priorities, and management of shared resources between partitions in certain cases. See \texttt{E.1(11)}.”
See the GLADE reference manual for full details on these aspects of multi-partition execution.

* “Whether the execution of the remote subprogram is immediately aborted as a result of cancellation. See E.4(13).”

See the GLADE reference manual for details on the effect of abort in a distributed application.

* “The range of type System.RPC.Partition_Id. See E.5(14).”

System.RPC.Partition_ID’Last is Integer’Last. See source file s-rpc.ads.

* “Implementation-defined interfaces in the PCS. See E.5(26).”

See the GLADE reference manual for a full description of all implementation defined interfaces.

* “The values of named numbers in the package Decimal. See F.2(7).”

<table>
<thead>
<tr>
<th>Named Number</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>‘Max_Scale’</td>
<td>+18</td>
</tr>
<tr>
<td>‘Min_Scale’</td>
<td>-18</td>
</tr>
<tr>
<td>‘Min_Delta’</td>
<td>1.0E-18</td>
</tr>
<tr>
<td>‘Max_Delta’</td>
<td>1.0E+18</td>
</tr>
<tr>
<td>‘Max_Decimal_Digits’</td>
<td>18</td>
</tr>
</tbody>
</table>

* “The value of Max_Picture_Length in the package Text_IO.Editing. See F.3.3(16).”

64

* “The value of Max_Picture_Length in the package Wide_Text_IO.Editing. See F.3.4(5).”

64

* “The accuracy actually achieved by the complex elementary functions and by other complex arithmetic operations. See G.1(1).”

Standard library functions are used for the complex arithmetic operations. Only fast math mode is currently supported.

* “The sign of a zero result (or a component thereof) from any operator or function in Numerics.Generic_Complex_Types, when Real’Signed_Zeros is True. See G.1.1(53).”

The signs of zero values are as recommended by the relevant implementation advice.

* “The sign of a zero result (or a component thereof) from any operator or function in Numerics.Generic_Complex_Elementary_Functions, when Real’Signed_Zeros is True. See G.1.2(45).”

The signs of zero values are as recommended by the relevant implementation advice.
* “Whether the strict mode or the relaxed mode is the default. See G.2(2).”

The strict mode is the default. There is no separate relaxed mode. GNAT provides a highly efficient implementation of strict mode.

* “The result interval in certain cases of fixed-to-float conversion. See G.2.1(10).”

For cases where the result interval is implementation dependent, the accuracy is that provided by performing all operations in 64-bit IEEE floating-point format.

* “The result of a floating point arithmetic operation in overflow situations, when the Machine_Overflows attribute of the result type is False. See G.2.1(13).”

Infinite and NaN values are produced as dictated by the IEEE floating-point standard. Note that on machines that are not fully compliant with the IEEE floating-point standard, such as Alpha, the ‘-mieee’ compiler flag must be used for achieving IEEE conforming behavior (although at the cost of a significant performance penalty), so infinite and NaN values are properly generated.

* “The result interval for division (or exponentiation by a negative exponent), when the floating point hardware implements division as multiplication by a reciprocal. See G.2.1(16).”

Not relevant, division is IEEE exact.

* “The definition of close result set, which determines the accuracy of certain fixed point multiplications and divisions. See G.2.3(5).”

Operations in the close result set are performed using IEEE long format floating-point arithmetic. The input operands are converted to floating-point, the operation is done in floating-point, and the result is converted to the target type.

* “Conditions on a ‘universal_real’ operand of a fixed point multiplication or division for which the result shall be in the perfect result set. See G.2.3(22).”

The result is only defined to be in the perfect result set if the result can be computed by a single scaling operation involving a scale factor representable in 64 bits.

* “The result of a fixed point arithmetic operation in overflow situations, when the Machine_Overflows attribute of the result type is False. See G.2.3(27).”

Not relevant, Machine_Overflows is True for fixed-point types.

* “The result of an elementary function reference in overflow situations, when the Machine_Overflows attribute of the result type is False. See G.2.4(4).”

IEEE infinite and Nan values are produced as appropriate.

* “The value of the angle threshold, within which certain elementary functions, complex arithmetic operations, and complex elementary functions yield results conforming to a maximum relative error bound. See G.2.4(10).”

Information on this subject is not yet available.

* “The accuracy of certain elementary functions for parameters beyond the angle threshold. See G.2.4(10).”

Information on this subject is not yet available.
The result of a complex arithmetic operation or complex elementary function reference in overflow situations, when the `Machine_Overflows` attribute of the corresponding real type is `False`. See G.2.6(5)."

IEEE infinite and Nan values are produced as appropriate.

The accuracy of certain complex arithmetic operations and certain complex elementary functions for parameters (or components thereof) beyond the angle threshold. See G.2.6(8)."

Information on those subjects is not yet available.

The accuracy requirements for the subprograms Solve, Inverse, Determinant, Eigenvalues and Eigensystem for type `Real_Matrix`. See G.3.1(81)."

Information on those subjects is not yet available.

The accuracy requirements for the subprograms Solve, Inverse, Determinant, Eigenvalues and Eigensystem for type `Complex_Matrix`. See G.3.2(149)."

Information on those subjects is not yet available.

The consequences of violating `No_Hidden_Indirect_Globals`. See H.4(23.9)."

Execution is erroneous in that case.
Chapter 8: Intrinsic Subprograms

8 Intrinsic Subprograms

GNAT allows a user application program to write the declaration:

```
pragma Import (Intrinsic, name);
```

providing that the name corresponds to one of the implemented intrinsic subprograms in GNAT, and that the parameter profile of the referenced subprogram meets the requirements. This chapter describes the set of implemented intrinsic subprograms, and the requirements on parameter profiles. Note that no body is supplied; as with other uses of pragma Import, the body is supplied elsewhere (in this case by the compiler itself). Note that any use of this feature is potentially non-portable, since the Ada standard does not require Ada compilers to implement this feature.

8.1 Intrinsic Operators

All the predefined numeric operators in package Standard in `pragma Import (Intrinsic,..)` declarations. In the binary operator case, the operands must have the same size. The operand or operands must also be appropriate for the operator. For example, for addition, the operands must both be floating-point or both be fixed-point, and the right operand for "**" must have a root type of Standard.Integer'Base. You can use an intrinsic operator declaration as in the following example:

```
type Int1 is new Integer;
type Int2 is new Integer;

function "+" (X1 : Int1; X2 : Int2) return Int1;
function "+" (X1 : Int1; X2 : Int2) return Int2;
pragma Import (Intrinsic, "+");
```

This declaration would permit ‘mixed mode’ arithmetic on items of the differing types Int1 and Int2. It is also possible to specify such operators for private types, if the full views are appropriate arithmetic types.

8.2 Compilation_ISO_Date

This intrinsic subprogram is used in the implementation of the library package GNAT.Source_Info. The only useful use of the intrinsic import in this case is the one in this unit, so an application program should simply call the function GNAT.Source_Info.Compilation_ISO_Date to obtain the date of the current compilation (in local time format YYYY-MM-DD).

8.3 Compilation_Date

Same as Compilation_ISO_Date, except the string is in the form MMM DD YYYY.

8.4 Compilation_Time

This intrinsic subprogram is used in the implementation of the library package GNAT.Source_Info. The only useful use of the intrinsic import in this case is the one
in this unit, so an application program should simply call the function `GNAT.Source_Info.Compilation_Time` to obtain the time of the current compilation (in local time format HH:MM:SS).

8.5 Enclosing_Entity

This intrinsic subprogram is used in the implementation of the library package `GNAT.Source_Info`. The only useful use of the intrinsic import in this case is the one in this unit, so an application program should simply call the function `GNAT.Source_Info.Enclosing_Entity` to obtain the name of the current subprogram, package, task, entry, or protected subprogram.

8.6 Exception_Information

This intrinsic subprogram is used in the implementation of the library package `GNAT.Current_Exception`. The only useful use of the intrinsic import in this case is the one in this unit, so an application program should simply call the function `GNAT.Current_Exception.Exception_Information` to obtain the exception information associated with the current exception.

8.7 Exception_Message

This intrinsic subprogram is used in the implementation of the library package `GNAT.Current_Exception`. The only useful use of the intrinsic import in this case is the one in this unit, so an application program should simply call the function `GNAT.Current_Exception.Exception_Message` to obtain the message associated with the current exception.

8.8 Exception_Name

This intrinsic subprogram is used in the implementation of the library package `GNAT.Current_Exception`. The only useful use of the intrinsic import in this case is the one in this unit, so an application program should simply call the function `GNAT.Current_Exception.Exception_Name` to obtain the name of the current exception.

8.9 File

This intrinsic subprogram is used in the implementation of the library package `GNAT.Source_Info`. The only useful use of the intrinsic import in this case is the one in this unit, so an application program should simply call the function `GNAT.Source_Info.File` to obtain the name of the current file.

8.10 Line

This intrinsic subprogram is used in the implementation of the library package `GNAT.Source_Info`. The only useful use of the intrinsic import in this case is the one in this unit, so an application program should simply call the function `GNAT.Source_Info.Line` to obtain the number of the current source line.
8.11 Shifts and Rotates

In standard Ada, the shift and rotate functions are available only for the predefined modular types in package *Interfaces*. However, in GNAT it is possible to define these functions for any integer type (signed or modular), as in this example:

```ada
function Shift_Left
    (Value : T;
     Amount : Natural) return T
with Import, Convention => Intrinsic;
```

The function name must be one of `Shift_Left`, `Shift_Right`, `Shift_Right_Arithmetic`, `Rotate_Left`, or `Rotate_Right`. `T` must be an integer type. `T'Size` must be 8, 16, 32 or 64 bits; if `T` is modular, the modulus must be `2**8`, `2**16`, `2**32` or `2**64`. The result type must be the same as the type of `Value`. The shift amount must be `Natural`. The formal parameter names can be anything.

A more convenient way of providing these shift operators is to use the `Provide_Shift_Operators` pragma, which provides the function declarations and corresponding `pragma Import`'s for all five shift functions. For signed types the semantics of these operators is to interpret the bitwise result of the corresponding operator for modular type. In particular, shifting a negative number may change its sign bit to positive.

8.12 Source_Location

This intrinsic subprogram is used in the implementation of the library routine `GNAT.Source_Info`. The only useful use of the intrinsic import in this case is the one in this unit, so an application program should simply call the function `GNAT.Source_Info.Source_Location` to obtain the current source file location.
9 Representation Clauses and Pragmas

This section describes the representation clauses accepted by GNAT, and their effect on the representation of corresponding data objects.

GNAT fully implements Annex C (Systems Programming). This means that all the implementation advice sections in chapter 13 are fully implemented. However, these sections only require a minimal level of support for representation clauses. GNAT provides much more extensive capabilities, and this section describes the additional capabilities provided.

9.1 Alignment Clauses

GNAT requires that all alignment clauses specify 0 or a power of 2, and all default alignments are always a power of 2. Specifying 0 is the same as specifying 1.

The default alignment values are as follows:

* ‘Elementary Types’.

For elementary types, the alignment is the minimum of the actual size of objects of the type divided by Storage_Unit, and the maximum alignment supported by the target. (This maximum alignment is given by the GNAT-specific attribute Standard'Maximum_Alignment; see [Attribute Maximum_Alignment], page 124.)

For example, for type Long_Float, the object size is 8 bytes, and the default alignment will be 8 on any target that supports alignments this large, but on some targets, the maximum alignment may be smaller than 8, in which case objects of type Long_Float will be maximally aligned.

* ‘Arrays’.

For arrays, the alignment is equal to the alignment of the component type for the normal case where no packing or component size is given. If the array is packed, and the packing is effective (see separate section on packed arrays), then the alignment will be either 4, 2, or 1 for long packed arrays or arrays whose length is not known at compile time, depending on whether the component size is divisible by 4, 2, or is odd. For short packed arrays, which are handled internally as modular types, the alignment will be as described for elementary types, e.g. a packed array of length 31 bits will have an object size of four bytes, and an alignment of 4.

* ‘Records’.

For the normal unpacked case, the alignment of a record is equal to the maximum alignment of any of its components. For tagged records, this includes the implicit access type used for the tag. If a pragma Pack is used and all components are packable (see separate section on pragma Pack), then the resulting alignment is 1, unless the layout of the record makes it profitable to increase it.

A special case is when:

* the size of the record is given explicitly, or a full record representation clause is given, and

* the size of the record is 2, 4, or 8 bytes.

In this case, an alignment is chosen to match the size of the record. For example, if we have:
type Small is record
  A, B : Character;
end record;
for Small'Size use 16;
then the default alignment of the record type Small is 2, not 1. This leads to more efficient code when the record is treated as a unit, and also allows the type to specified as Atomic on architectures requiring strict alignment.

An alignment clause may specify a larger alignment than the default value up to some maximum value dependent on the target (obtainable by using the attribute reference Standard'Maximum_Alignment). It may also specify a smaller alignment than the default value for enumeration, integer and fixed point types, as well as for record types, for example

type V is record
  A : Integer;
end record;
for V'alignment use 1;
The default alignment for the type V is 4, as a result of the Integer field in the record, but it is permissible, as shown, to override the default alignment of the record with a smaller value.

Note that according to the Ada standard, an alignment clause applies only to the first named subtype. If additional subtypes are declared, then the compiler is allowed to choose any alignment it likes, and there is no way to control this choice. Consider:

type R is range 1 .. 10_000;
for R'Alignment use 1;
subtype RS is R range 1 .. 1000;
The alignment clause specifies an alignment of 1 for the first named subtype R but this does not necessarily apply to RS. When writing portable Ada code, you should avoid writing code that explicitly or implicitly relies on the alignment of such subtypes.

For the GNAT compiler, if an explicit alignment clause is given, this value is also used for any subsequent subtypes. So for GNAT, in the above example, you can count on the alignment of RS being 1. But this assumption is non-portable, and other compilers may choose different alignments for the subtype RS.

9.2 Size Clauses

The default size for a type T is obtainable through the language-defined attribute T'Size and also through the equivalent GNAT-defined attribute T'Value_Size. For objects of type T, GNAT will generally increase the type size so that the object size (obtainable through the GNAT-defined attribute T'Object_Size) is a multiple of T'Alignment * Storage_Unit. For example:

type Smallint is range 1 .. 6;

type Rec is record
  Y1 : integer;
Y2 : boolean;
end record;

In this example, Smallint’Size = Smallint’Value_Size = 3, as specified by the RM rules, but objects of this type will have a size of 8 (Smallint’Object_Size = 8), since objects by default occupy an integral number of storage units. On some targets, notably older versions of the Digital Alpha, the size of stand alone objects of this type may be 32, reflecting the inability of the hardware to do byte load/stores.

Similarly, the size of type Rec is 40 bits (Rec’Size = Rec’Value_Size = 40), but the alignment is 4, so objects of this type will have their size increased to 64 bits so that it is a multiple of the alignment (in bits). This decision is in accordance with the specific Implementation Advice in RM 13.3(43):

“A Size clause should be supported for an object if the specified Size is at least as large as its subtype’s Size, and corresponds to a size in storage elements that is a multiple of the object’s Alignment (if the Alignment is nonzero).”

An explicit size clause may be used to override the default size by increasing it. For example, if we have:

type My_Boolean is new Boolean;
for My_Boolean'Size use 32;

then values of this type will always be 32-bit long. In the case of discrete types, the size can be increased up to 64 bits on 32-bit targets and 128 bits on 64-bit targets, with the effect that the entire specified field is used to hold the value, sign- or zero-extended as appropriate. If more than 64 bits or 128 bits resp. is specified, then padding space is allocated after the value, and a warning is issued that there are unused bits.

Similarly the size of records and arrays may be increased, and the effect is to add padding bits after the value. This also causes a warning message to be generated.

The largest Size value permitted in GNAT is $2^{31}-1$. Since this is a Size in bits, this corresponds to an object of size 256 megabytes (minus one). This limitation is true on all targets. The reason for this limitation is that it improves the quality of the code in many cases if it is known that a Size value can be accommodated in an object of type Integer.

9.3 Storage_Size Clauses

For tasks, the Storage_Size clause specifies the amount of space to be allocated for the task stack. This cannot be extended, and if the stack is exhausted, then Storage_Error will be raised (if stack checking is enabled). Use a Storage_Size attribute definition clause, or a Storage_Size pragma in the task definition to set the appropriate required size. A useful technique is to include in every task definition a pragma of the form:

pragma Storage_Size (Default_Stack_Size);

Then Default_Stack_Size can be defined in a global package, and modified as required. Any tasks requiring stack sizes different from the default can have an appropriate alternative reference in the pragma.

You can also use the ‘-d’ binder switch to modify the default stack size.

For access types, the Storage_Size clause specifies the maximum space available for allocation of objects of the type. If this space is exceeded then Storage_Error will be raised by
an allocation attempt. In the case where the access type is declared local to a subprogram, the use of a Storage_Size clause triggers automatic use of a special predefined storage pool (System.Pool_Size) that ensures that all space for the pool is automatically reclaimed on exit from the scope in which the type is declared.

A special case recognized by the compiler is the specification of a Storage_Size of zero for an access type. This means that no items can be allocated from the pool, and this is recognized at compile time, and all the overhead normally associated with maintaining a fixed size storage pool is eliminated. Consider the following example:

```pascal
procedure p is
  type R is array (Natural) of Character;
  type P is access all R;
  for P'Storage_Size use 0;
  -- Above access type intended only for interfacing purposes
  y : P;

  procedure g (m : P);
  pragma Import (C, g);
  -- ...

begin
    -- ...
    y := new R;
end;
```

As indicated in this example, these dummy storage pools are often useful in connection with interfacing where no object will ever be allocated. If you compile the above example, you get the warning:

```
  p.adb:16:09: warning: allocation from empty storage pool
  p.adb:16:09: warning: Storage_Error will be raised at run time
```

Of course in practice, there will not be any explicit allocators in the case of such an access declaration.

### 9.4 Size of Variant Record Objects

In the case of variant record objects, there is a question whether Size gives information about a particular variant, or the maximum size required for any variant. Consider the following program

```pascal
with Text_IO; use Text_IO;
procedure q is
  type R1 (A : Boolean := False) is record
    case A is
      when True => X : Character;
      when False => null;
    end case;
  end record;
```
V1 : R1 (False);
V2 : R1;

begin
  Put_Line (Integer'Image (V1'Size));
  Put_Line (Integer'Image (V2'Size));
end q;

Here we are dealing with a variant record, where the True variant requires 16 bits, and the False variant requires 8 bits. In the above example, both V1 and V2 contain the False variant, which is only 8 bits long. However, the result of running the program is:

8
16

The reason for the difference here is that the discriminant value of V1 is fixed, and will always be False. It is not possible to assign a True variant value to V1, therefore 8 bits is sufficient. On the other hand, in the case of V2, the initial discriminant value is False (from the default), but it is possible to assign a True variant value to V2, therefore 16 bits must be allocated for V2 in the general case, even fewer bits may be needed at any particular point during the program execution.

As can be seen from the output of this program, the 'Size attribute applied to such an object in GNAT gives the actual allocated size of the variable, which is the largest size of any of the variants. The Ada Reference Manual is not completely clear on what choice should be made here, but the GNAT behavior seems most consistent with the language in the RM.

In some cases, it may be desirable to obtain the size of the current variant, rather than the size of the largest variant. This can be achieved in GNAT by making use of the fact that in the case of a subprogram parameter, GNAT does indeed return the size of the current variant (because a subprogram has no way of knowing how much space is actually allocated for the actual).

Consider the following modified version of the above program:

with Text_IO; use Text_IO;
procedure q is
  type R1 (A : Boolean := False) is record
    case A is
      when True  => X : Character;
      when False => null;
    end case;
  end record;

  V2 : R1;

  function Size (V : R1) return Integer is
  begin
    return V'Size;
  end Size;

V1 : R1 (False);
V2 : R1;

begin
  Put_Line (Integer'Image (V1'Size));
  Put_Line (Integer'Image (V2'Size));
end q;
Chapter 9: Representation Clauses and Pragmas

```
begin
  Put_Line (Integer'Image (V2'Size));
  Put_Line (Integer'Image (Size (V2)));
  V2 := (True, 'x');
  Put_Line (Integer'Image (V2'Size));
  Put_Line (Integer'Image (Size (V2)));
end q;
```

The output from this program is

```
16
8
16
16
```

Here we see that while the 'Size attribute always returns the maximum size, regardless of the current variant value, the Size function does indeed return the size of the current variant value.

### 9.5 Biased Representation

In the case of scalars with a range starting at other than zero, it is possible in some cases to specify a size smaller than the default minimum value, and in such cases, GNAT uses an unsigned biased representation, in which zero is used to represent the lower bound, and successive values represent successive values of the type.

For example, suppose we have the declaration:

```
type Small is range -7 .. -4;
for Small'Size use 2;
```

Although the default size of type Small is 4, the Size clause is accepted by GNAT and results in the following representation scheme:

```
-7 is represented as 2#00#
-6 is represented as 2#01#
-5 is represented as 2#10#
-4 is represented as 2#11#
```

Biased representation is only used if the specified Size clause cannot be accepted in any other manner. These reduced sizes that force biased representation can be used for all discrete types except for enumeration types for which a representation clause is given.

### 9.6 Value_Size and Object_Size Clauses

In Ada 95 and Ada 2005, T'Size for a type T is the minimum number of bits required to hold values of type T. Although this interpretation was allowed in Ada 83, it was not required, and this requirement in practice can cause some significant difficulties. For example, in most Ada 83 compilers, Natural'Size was 32. However, in Ada 95 and Ada 2005, Natural'Size is typically 31. This means that code may change in behavior when moving from Ada 83 to Ada 95 or Ada 2005. For example, consider:

```
type Rec is record
  A : Natural;
```
B : Natural;
end record;

for Rec use record
A at 0 range 0 .. Natural’Size - 1;
B at 0 range Natural’Size .. 2 * Natural’Size - 1;
end record;

In the above code, since the typical size of Natural objects is 32 bits and Natural’Size is 31, the above code can cause unexpected inefficient packing in Ada 95 and Ada 2005, and in general there are cases where the fact that the object size can exceed the size of the type causes surprises.

To help get around this problem GNAT provides two implementation defined attributes, Value_Size and Object_Size. When applied to a type, these attributes yield the size of the type (corresponding to the RM defined size attribute), and the size of objects of the type respectively.

The Object_Size is used for determining the default size of objects and components. This size value can be referred to using the Object_Size attribute. The phrase ‘is used’ here means that it is the basis of the determination of the size. The backend is free to pad this up if necessary for efficiency, e.g., an 8-bit stand-alone character might be stored in 32 bits on a machine with no efficient byte access instructions such as the Alpha.

The default rules for the value of Object_Size for discrete types are as follows:

* The Object_Size for base subtypes reflect the natural hardware size in bits (run the compiler with ‘-gnatS’ to find those values for numeric types). Enumeration types and fixed-point base subtypes have 8, 16, 32, or 64 bits for this size, depending on the range of values to be stored.

* The Object_Size of a subtype is the same as the Object_Size of the type from which it is obtained.

* The Object_Size of a derived base type is copied from the parent base type, and the Object_Size of a derived first subtype is copied from the parent first subtype.

The Value_Size attribute is the (minimum) number of bits required to store a value of the type. This value is used to determine how tightly to pack records or arrays with components of this type, and also affects the semantics of unchecked conversion (unchecked conversions where the Value_Size values differ generate a warning, and are potentially target dependent).

The default rules for the value of Value_Size are as follows:

* The Value_Size for a base subtype is the minimum number of bits required to store all values of the type (including the sign bit only if negative values are possible).

* If a subtype statically matches the first subtype of a given type, then it has by default the same Value_Size as the first subtype. (This is a consequence of RM 13.1(14): “if two subtypes statically match, then their subtype-specific aspects are the same”.)

* All other subtypes have a Value_Size corresponding to the minimum number of bits required to store all values of the subtype. For dynamic bounds, it is assumed that the value can range down or up to the corresponding bound of the ancestor.
The RM defined attribute `Size` corresponds to the `Value_Size` attribute.

The `Size` attribute may be defined for a first-named subtype. This sets the `Value_Size` of the first-named subtype to the given value, and the `Object_Size` of this first-named subtype to the given value padded up to an appropriate boundary. It is a consequence of the default rules above that this `Object_Size` will apply to all further subtypes. On the other hand, `Value_Size` is affected only for the first subtype, any dynamic subtypes obtained from it directly, and any statically matching subtypes. The `Value_Size` of any other static subtypes is not affected.

`Value_Size` and `Object_Size` may be explicitly set for any subtype using an attribute definition clause. Note that the use of these attributes can cause the RM 13.1(14) rule to be violated. If two access types reference aliased objects whose subtypes have differing `Object_Size` values as a result of explicit attribute definition clauses, then it is illegal to convert from one access subtype to the other. For a more complete description of this additional legality rule, see the description of the `Object_Size` attribute.

To get a feel for the difference, consider the following examples (note that in each case the base is `Short_Short_Integer` with a size of 8):

<table>
<thead>
<tr>
<th>Type or subtype declaration</th>
<th>Object_Size</th>
<th>Value_Size</th>
</tr>
</thead>
<tbody>
<tr>
<td>type x1 is range 0 .. 5;</td>
<td>8</td>
<td>3</td>
</tr>
<tr>
<td>type x2 is range 0 .. 5; for x2’size use 12;</td>
<td>16</td>
<td>12</td>
</tr>
<tr>
<td>subtype x3 is x2 range 0 .. 3;</td>
<td>16</td>
<td>2</td>
</tr>
<tr>
<td>subtype x4 is x2’base range 0 .. 10;</td>
<td>8</td>
<td>4</td>
</tr>
<tr>
<td>dynamic : x2’Base range -64 .. +63;</td>
<td></td>
<td></td>
</tr>
<tr>
<td>subtype x5 is x2 range 0 .. dynamic;</td>
<td>16</td>
<td>3*</td>
</tr>
<tr>
<td>subtype x6 is x2’base range 0 .. dynamic;</td>
<td>8</td>
<td>7*</td>
</tr>
</tbody>
</table>

Note: the entries marked ‘*’ are not actually specified by the Ada Reference Manual, which has nothing to say about size in the dynamic case. What GNAT does is to allocate sufficient bits to accommodate any possible dynamic values for the bounds at run-time.

So far, so good, but GNAT has to obey the RM rules, so the question is under what conditions must the RM `Size` be used. The following is a list of the occasions on which the RM `Size` must be used:

* Component size for packed arrays or records
* Value of the attribute `Size` for a type
* Warning about sizes not matching for unchecked conversion

For record types, the `Object_Size` is always a multiple of the alignment of the type (this is true for all types). In some cases the `Value_Size` can be smaller. Consider:
type R is record  
    X : Integer;  
    Y : Character;  
end record;

On a typical 32-bit architecture, the X component will occupy four bytes and the Y component will occupy one byte, for a total of 5 bytes. As a result R'Value_Size will be 40 (bits) since this is the minimum size required to store a value of this type. For example, it is permissible to have a component of type R in an array whose component size is specified to be 40 bits.

However, R'Object_Size will be 64 (bits). The difference is due to the alignment requirement for objects of the record type. The X component will require four-byte alignment because that is what type Integer requires, whereas the Y component, a Character, will only require 1-byte alignment. Since the alignment required for X is the greatest of all the components' alignments, that is the alignment required for the enclosing record type, i.e., 4 bytes or 32 bits. As indicated above, the actual object size must be rounded up so that it is a multiple of the alignment value. Therefore, 40 bits rounded up to the next multiple of 32 yields 64 bits.

For all other types, the Object_Size and Value_Size are the same (and equivalent to the RM attribute Size). Only Size may be specified for such types.

Note that Value_Size can be used to force biased representation for a particular subtype. Consider this example:

    type R is (A, B, C, D, E, F);  
    subtype RAB is R range A .. B;  
    subtype REF is R range E .. F;

By default, RAB has a size of 1 (sufficient to accommodate the representation of A and B, 0 and 1), and REF has a size of 3 (sufficient to accommodate the representation of E and F, 4 and 5). But if we add the following Value_Size attribute definition clause:

    for REF'Value_Size use 1;

then biased representation is forced for REF, and 0 will represent E and 1 will represent F. A warning is issued when a Value_Size attribute definition clause forces biased representation. This warning can be turned off using -gnatw.B.

9.7 Component_Size Clauses

Normally, the value specified in a component size clause must be consistent with the subtype of the array component with regard to size and alignment. In other words, the value specified must be at least equal to the size of this subtype, and must be a multiple of the alignment value.

In addition, component size clauses are allowed which cause the array to be packed, by specifying a smaller value. A first case is for component size values in the range 1 through 63 on 32-bit targets, and 1 through 127 on 64-bit targets. The value specified may not be smaller than the Size of the subtype. GNAT will accurately honor all packing requests in this range. For example, if we have:

    type r is array (1 .. 8) of Natural;  
    for r'Component_Size use 31;
then the resulting array has a length of 31 bytes (248 bits = 8 * 31). Of course access to the components of such an array is considerably less efficient than if the natural component size of 32 is used. A second case is when the subtype of the component is a record type padded because of its default alignment. For example, if we have:

```pascal
type r is record
  i : Integer;
  j : Integer;
  b : Boolean;
end record;

type a is array (1 .. 8) of r;
for a'Component_Size use 72;
```

then the resulting array has a length of 72 bytes, instead of 96 bytes if the alignment of the record (4) was obeyed.

Note that there is no point in giving both a component size clause and a pragma Pack for the same array type. If such duplicate clauses are given, the pragma Pack will be ignored.

### 9.8 Bit_Order Clauses

For record subtypes, GNAT permits the specification of the Bit_Order attribute. The specification may either correspond to the default bit order for the target, in which case the specification has no effect and places no additional restrictions, or it may be for the non-standard setting (that is the opposite of the default).

In the case where the non-standard value is specified, the effect is to renumber bits within each byte, but the ordering of bytes is not affected. There are certain restrictions placed on component clauses as follows:

* Components fitting within a single storage unit.

These are unrestricted, and the effect is merely to renumber bits. For example if we are on a little-endian machine with Low_Order_First being the default, then the following two declarations have exactly the same effect:

```pascal
type R1 is record
  A : Boolean;
  B : Integer range 1 .. 120;
end record;

for R1 use record
  A at 0 range 0 .. 0;
  B at 0 range 1 .. 7;
end record;

type R2 is record
  A : Boolean;
  B : Integer range 1 .. 120;
end record;

for R2'Bit_Order use High_Order_First;
```
Chapter 9: Representation Clauses and Pragmas

for R2 use record
   A at 0 range 7 .. 7;
   B at 0 range 0 .. 6;
end record;

The useful application here is to write the second declaration with the `Bit_Order` attribute definition clause, and know that it will be treated the same, regardless of whether the target is little-endian or big-endian.

* Components occupying an integral number of bytes.

These are components that exactly fit in two or more bytes. Such component declarations are allowed, but have no effect, since it is important to realize that the `Bit_Order` specification does not affect the ordering of bytes. In particular, the following attempt at getting an endian-independent integer does not work:

```
for R2'Bit_Order use High_Order_First;
```

This declaration will result in a little-endian integer on a little-endian machine, and a big-endian integer on a big-endian machine. If byte flipping is required for interoperability between big- and little-endian machines, this must be explicitly programmed. This capability is not provided by `Bit_Order`.

* Components that are positioned across byte boundaries.

but do not occupy an integral number of bytes. Given that bytes are not reordered, such fields would occupy a non-contiguous sequence of bits in memory, requiring non-trivial code to reassemble. They are for this reason not permitted, and any component clause specifying such a layout will be flagged as illegal by GNAT.

Since the misconception that `Bit_Order` automatically deals with all endian-related incompatibilities is a common one, the specification of a component field that is an integral number of bytes will always generate a warning. This warning may be suppressed using `pragma Warnings (Off)` if desired. The following section contains additional details regarding the issue of byte ordering.

9.9 Effect of `Bit_Order` on Byte Ordering

In this section we will review the effect of the `Bit_Order` attribute definition clause on byte ordering. Briefly, it has no effect at all, but a detailed example will be helpful. Before giving this example, let us review the precise definition of the effect of defining `Bit_Order`. The effect of a non-standard bit order is described in section 13.5.3 of the Ada Reference Manual:

"2 A bit ordering is a method of interpreting the meaning of the storage place attributes."
To understand the precise definition of storage place attributes in this context, we visit section 13.5.1 of the manual:

“13 A record_representation_clause (without the mod_clause) specifies the layout. The storage place attributes (see 13.5.2) are taken from the values of the position, first_bit, and last_bit expressions after normalizing those values so that first_bit is less than Storage_Unit.”

The critical point here is that storage places are taken from the values after normalization, not before. So the Bit_Order interpretation applies to normalized values. The interpretation is described in the later part of the 13.5.3 paragraph:

“2 A bit ordering is a method of interpreting the meaning of the storage place attributes. High_Order_First (known in the vernacular as ‘big endian’) means that the first bit of a storage element (bit 0) is the most significant bit (interpreting the sequence of bits that represent a component as an unsigned integer value). Low_Order_First (known in the vernacular as ‘little endian’) means the opposite: the first bit is the least significant.”

Note that the numbering is with respect to the bits of a storage unit. In other words, the specification affects only the numbering of bits within a single storage unit.

We can make the effect clearer by giving an example.

Suppose that we have an external device which presents two bytes, the first byte presented, which is the first (low addressed byte) of the two byte record is called Master, and the second byte is called Slave.

The left most (most significant) bit is called Control for each byte, and the remaining 7 bits are called V1, V2, .. V7, where V7 is the rightmost (least significant) bit.

On a big-endian machine, we can write the following representation clause

```haskell
type Data is record
  Master_Control : Bit;
  Master_V1 : Bit;
  Master_V2 : Bit;
  Master_V3 : Bit;
  Master_V4 : Bit;
  Master_V5 : Bit;
  Master_V6 : Bit;
  Master_V7 : Bit;
  Slave_Control : Bit;
  Slave_V1 : Bit;
  Slave_V2 : Bit;
  Slave_V3 : Bit;
  Slave_V4 : Bit;
  Slave_V5 : Bit;
  Slave_V6 : Bit;
  Slave_V7 : Bit;
end record;

for Data use record
```

```
Master_Control at 0 range 0 .. 0;
Master_V1 at 0 range 1 .. 1;
Master_V2 at 0 range 2 .. 2;
Master_V3 at 0 range 3 .. 3;
Master_V4 at 0 range 4 .. 4;
Master_V5 at 0 range 5 .. 5;
Master_V6 at 0 range 6 .. 6;
Master_V7 at 0 range 7 .. 7;
Slave_Control at 1 range 0 .. 0;
Slave_V1 at 1 range 1 .. 1;
Slave_V2 at 1 range 2 .. 2;
Slave_V3 at 1 range 3 .. 3;
Slave_V4 at 1 range 4 .. 4;
Slave_V5 at 1 range 5 .. 5;
Slave_V6 at 1 range 6 .. 6;
Slave_V7 at 1 range 7 .. 7;
end record;

Now if we move this to a little endian machine, then the bit ordering within the byte is backwards, so we have to rewrite the record rep clause as:

for Data use record
  Master_Control at 0 range 7 .. 7;
  Master_V1 at 0 range 6 .. 6;
  Master_V2 at 0 range 5 .. 5;
  Master_V3 at 0 range 4 .. 4;
  Master_V4 at 0 range 3 .. 3;
  Master_V5 at 0 range 2 .. 2;
  Master_V6 at 0 range 1 .. 1;
  Master_V7 at 0 range 0 .. 0;
  Slave_Control at 1 range 7 .. 7;
  Slave_V1 at 1 range 6 .. 6;
  Slave_V2 at 1 range 5 .. 5;
  Slave_V3 at 1 range 4 .. 4;
  Slave_V4 at 1 range 3 .. 3;
  Slave_V5 at 1 range 2 .. 2;
  Slave_V6 at 1 range 1 .. 1;
  Slave_V7 at 1 range 0 .. 0;
end record;

It is a nuisance to have to rewrite the clause, especially if the code has to be maintained on both machines. However, this is a case that we can handle with the Bit_Order attribute if it is implemented. Note that the implementation is not required on byte addressed machines, but it is indeed implemented in GNAT. This means that we can simply use the first record clause, together with the declaration

for Data’Bit_Order use High_Order_First;

and the effect is what is desired, namely the layout is exactly the same, independent of whether the code is compiled on a big-endian or little-endian machine.
The important point to understand is that byte ordering is not affected. A Bit_Order attribute definition never affects which byte a field ends up in, only where it ends up in that byte. To make this clear, let us rewrite the record rep clause of the previous example as:

```haskell
for Data’Bit_Order use High_Order_First;
for Data use record
    Master_Control at 0 range 0 .. 0;
    Master_V1 at 0 range 1 .. 1;
    Master_V2 at 0 range 2 .. 2;
    Master_V3 at 0 range 3 .. 3;
    Master_V4 at 0 range 4 .. 4;
    Master_V5 at 0 range 5 .. 5;
    Master_V6 at 0 range 6 .. 6;
    Master_V7 at 0 range 7 .. 7;
    Slave_Control at 0 range 8 .. 8;
    Slave_V1 at 0 range 9 .. 9;
    Slave_V2 at 0 range 10 .. 10;
    Slave_V3 at 0 range 11 .. 11;
    Slave_V4 at 0 range 12 .. 12;
    Slave_V5 at 0 range 13 .. 13;
    Slave_V6 at 0 range 14 .. 14;
    Slave_V7 at 0 range 15 .. 15;
end record;
```

This is exactly equivalent to saying (a repeat of the first example):

```haskell
for Data’Bit_Order use High_Order_First;
for Data use record
    Master_Control at 0 range 0 .. 0;
    Master_V1 at 0 range 1 .. 1;
    Master_V2 at 0 range 2 .. 2;
    Master_V3 at 0 range 3 .. 3;
    Master_V4 at 0 range 4 .. 4;
    Master_V5 at 0 range 5 .. 5;
    Master_V6 at 0 range 6 .. 6;
    Master_V7 at 0 range 7 .. 7;
    Slave_Control at 1 range 0 .. 0;
    Slave_V1 at 1 range 1 .. 1;
    Slave_V2 at 1 range 2 .. 2;
    Slave_V3 at 1 range 3 .. 3;
    Slave_V4 at 1 range 4 .. 4;
    Slave_V5 at 1 range 5 .. 5;
    Slave_V6 at 1 range 6 .. 6;
    Slave_V7 at 1 range 7 .. 7;
end record;
```

Why are they equivalent? Well take a specific field, the Slave_V2 field. The storage place attributes are obtained by normalizing the values given so that the First_Bit value is less than 8. After normalizing the values (0,10,10) we get (1,2,2) which is exactly what we specified in the other case.
Now one might expect that the \texttt{Bit\_Order} attribute might affect bit numbering within the entire record component (two bytes in this case, thus affecting which byte fields end up in), but that is not the way this feature is defined, it only affects numbering of bits, not which byte they end up in.

Consequently it never makes sense to specify a starting bit number greater than 7 (for a byte addressable field) if an attribute definition for \texttt{Bit\_Order} has been given, and indeed it may be actively confusing to specify such a value, so the compiler generates a warning for such usage.

If you do need to control byte ordering then appropriate conditional values must be used. If in our example, the slave byte came first on some machines we might write:

\begin{verbatim}
Master_Byte_First constant Boolean := ...;

Master_Byte : constant Natural :=
    1 - Boolean'Pos (Master_Byte_First);
Slave_Byte : constant Natural :=
    Boolean'Pos (Master_Byte_First);

for Data’Bit_Order use High_Order_First;
for Data use record
    Master_Control at Master_Byte range 0 .. 0;
    Master_V1 at Master_Byte range 1 .. 1;
    Master_V2 at Master_Byte range 2 .. 2;
    Master_V3 at Master_Byte range 3 .. 3;
    Master_V4 at Master_Byte range 4 .. 4;
    Master_V5 at Master_Byte range 5 .. 5;
    Master_V6 at Master_Byte range 6 .. 6;
    Master_V7 at Master_Byte range 7 .. 7;
Slave_Control at Slave_Byte range 0 .. 0;
Slave_V1 at Slave_Byte range 1 .. 1;
Slave_V2 at Slave_Byte range 2 .. 2;
Slave_V3 at Slave_Byte range 3 .. 3;
Slave_V4 at Slave_Byte range 4 .. 4;
Slave_V5 at Slave_Byte range 5 .. 5;
Slave_V6 at Slave_Byte range 6 .. 6;
Slave_V7 at Slave_Byte range 7 .. 7;
end record;
\end{verbatim}

Now to switch between machines, all that is necessary is to set the boolean constant \texttt{Master\_Byte\_First} in an appropriate manner.

\subsection*{9.10 Pragma Pack for Arrays}

Pragma \texttt{Pack} applied to an array has an effect that depends upon whether the component type is ‘packable’. For a component type to be ‘packable’, it must be one of the following cases:

- Any elementary type.
- Any small packed array type with a static size.
* Any small simple record type with a static size.

For all these cases, if the component subtype size is in the range 1 through 63 on 32-bit targets, and 1 through 127 on 64-bit targets, then the effect of the pragma `Pack` is exactly as though a component size were specified giving the component subtype size.

All other types are non-packable, they occupy an integral number of storage units and the only effect of pragma `Pack` is to remove alignment gaps.

For example if we have:

```ada
type r is range 0 .. 17;

type ar is array (1 .. 8) of r;
pragma Pack (ar);
```

Then the component size of `ar` will be set to 5 (i.e., to `r'size`, and the size of the array `ar` will be exactly 40 bits).

Note that in some cases this rather fierce approach to packing can produce unexpected effects. For example, in Ada 95 and Ada 2005, subtype `Natural` typically has a size of 31, meaning that if you pack an array of `Natural`, you get 31-bit close packing, which saves a few bits, but results in far less efficient access. Since many other Ada compilers will ignore such a packing request, GNAT will generate a warning on some uses of pragma `Pack` that it guesses might not be what is intended. You can easily remove this warning by using an explicit `Component_Size` setting instead, which never generates a warning, since the intention of the programmer is clear in this case.

GNAT treats packed arrays in one of two ways. If the size of the array is known at compile time and is at most 64 bits on 32-bit targets, and at most 128 bits on 64-bit targets, then internally the array is represented as a single modular type, of exactly the appropriate number of bits. If the length is greater than 64 bits on 32-bit targets, and greater than 128 bits on 64-bit targets, or is not known at compile time, then the packed array is represented as an array of bytes, and its length is always a multiple of 8 bits.

Note that to represent a packed array as a modular type, the alignment must be suitable for the modular type involved. For example, on typical machines a 32-bit packed array will be represented by a 32-bit modular integer with an alignment of four bytes. If you explicitly override the default alignment with an alignment clause that is too small, the modular representation cannot be used. For example, consider the following set of declarations:

```ada
type R is range 1 .. 3;
type S is array (1 .. 31) of R;
for S'Component_Size use 2;
for S'Size use 62;
for S'Alignment use 1;
```

If the alignment clause were not present, then a 62-bit modular representation would be chosen (typically with an alignment of 4 or 8 bytes depending on the target). But the default alignment is overridden with the explicit alignment clause. This means that the modular representation cannot be used, and instead the array of bytes representation must be used, meaning that the length must be a multiple of 8. Thus the above set of declarations will result in a diagnostic rejecting the size clause and noting that the minimum size allowed is 64.
One special case that is worth noting occurs when the base type of the component size is 8/16/32 and the subtype is one bit less. Notably this occurs with subtype \texttt{Natural}. Consider:

\begin{verbatim}
  type Arr is array (1 .. 32) of Natural;
  pragma Pack (Arr);
\end{verbatim}

In all commonly used Ada 83 compilers, this pragma Pack would be ignored, since typically \texttt{Natural'Size} is 32 in Ada 83, and in any case most Ada 83 compilers did not attempt 31 bit packing.

In Ada 95 and Ada 2005, \texttt{Natural'Size} is required to be 31. Furthermore, GNAT really does pack 31-bit subtype to 31 bits. This may result in a substantial unintended performance penalty when porting legacy Ada 83 code. To help prevent this, GNAT generates a warning in such cases. If you really want 31 bit packing in a case like this, you can set the component size explicitly:

\begin{verbatim}
  type Arr is array (1 .. 32) of Natural;
  for Arr'Component_Size use 31;
\end{verbatim}

Here 31-bit packing is achieved as required, and no warning is generated, since in this case the programmer intention is clear.

\section*{9.11 Pragma Pack for Records}

Pragma \texttt{Pack} applied to a record will pack the components to reduce wasted space from alignment gaps and by reducing the amount of space taken by components. We distinguish between ‘packable’ components and ‘non-packable’ components. Components of the following types are considered packable:

\begin{itemize}
  \item Components of an elementary type are packable unless they are aliased, independent or atomic.
  \item Small packed arrays, where the size is statically known, are represented internally as modular integers, and so they are also packable.
  \item Small simple records, where the size is statically known, are also packable.
\end{itemize}

For all these cases, if the \texttt{'Size} value is in the range 1 through 64 on 32-bit targets, and 1 through 128 on 64-bit targets, the components occupy the exact number of bits corresponding to this value and are packed with no padding bits, i.e. they can start on an arbitrary bit boundary.

All other types are non-packable, they occupy an integral number of storage units and the only effect of pragma \texttt{Pack} is to remove alignment gaps.

For example, consider the record

\begin{verbatim}
  type Rb1 is array (1 .. 13) of Boolean;
  pragma Pack (Rb1);

  type Rb2 is array (1 .. 65) of Boolean;
  pragma Pack (Rb2);

  type AF is new Float with Atomic;
\end{verbatim}
type X2 is record
  L1 : Boolean;
  L2 : Duration;
  L3 : AF;
  L4 : Boolean;
  L5 : Rb1;
  L6 : Rb2;
end record;
pragma Pack (X2);

The representation for the record X2 is as follows on 32-bit targets:

for X2'Size use 224;
for X2 use record
  L1 at 0 range 0 .. 0;
  L2 at 0 range 1 .. 64;
  L3 at 12 range 0 .. 31;
  L4 at 16 range 0 .. 0;
  L5 at 16 range 1 .. 13;
  L6 at 18 range 0 .. 71;
end record;

Studying this example, we see that the packable fields L1 and L2 are of length equal to their sizes, and placed at specific bit boundaries (and not byte boundaries) to eliminate padding. But L3 is of a non-packable float type (because it is aliased), so it is on the next appropriate alignment boundary.

The next two fields are fully packable, so L4 and L5 are minimally packed with no gaps. However, type Rb2 is a packed array that is longer than 64 bits, so it is itself non-packable on 32-bit targets. Thus the L6 field is aligned to the next byte boundary, and takes an integral number of bytes, i.e., 72 bits.

9.12 Record Representation Clauses

Record representation clauses may be given for all record types, including types obtained by record extension. Component clauses are allowed for any static component. The restrictions on component clauses depend on the type of the component.

For all components of an elementary type, the only restriction on component clauses is that the size must be at least the 'Size value of the type (actually the Value_Size). There are no restrictions due to alignment, and such components may freely cross storage boundaries.

Packed arrays with a size up to and including 64 bits on 32-bit targets, and up to and including 128 bits on 64-bit targets, are represented internally using a modular type with the appropriate number of bits, and thus the same lack of restriction applies. For example, if you declare:

    type R is array (1 .. 49) of Boolean;
    pragma Pack (R);
    for R'Size use 49;

then a component clause for a component of type R may start on any specified bit boundary, and may specify a value of 49 bits or greater.
For packed bit arrays that are longer than 64 bits on 32-bit targets, and longer than 128 bits on 64-bit targets, there are two cases. If the component size is a power of 2 (1, 2, 4, 8, 16, 32, 64 bits), including the important case of single bits or boolean values, then there are no limitations on placement of such components, and they may start and end at arbitrary bit boundaries.

If the component size is not a power of 2 (e.g., 3 or 5), then an array of this type must always be placed on a storage unit (byte) boundary and occupy an integral number of storage units (bytes). Any component clause that does not meet this requirement will be rejected.

Any aliased component, or component of an aliased type, must have its normal alignment and size. A component clause that does not meet this requirement will be rejected.

The tag field of a tagged type always occupies an address sized field at the start of the record. No component clause may attempt to overlay this tag. When a tagged type appears as a component, the tag field must have proper alignment.

In the case of a record extension T1, of a type T, no component clause applied to the type T1 can specify a storage location that would overlap the first T\'Object_Size bits of the record.

For all other component types, including non-bit-packed arrays, the component can be placed at an arbitrary bit boundary, so for example, the following is permitted:

```plaintext
type R is array (1 .. 10) of Boolean;
for R\'Size use 80;

type Q is record
   G, H : Boolean;
   L, M : R;
end record;

for Q use record
   G at 0 range 0 .. 0;
   H at 0 range 1 .. 1;
   L at 0 range 2 .. 81;
   R at 0 range 82 .. 161;
end record;
```

### 9.13 Handling of Records with Holes

As a result of alignment considerations, records may contain “holes” or gaps which do not correspond to the data bits of any of the components. Record representation clauses can also result in holes in records.

GNAT does not attempt to clear these holes, so in record objects, they should be considered to hold undefined rubbish. The generated equality routine just tests components so does not access these undefined bits, and assignment and copy operations may or may not preserve the contents of these holes (for assignments, the holes in the target will in practice contain either the bits that are present in the holes in the source, or the bits that were present in the target before the assignment).
If it is necessary to ensure that holes in records have all zero bits, then record objects for which this initialization is desired should be explicitly set to all zero values using `Unchecked_Conversion` or address overlays. For example

```pascal
type HRec is record
    C : Character;
    I : Integer;
end record;
```

On typical machines, integers need to be aligned on a four-byte boundary, resulting in three bytes of undefined rubbish following the 8-bit field for C. To ensure that the hole in a variable of type HRec is set to all zero bits, you could for example do:

```pascal
type Base is record
    Dummy1, Dummy2 : Integer := 0;
end record;

BaseVar : Base;
RealVar : HRec;
for RealVar'Address use BaseVar'Address;
```

Now the 8-bytes of the value of RealVar start out containing all zero bits. A safer approach is to just define dummy fields, avoiding the holes, as in:

```pascal
type HRec is record
    C : Character;
    Dummy1 : Short_Short_Integer := 0;
    Dummy2 : Short_Short_Integer := 0;
    Dummy3 : Short_Short_Integer := 0;
    I : Integer;
end record;
```

And to make absolutely sure that the intent of this is followed, you can use representation clauses:

```pascal
for Hrec use record
    C at 0 range 0 .. 7;
    Dummy1 at 1 range 0 .. 7;
    Dummy2 at 2 range 0 .. 7;
    Dummy3 at 3 range 0 .. 7;
    I at 4 range 0 .. 31;
end record;
for Hrec'Size use 64;
```

### 9.14 Enumeration Clauses

The only restriction on enumeration clauses is that the range of values must be representable. For the signed case, if one or more of the representation values are negative, all values must be in the range:

```
System.Min_Int .. System.Max_Int
```

For the unsigned case, where all values are nonnegative, the values must be in the range:

```
0 .. System.Max_Binary_Modulus;
```
A ‘confirming’ representation clause is one in which the values range from 0 in sequence, i.e., a clause that confirms the default representation for an enumeration type. Such a confirming representation is permitted by these rules, and is specially recognized by the compiler so that no extra overhead results from the use of such a clause.

If an array has an index type which is an enumeration type to which an enumeration clause has been applied, then the array is stored in a compact manner. Consider the declarations:

```literate
type r is (A, B, C);
for r use (A => 1, B => 5, C => 10);
type t is array (r) of Character;
```

The array type t corresponds to a vector with exactly three elements and has a default size equal to $3 \times \text{Character'Size}$. This ensures efficient use of space, but means that accesses to elements of the array will incur the overhead of converting representation values to the corresponding positional values, (i.e., the value delivered by the `Pos` attribute).

### 9.15 Address Clauses

The reference manual allows a general restriction on representation clauses, as found in RM 13.1(22):

> "An implementation need not support representation items containing nonstatic expressions, except that an implementation should support a representation item for a given entity if each nonstatic expression in the representation item is a name that statically denotes a constant declared before the entity."

In practice this is applicable only to address clauses, since this is the only case in which a nonstatic expression is permitted by the syntax. As the AARM notes in sections 13.1 (22.a-22.h):

22.a Reason: This is to avoid the following sort of thing:
22.b X : Integer := F(...); Y : Address := G(...); for X'Address use Y;
22.c In the above, we have to evaluate the initialization expression for X before we know where to put the result. This seems like an unreasonable implementation burden.
22.d The above code should instead be written like this:
22.e Y : constant Address := G(...); X : Integer := F(...); for X'Address use Y;
22.f This allows the expression ‘Y’ to be safely evaluated before X is created.
22.g The constant could be a formal parameter of mode in.
22.h An implementation can support other nonstatic expressions if it wants to. Expressions of type Address are hardly ever static, but their value might be known at compile time anyway in many cases.

GNAT does indeed permit many additional cases of nonstatic expressions. In particular, if the type involved is elementary there are no restrictions (since in this case, holding a temporary copy of the initialization value, if one is present, is inexpensive). In addition, if there is no implicit or explicit initialization, then there are no restrictions. GNAT will reject only the case where all three of these conditions hold:
* The type of the item is non-elementary (e.g., a record or array).
* There is explicit or implicit initialization required for the object. Note that access values are always implicitly initialized.
* The address value is nonstatic. Here GNAT is more permissive than the RM, and allows the address value to be the address of a previously declared stand-alone variable, as long as it does not itself have an address clause.

```
Anchor : Some_Initialized_Type;
Overlay : Some_Initialized_Type;
for Overlay’Address use Anchor’Address;
```

However, the prefix of the address clause cannot be an array component, or a component of a discriminated record.

As noted above in section 22.h, address values are typically nonstatic. In particular the `To_Address` function, even if applied to a literal value, is a nonstatic function call. To avoid this minor annoyance, GNAT provides the implementation defined attribute `‘To_Address`. The following two expressions have identical values:

```
To_Address (16#1234_0000#)
System’To_Address (16#1234_0000#);
```

except that the second form is considered to be a static expression, and thus when used as an address clause value is always permitted.

Additionally, GNAT treats as static an address clause that is an unchecked_conversion of a static integer value. This simplifies the porting of legacy code, and provides a portable equivalent to the GNAT attribute `To_Address`.

Another issue with address clauses is the interaction with alignment requirements. When an address clause is given for an object, the address value must be consistent with the alignment of the object (which is usually the same as the alignment of the type of the object). If an address clause is given that specifies an inappropriately aligned address value, then the program execution is erroneous.

Since this source of erroneous behavior can have unfortunate effects on machines with strict alignment requirements, GNAT checks (at compile time if possible, generating a warning, or at execution time with a run-time check) that the alignment is appropriate. If the run-time check fails, then `Program_Error` is raised. This run-time check is suppressed if range checks are suppressed, or if the special GNAT check Alignment_Check is suppressed, or if `pragma Restrictions (No_Elaboration_Code)` is in effect. It is also suppressed by default on non-strict alignment machines (such as the x86).

In some cases, GNAT does not support an address specification (using either form of aspect specification syntax) for the declaration of an object that has an indefinite nominal subtype. An object declaration has an indefinite nominal subtype if it takes its bounds (for an array type), discriminant values (for a discriminated type whose discriminants lack defaults), or tag (for a class-wide type) from its initial value, as in

```
X : String := Some_Function_Call;
-- String has no constraint, so bounds for X come from function call
```

This restriction does not apply if the size of the object’s initial value is known at compile time and the type of the object is not class-wide.
An address clause cannot be given for an exported object. More understandably the real restriction is that objects with an address clause cannot be exported. This is because such variables are not defined by the Ada program, so there is no external object to export.

It is permissible to give an address clause and a pragma Import for the same object. In this case, the variable is not really defined by the Ada program, so there is no external symbol to be linked. The link name and the external name are ignored in this case. The reason that we allow this combination is that it provides a useful idiom to avoid unwanted initializations on objects with address clauses.

When an address clause is given for an object that has implicit or explicit initialization, then by default initialization takes place. This means that the effect of the object declaration is to overwrite the memory at the specified address. This is almost always not what the programmer wants, so GNAT will output a warning:

```
with System;
package G is
  type R is record
    M : Integer := 0;
  end record;

  Ext : R;
  for Ext'Address use System'To_Address (16#1234_1234#);

  >>> warning: implicit initialization of "Ext" may modify overlaid storage
  >>> warning: use pragma Import for "Ext" to suppress initialization (RM B(24))

end G;
```

As indicated by the warning message, the solution is to use a (dummy) pragma Import to suppress this initialization. The pragma tell the compiler that the object is declared and initialized elsewhere. The following package compiles without warnings (and the initialization is suppressed):

```
with System;
package G is
  type R is record
    M : Integer := 0;
  end record;

  Ext : R;
  for Ext'Address use System'To_Address (16#1234_1234#);
  pragma Import (Ada, Ext);
end G;
```

A final issue with address clauses involves their use for overlaying variables, as in the following example:

```
A : Integer;
B : Integer;
```
for B’Address use A’Address;

or alternatively, using the form recommended by the RM:

A : Integer;
Addr : constant Address := A’Address;
B : Integer;
for B’Address use Addr;

In both of these cases, A and B become aliased to one another via the address clause. This use of address clauses to overlay variables, achieving an effect similar to unchecked conversion was erroneous in Ada 83, but in Ada 95 and Ada 2005 the effect is implementation defined. Furthermore, the Ada RM specifically recommends that in a situation like this, B should be subject to the following implementation advice (RM 13.3(19)):

“19 If the Address of an object is specified, or it is imported or exported, then the implementation should not perform optimizations based on assumptions of no aliases.”

GNAT follows this recommendation, and goes further by also applying this recommendation to the overlaid variable (A in the above example) in this case. This means that the overlay works “as expected”, in that a modification to one of the variables will affect the value of the other.

More generally, GNAT interprets this recommendation conservatively for address clauses: in the cases other than overlays, it considers that the object is effectively subject to pragma Volatile and implements the associated semantics.

Note that when address clause overlays are used in this way, there is an issue of unintentional initialization, as shown by this example:

```ada
package Overwrite_Record is
  type R is record
     A : Character := 'C';
     B : Character := 'A';
  end record;
  X : Short_Integer := 3;
  Y : R;
  for Y’Address use X’Address;
end Overwrite_Record;
```

Here the default initialization of Y will clobber the value of X, which justifies the warning. The warning notes that this effect can be eliminated by adding a pragma Import which suppresses the initialization:

```ada
package Overwrite_Record is
  type R is record
     A : Character := 'C';
     B : Character := 'A';
  end record;
end Overwrite_Record;
```
X : Short_Integer := 3;
Y : R;
for Y’Address use X’Address;
pragma Import (Ada, Y);
end Overwrite_Record;

Note that the use of *pragma Initialize_Scalars* may cause variables to be initialized when they would not otherwise have been in the absence of the use of this pragma. This may cause an overlay to have this unintended clobbering effect. The compiler avoids this for scalar types, but not for composite objects (where in general the effect of *Initialize_Scalars* is part of the initialization routine for the composite object):

```
pragma Initialize_Scalars;
with Ada.Text_IO; use Ada.Text_IO;
procedure Overwrite_Array is
  type Arr is array (1 .. 5) of Integer;
  X : Arr := (others => 1);
  A : Arr;
  for A’Address use X’Address;
begin
  if X /= Arr’(others => 1) then
    Put_Line ("X was clobbered");
  else
    Put_Line ("X was not clobbered");
  end if;
end Overwrite_Array;
```

The above program generates the warning as shown, and at execution time, prints *X was clobbered*. If the *pragma Import* is added as suggested:

```
pragma Initialize_Scalars;
with Ada.Text_IO; use Ada.Text_IO;
procedure Overwrite_Array is
  type Arr is array (1 .. 5) of Integer;
  X : Arr := (others => 1);
  A : Arr;
  for A’Address use X’Address;
  pragma Import (Ada, A);
begin
  if X /= Arr’(others => 1) then
    Put_Line ("X was clobbered");
  else
    Put_Line ("X was not clobbered");
  end if;
end Overwrite_Array;
```
then the program compiles without the warning and when run will generate the output X was not clobbered.

9.16 Use of Address Clauses for Memory-Mapped I/O

A common pattern is to use an address clause to map an atomic variable to a location in memory that corresponds to a memory-mapped I/O operation or operations, for example:

```ada
type Mem_Word is record
  A,B,C,D : Byte;
end record;
pragma Atomic (Mem_Word);
for Mem_Word_Size use 32;

Mem : Mem_Word;
for Mem'Address use some-address;

Temp := Mem;
Temp.A := 32;
Mem := Temp;
```

For a full access (reference or modification) of the variable (Mem) in this case, as in the above examples, GNAT guarantees that the entire atomic word will be accessed, in accordance with the RM C.6(15) clause.

A problem arises with a component access such as:

```ada
Mem.A := 32;
```

Note that the component A is not declared as atomic. This means that it is not clear what this assignment means. It could correspond to full word read and write as given in the first example, or on architectures that supported such an operation it might be a single byte store instruction. The RM does not have anything to say in this situation, and GNAT does not make any guarantee. The code generated may vary from target to target. GNAT will issue a warning in such a case:

```ada
| >>> warning: access to non-atomic component of atomic array,
| may cause unexpected accesses to atomic object
```

It is best to be explicit in this situation, by either declaring the components to be atomic if you want the byte store, or explicitly writing the full word access sequence if that is what the hardware requires. Alternatively, if the full word access sequence is required, GNAT also provides the pragma **Volatile_Full_Access** which can be used in lieu of pragma **Atomic** and will give the additional guarantee.

9.17 Effect of Convention on Representation

Normally the specification of a foreign language convention for a type or an object has no effect on the chosen representation. In particular, the representation chosen for data in GNAT generally meets the standard system conventions, and for example records are laid
out in a manner that is consistent with C. This means that specifying convention C (for example) has no effect.

There are four exceptions to this general rule:

* ‘Convention Fortran and array subtypes’

  If pragma Convention Fortran is specified for an array subtype, then in accordance with the implementation advice in section 3.6.2(11) of the Ada Reference Manual, the array will be stored in a Fortran-compatible column-major manner, instead of the normal default row-major order.

* ‘Convention C and enumeration types’

  GNAT normally stores enumeration types in 8, 16, or 32 bits as required to accommodate all values of the type. For example, for the enumeration type declared by:

  ```ada
  type Color is (Red, Green, Blue);
  ```

  8 bits is sufficient to store all values of the type, so by default, objects of type Color will be represented using 8 bits. However, normal C convention is to use 32 bits for all enum values in C, since enum values are essentially of type int. If pragma Convention C is specified for an Ada enumeration type, then the size is modified as necessary (usually to 32 bits) to be consistent with the C convention for enum values.

  Note that this treatment applies only to types. If Convention C is given for an enumeration object, where the enumeration type is not Convention C, then Object_Size bits are allocated. For example, for a normal enumeration type, with less than 256 elements, only 8 bits will be allocated for the object. Since this may be a surprise in terms of what C expects, GNAT will issue a warning in this situation. The warning can be suppressed by giving an explicit size clause specifying the desired size.

* ‘Convention C/Fortran and Boolean types’

  In C, the usual convention for boolean values, that is values used for conditions, is that zero represents false, and nonzero values represent true. In Ada, the normal convention is that two specific values, typically 0/1, are used to represent false/true respectively.

  Fortran has a similar convention for LOGICAL values (any nonzero value represents true).

  To accommodate the Fortran and C conventions, if a pragma Convention specifies C or Fortran convention for a derived Boolean, as in the following example:

  ```ada
  type C_Switch is new Boolean;
  pragma Convention (C, C_Switch);
  ```

  then the GNAT generated code will treat any nonzero value as true. For truth values generated by GNAT, the conventional value 1 will be used for True, but when one of these values is read, any nonzero value is treated as True.

### 9.18 Conventions and Anonymous Access Types

The RM is not entirely clear on convention handling in a number of cases, and in particular, it is not clear on the convention to be given to anonymous access types in general, and in particular what is to be done for the case of anonymous access-to-subprogram.

In GNAT, we decide that if an explicit Convention is applied to an object or component, and its type is such an anonymous type, then the convention will apply to this anonymous
type as well. This seems to make sense since it is anomalous in any case to have a different
convention for an object and its type, and there is clearly no way to explicitly specify a
convention for an anonymous type, since it doesn’t have a name to specify!

Furthermore, we decide that if a convention is applied to a record type, then this convention
is inherited by any of its components that are of an anonymous access type which do not
have an explicitly specified convention.

The following program shows these conventions in action:

```ada
package ConvComp is
  type Foo is range 1 .. 10;
  type T1 is record
    A : access function (X : Foo) return Integer;
    B : Integer;
  end record;
  pragma Convention (C, T1);

  type T2 is record
    A : access function (X : Foo) return Integer;
    pragma Convention (C, A);
    B : Integer;
  end record;
  pragma Convention (COBOL, T2);

  type T3 is record
    A : access function (X : Foo) return Integer;
    pragma Convention (COBOL, A);
    B : Integer;
  end record;
  pragma Convention (C, T3);

  type T4 is record
    A : access function (X : Foo) return Integer;
    B : Integer;
  end record;
  pragma Convention (COBOL, T4);

  function F (X : Foo) return Integer;
  pragma Convention (C, F);

  function F (X : Foo) return Integer is (13);

  TV1 : T1 := (F'Access, 12);   -- OK
  TV2 : T2 := (F'Access, 13);   -- OK
  TV3 : T3 := (F'Access, 13);   -- ERROR
  | >>> subprogram "F" has wrong convention
```
9.19 Determining the Representations chosen by GNAT

Although the descriptions in this section are intended to be complete, it is often easier to simply experiment to see what GNAT accepts and what the effect is on the layout of types and objects.

As required by the Ada RM, if a representation clause is not accepted, then it must be rejected as illegal by the compiler. However, when a representation clause or pragma is accepted, there can still be questions of what the compiler actually does. For example, if a partial record representation clause specifies the location of some components and not others, then where are the non-specified components placed? Or if pragma Pack is used on a record, then exactly where are the resulting fields placed? The section on pragma Pack in this chapter can be used to answer the second question, but it is often easier to just see what the compiler does.

For this purpose, GNAT provides the option `-gnatR`. If you compile with this option, then the compiler will output information on the actual representations chosen, in a format similar to source representation clauses. For example, if we compile the package:

```ada
package q is
    type r (x : boolean) is tagged record
        case x is
            when True => S : String (1 .. 100);
            when False => null;
        end case;
    end record;

    type r2 is new r (false) with record
        y2 : integer;
    end record;

    for r2 use record
        y2 at 16 range 0 .. 31;
    end record;

    type x is record
        y : character;
    end record;

    type x1 is array (1 .. 10) of x;
    for x1'component_size use 11;

    type ia is access integer;

end q;
```
type Rb1 is array (1 .. 13) of Boolean;
pragma Pack (rb1);

type Rb2 is array (1 .. 65) of Boolean;
pragma Pack (rb2);

type x2 is record
  l1 : Boolean;
  l2 : Duration;
  l3 : Float;
  l4 : Boolean;
  l5 : Rb1;
  l6 : Rb2;
end record;
pragma Pack (x2);
end q;

using the switch `-gnatR` we obtain the following output:

Representation information for unit q
-------------------------------------
for r'Size use ??;
for r'Alignment use 4;
for r use record
  x at 4 range 0 .. 7;
  _tag at 0 range 0 .. 31;
  s at 5 range 0 .. 799;
end record;

for r2'Size use 160;
for r2'Alignment use 4;
for r2 use record
  x at 4 range 0 .. 7;
  _tag at 0 range 0 .. 31;
  _parent at 0 range 0 .. 63;
  y2 at 16 range 0 .. 31;
end record;

for x'Size use 8;
for x'Alignment use 1;
for x use record
  y at 0 range 0 .. 7;
end record;

for x1'Size use 112;
for x1'Alignment use 1;
for x1'Component_Size use 11;

for rb1'Size use 13;
for rb1'Alignment use 2;
for rb1'Component_Size use 1;

for rb2'Size use 72;
for rb2'Alignment use 1;
for rb2'Component_Size use 1;

for x2'Size use 224;
for x2'Alignment use 4;
for x2 use record
  11 at 0 range 0 .. 0;
  12 at 0 range 1 .. 64;
  13 at 12 range 0 .. 31;
  14 at 16 range 0 .. 0;
  15 at 16 range 1 .. 13;
  16 at 18 range 0 .. 71;
end record;

The Size values are actually the Object_Size, i.e., the default size that will be allocated for objects of the type. The ?? size for type r indicates that we have a variant record, and the actual size of objects will depend on the discriminant value.

The Alignment values show the actual alignment chosen by the compiler for each record or array type.

The record representation clause for type r shows where all fields are placed, including the compiler generated tag field (whose location cannot be controlled by the programmer).

The record representation clause for the type extension r2 shows all the fields present, including the parent field, which is a copy of the fields of the parent type of r2, i.e., r1.

The component size and size clauses for types rb1 and rb2 show the exact effect of pragma Pack on these arrays, and the record representation clause for type x2 shows how pragma Pack affects this record type.

In some cases, it may be useful to cut and paste the representation clauses generated by the compiler into the original source to fix and guarantee the actual representation to be used.
10 Standard Library Routines

The Ada Reference Manual contains in Annex A a full description of an extensive set of standard library routines that can be used in any Ada program, and which must be provided by all Ada compilers. They are analogous to the standard C library used by C programs. GNAT implements all of the facilities described in annex A, and for most purposes the description in the Ada Reference Manual, or appropriate Ada text book, will be sufficient for making use of these facilities.

In the case of the input-output facilities, [The Implementation of Standard I/O], page 236, gives details on exactly how GNAT interfaces to the file system. For the remaining packages, the Ada Reference Manual should be sufficient. The following is a list of the packages included, together with a brief description of the functionality that is provided.

For completeness, references are included to other predefined library routines defined in other sections of the Ada Reference Manual (these are cross-indexed from Annex A). For further details see the relevant package declarations in the run-time library. In particular, a few units are not implemented, as marked by the presence of pragma Unimplemented_Unit, and in this case the package declaration contains comments explaining why the unit is not implemented.

Ada ‘(A.2)’
This is a parent package for all the standard library packages. It is usually included implicitly in your program, and itself contains no useful data or routines.

Ada.Assertions ‘(11.4.2)’
Assertions provides the Assert subprograms, and also the declaration of the Assertion_Error exception.

Ada.Asynchronous_Task_Control ‘(D.11)’
Asynchronous_Task_Control provides low level facilities for task synchronization. It is typically not implemented. See package spec for details.

Ada.Calendar ‘(9.6)’
Calendar provides time of day access, and routines for manipulating times and durations.

Ada.Calendar.Arithmetic ‘(9.6.1)’
This package provides additional arithmetic operations for Calendar.

Ada.Calendar.Formatting ‘(9.6.1)’
This package provides formatting operations for Calendar.

Ada.Calendar.Time_Zones ‘(9.6.1)’
This package provides additional Calendar facilities for handling time zones.

Ada.Characters ‘(A.3.1)’
This is a dummy parent package that contains no useful entities.

Ada.Characters.Conversions ‘(A.3.2)’
This package provides character conversion functions.
Ada.Characters.Handling ‘(A.3.2)’
This package provides some basic character handling capabilities, including
classification functions for classes of characters (e.g., test for letters, or digits).

Ada.Characters.Latin_1 ‘(A.3.3)’
This package includes a complete set of definitions of the characters that ap-
pear in type CHARACTER. It is useful for writing programs that will run in
international environments. For example, if you want an upper case E with an
acute accent in a string, it is often better to use the definition of UC_E_Acute in
this package. Then your program will print in an understandable manner even
if your environment does not support these extended characters.

Ada.Command_Line ‘(A.15)’
This package provides access to the command line parameters and the name
of the current program (analogous to the use of argc and argv in C), and
also allows the exit status for the program to be set in a system-independent
manner.

Ada.Complex_Text_IO ‘(G.1.3)’
This package provides text input and output of complex numbers.

Ada.Containers ‘(A.18.1)’
A top level package providing a few basic definitions used by all the following
specific child packages that provide specific kinds of containers.

Ada.Containers.Bounded_Synchronized_ Queues ‘(A.18.29)’
Ada.Containers.Doubly_Linked_Lists ‘(A.18.3)’
Ada.Containers.Generic_Array_Sort ‘(A.18.26)’
Ada.Containers.Generic_Constrained_Array_Sort ‘(A.18.26)’
Ada.Containers.Hashed_Maps ‘(A.18.5)’
Ada.Containers.Hashed_Sets ‘(A.18.8)’
Ada.Containers.Indefinite_Hashed_Sets ‘(A.18.15)’
Ada.Containers.Indefinite_Holders ‘(A.18.18)’
Ada.Containers.Indefinite_Ordered_Sets ‘(A.18.16)’
Ada.Containers.Indefinite_Vectors ‘(A.18.11)’
Ada.Containers.Multiway_Trees ‘(A.18.10)’
Ada.Containers.Ordered_Maps ‘(A.18.6)’
Ada.Containers.Ordered_Sets ‘(A.18.9)’
Ada.Containers.Synchronized_Queue_Interfaces ‘(A.18.27)’
This package provides operations on directories.

This package provides additional directory operations handling hierarchical file names.

This is an implementation defined package for additional directory operations, which is not implemented in GNAT.

This package provides constants describing the range of decimal numbers implemented, and also a decimal divide routine (analogous to the COBOL verb DIVIDE . . . GIVING . . . REMAINDER . . .)

This package provides input-output using a model of a set of records of fixed-length, containing an arbitrary definite Ada type, indexed by an integer record number.

A parent package containing definitions for task dispatching operations.

Not implemented in GNAT.

Not implemented in GNAT.

Not implemented in GNAT.

This package allows the priorities of a task to be adjusted dynamically as the task is running.

This package provides facilities for accessing environment variables.

This package provides additional information on exceptions, and also contains facilities for treating exceptions as data objects, and raising exceptions with associated messages.

This package provides CPU clock functionalities. It is not implemented on all targets (see package spec for details).

Not implemented in GNAT.
Ada.Execution_Time.Timers ‘(D.14.1)’
Not implemented in GNAT.

Ada.Finalization ‘(7.6)’
This package contains the declarations and subprograms to support the use of
controlled types, providing for automatic initialization and finalization (analogous
to the constructors and destructors of C++).

Ada.Float_Text_IO ‘(A.10.9)’
A library level instantiation of Text_IO.Float_IO for type Float.

Ada.Float_Wide_Text_IO ‘(A.10.9)’
A library level instantiation of Wide_Text_IO.Float_IO for type Float.

Ada.Float_Wide_Wide_Text_IO ‘(A.10.9)’
A library level instantiation of Wide_Wide_Text_IO.Float_IO for type Float.

Ada.Integer_Text_IO ‘(A.10.9)’
A library level instantiation of Text_IO.Integer_IO for type Integer.

Ada.Integer_Wide_Text_IO ‘(A.10.9)’
A library level instantiation of Wide_Text_IO.Integer_IO for type Integer.

Ada.Integer_Wide_Wide_Text_IO ‘(A.10.9)’
A library level instantiation of Wide_Wide_Text_IO.Integer_IO for type Integer.

Ada.Interrupts ‘(C.3.2)’
This package provides facilities for interfacing to interrupts, which includes the
set of signals or conditions that can be raised and recognized as interrupts.

Ada.Interrupts.Names ‘(C.3.2)’
This package provides the set of interrupt names (actually signal or condition
names) that can be handled by GNAT.

Ada.IO_Exceptions ‘(A.13)’
This package defines the set of exceptions that can be raised by use of the
standard IO packages.

Ada.Iterator_Interfaces ‘(5.5.1)’
This package provides a generic interface to generalized iterators.

Ada.Locales ‘(A.19)’
This package provides declarations providing information (Language and Coun-
try) about the current locale.

Ada.Numerics
This package contains some standard constants and exceptions used throughout
the numerics packages. Note that the constants pi and e are defined here, and
it is better to use these definitions than rolling your own.

Ada.Numerics.Complex_Arrays ‘(G.3.2)’
Provides operations on arrays of complex numbers.

Ada.Numerics.Complex_Elementary_Functions
Provides the implementation of standard elementary functions (such as log
and trigonometric functions) operating on complex numbers using the stan-
standard Float and the Complex and Imaginary types created by the package Numerics.Complex_Types.

Ada.Numerics.Complex_Types
This is a predefined instantiation of Numerics.Generic_Complex_Types using Standard.Float to build the type Complex and Imaginary.

Ada.Numerics.Discrete_Random
This generic package provides a random number generator suitable for generating uniformly distributed values of a specified discrete subtype.

Ada.Numerics.Float_Random
This package provides a random number generator suitable for generating uniformly distributed floating point values in the unit interval.

Ada.Numerics.Generic_Complex_Elementary_Functions
This is a generic version of the package that provides the implementation of standard elementary functions (such as log and trigonometric functions) for an arbitrary complex type.

The following predefined instantiations of this package are provided:
  * Short_Float
    Ada.Numerics.Short_Complex_Elementary_Functions
  * Float
    Ada.Numerics.Complex_Elementary_Functions
  * Long_Float
    Ada.Numerics.Long_Complex_Elementary_Functions

Ada.Numerics.Generic_Complex_Types
This is a generic package that allows the creation of complex types, with associated complex arithmetic operations.

The following predefined instantiations of this package exist
  * Short_Float
    Ada.Numerics.Short_Complex_Complex_Types
  * Float
    Ada.Numerics.Complex_Complex_Types
  * Long_Float
    Ada.Numerics.Long_Complex_Complex_Types

Ada.Numerics.Generic_Elementary_Functions
This is a generic package that provides the implementation of standard elementary functions (such as log and trigonometric functions) for an arbitrary float type.

The following predefined instantiations of this package exist
  * Short_Float
    Ada.Numerics.Short_Elementary_Functions
  * Float
    Ada.Numerics.Elementary_Functions
Chapter 10: Standard Library Routines

* Long_Float
  Ada.Numerics.Long_Elementary_Functions

Ada.Numerics.Generic_Real_Arrays ‘(G.3.1)’
  Generic operations on arrays of reals

Ada.Numerics.Real_Arrays ‘(G.3.1)’
  Preinstantiation of Ada.Numerics.Generic_Real_Arrays (Float).

Ada.Real_Time ‘(D.8)’
  This package provides facilities similar to those of Calendar, but operating with
  a finer clock suitable for real time control. Note that annex D requires that there
  be no backward clock jumps, and GNAT generally guarantees this behavior,
  but of course if the external clock on which the GNAT runtime depends is
  deliberately reset by some external event, then such a backward jump may
  occur.

Ada.Real_Time.Timing_Events ‘(D.15)’
  Not implemented in GNAT.

Ada.Sequential_IO ‘(A.8.1)’
  This package provides input-output facilities for sequential files, which can con-
  tain a sequence of values of a single type, which can be any Ada type, including
  indefinite (unconstrained) types.

Ada.Storage_IO ‘(A.9)’
  This package provides a facility for mapping arbitrary Ada types to and from
  a storage buffer. It is primarily intended for the creation of new IO packages.

Ada.Strings ‘(A.4.1)’
  This package provides some basic constants used by the string handling pack-
  ages.

Ada.Strings.Bounded ‘(A.4.4)’
  This package provides facilities for handling variable length strings. The
  bounded model requires a maximum length. It is thus somewhat more limited
  than the unbounded model, but avoids the use of dynamic allocation or
  finalization.

Ada.Strings.Bounded.Equal_Case_Insensitive ‘(A.4.10)’
  Provides case-insensitive comparisons of bounded strings

  This package provides a generic hash function for bounded strings
This package provides a generic hash function for bounded strings that converts
the string to be hashed to lower case.

Ada.Strings.Bounded.Less_Case_Insensitive ‘(A.4.10)’
This package provides a comparison function for bounded strings that works in
a case insensitive manner by converting to lower case before the comparison.

Ada.Strings.Fixed ‘(A.4.3)’
This package provides facilities for handling fixed length strings.

Ada.Strings.Fixed.Equal_Case_Insensitive ‘(A.4.10)’
This package provides an equality function for fixed strings that compares the
strings after converting both to lower case.

Ada.Strings.Fixed.Hash_Case_Insensitive ‘(A.4.9)’
This package provides a case insensitive hash function for fixed strings that
converts the string to lower case before computing the hash.

Ada.Strings.Fixed.Less_Case_Insensitive ‘(A.4.10)’
This package provides a comparison function for fixed strings that works in a
case insensitive manner by converting to lower case before the comparison.

Ada.Strings.Hash ‘(A.4.9)’
This package provides a hash function for strings.

Ada.Strings.Hash_Case_Insensitive ‘(A.4.9)’
This package provides a hash function for strings that is case insensitive. The
string is converted to lower case before computing the hash.

Ada.Strings.Less_Case_Insensitive ‘(A.4.10)’
This package provides a comparison function for strings that works in a case
insensitive manner by converting to lower case before the comparison.

Ada.Strings.Maps ‘(A.4.2)’
This package provides facilities for handling character mappings and arbitrarily
defined subsets of characters. For instance it is useful in defining specialized
translation tables.

Ada.Strings.Maps.Constants ‘(A.4.6)’
This package provides a standard set of predefined mappings and predefined
character sets. For example, the standard upper to lower case conversion table
is found in this package. Note that upper to lower case conversion is non-trivial
if you want to take the entire set of characters, including extended characters
like E with an acute accent, into account. You should use the mappings in this
package (rather than adding 32 yourself) to do case mappings.

Ada.Strings.Unbounded ‘(A.4.5)’
This package provides facilities for handling variable length strings. The un-
bounded model allows arbitrary length strings, but requires the use of dynamic
allocation and finalization.

Ada.Strings.Unbounded.Equal_Case_Insensitive ‘(A.4.10)’
Provides case-insensitive comparisons of unbounded strings
Ada.Strings.Unbounded.Hash '(A.4.9)'
This package provides a generic hash function for unbounded strings

Ada.Strings.Unbounded.Hash_Case_Insensitive '(A.4.9)'
This package provides a generic hash function for unbounded strings that converts the string to be hashed to lower case.

Ada.Strings.Unbounded.Less_Case_Insensitive '(A.4.10)'
This package provides a comparison function for unbounded strings that works in a case insensitive manner by converting to lower case before the comparison.

Ada.Strings.UTF_Encoding '(A.4.11)'
This package provides basic definitions for dealing with UTF-encoded strings.

Ada.Strings.UTF_Encoding.Conversions '(A.4.11)'
This package provides conversion functions for UTF-encoded strings.

Ada.Strings.UTF_Encoding.Strings '(A.4.11)'
Ada.Strings.UTF_Encoding.Wide_Strings '(A.4.11)'
Ada.Strings.UTF_Encoding.Wide_Wide_Strings '(A.4.11)'
These packages provide facilities for handling UTF encodings for Strings, Wide_Strings and Wide_Wide_Strings.

Ada.Strings.Wide_Bounded '(A.4.7)'
Ada.Strings.Wide_Fixed '(A.4.7)'
Ada.Strings.Wide_Maps '(A.4.7)'
Ada.Strings.Wide_Unbounded '(A.4.7)'
These packages provide analogous capabilities to the corresponding packages without Wide_ in the name, but operate with the types Wide_String and Wide_Character instead of String and Character. Versions of all the child packages are available.

Ada.Strings.Wide_Wide_Bounded '(A.4.7)'
Ada.Strings.Wide_Wide_Fixed '(A.4.7)'
Ada.Strings.Wide_Wide_Maps '(A.4.7)'
Ada.Strings.Wide_Wide_Unbounded '(A.4.7)'
These packages provide analogous capabilities to the corresponding packages without Wide_ in the name, but operate with the types Wide_Wide_String and Wide_Wide_Character instead of String and Character.

Ada.Synchronous_Barriers '(D.10.1)'
This package provides facilities for synchronizing tasks at a low level with barriers.

Ada.Synchronous_Task_Control '(D.10)'
This package provides some standard facilities for controlling task communication in a synchronous manner.

Ada.Synchronous_Task_Control.EDF '(D.10)'
Not implemented in GNAT.
Ada.Tags
This package contains definitions for manipulation of the tags of tagged values.

Ada.Tags.Generic_Dispatching_Constructor ‘(3.9)’
This package provides a way of constructing tagged class-wide values given only the tag value.

Ada.Task_Attributes ‘(C.7.2)’
This package provides the capability of associating arbitrary task-specific data with separate tasks.

Ada.Task_Identification ‘(C.7.1)’
This package provides capabilities for task identification.

Ada.Task_Termination ‘(C.7.3)’
This package provides control over task termination.

Ada.Text_IO
This package provides basic text input-output capabilities for character, string and numeric data. The subpackages of this package are listed next. Note that although these are defined as subpackages in the RM, they are actually transparently implemented as child packages in GNAT, meaning that they are only loaded if needed.

Ada.Text_IO.Decimal_IO
Provides input-output facilities for decimal fixed-point types

Ada.Text_IO.Enumeration_IO
Provides input-output facilities for enumeration types.

Ada.Text_IO.Fixed_IO
Provides input-output facilities for ordinary fixed-point types.

Ada.Text_IO.Float_IO
Provides input-output facilities for float types. The following predefined instantiations of this generic package are available:

* Short_Float
  Short_Float_Text_IO
* Float
  Float_Text_IO
* Long_Float
  Long_Float_Text_IO

Ada.Text_IO.Integer_IO
Provides input-output facilities for integer types. The following predefined instantiations of this generic package are available:

* Short_Short_Integer
  Ada.Short_Short_Integer_Text_IO
* Short_Integer
  Ada.Short_Integer_Text_IO
* Integer
  Ada.Integer_Text_IO
* Long_Integer
  Ada.Long_Integer_Text_IO
* Long_Long_Integer
  Ada.Long_Long_Integer_Text_IO

Ada.Text_IO.Modular_IO
  Provides input-output facilities for modular (unsigned) types.

Ada.Text_IO.Bounded_IO (A.10.11)
  Provides input-output facilities for bounded strings.

Ada.Text_IO.Complex_IO (G.1.3)
  This package provides basic text input-output capabilities for complex data.

Ada.Text_IO.Editing (F.3.3)
  This package contains routines for edited output, analogous to the use of pictures in COBOL. The picture formats used by this package are a close copy of the facility in COBOL.

Ada.Text_IO.Text_Streams (A.12.2)
  This package provides a facility that allows Text_IO files to be treated as streams, so that the stream attributes can be used for writing arbitrary data, including binary data, to Text_IO files.

Ada.Text_IO.Unbounded_IO (A.10.12)
  This package provides input-output facilities for unbounded strings.

Ada.Unchecked_Conversion (13.9)
  This generic package allows arbitrary conversion from one type to another of the same size, providing for breaking the type safety in special circumstances.

  If the types have the same Size (more accurately the same Value_Size), then the effect is simply to transfer the bits from the source to the target type without any modification. This usage is well defined, and for simple types whose representation is typically the same across all implementations, gives a portable method of performing such conversions.

  If the types do not have the same size, then the result is implementation defined, and thus may be non-portable. The following describes how GNAT handles such unchecked conversion cases.

  If the types are of different sizes, and are both discrete types, then the effect is of a normal type conversion without any constraint checking. In particular if the result type has a larger size, the result will be zero or sign extended. If the result type has a smaller size, the result will be truncated by ignoring high order bits.

  If the types are of different sizes, and are not both discrete types, then the conversion works as though pointers were created to the source and target, and the pointer value is converted. The effect is that bits are copied from successive
low order storage units and bits of the source up to the length of the target type.
A warning is issued if the lengths differ, since the effect in this case is implementation dependent, and the above behavior may not match that of some other compiler.
A pointer to one type may be converted to a pointer to another type using unchecked conversion. The only case in which the effect is undefined is when one or both pointers are pointers to unconstrained array types. In this case, the bounds information may get incorrectly transferred, and in particular, GNAT uses double size pointers for such types, and it is meaningless to convert between such pointer types. GNAT will issue a warning if the alignment of the target designated type is more strict than the alignment of the source designated type (since the result may be unaligned in this case).
A pointer other than a pointer to an unconstrained array type may be converted to and from System.Address. Such usage is common in Ada 83 programs, but note that Ada.Address_To_Access_Conversions is the preferred method of performing such conversions in Ada 95 and Ada 2005. Neither unchecked conversion nor Ada.Address_To_Access_Conversions should be used in conjunction with pointers to unconstrained objects, since the bounds information cannot be handled correctly in this case.

Ada.Unchecked_Deallocation (13.11.2)
This generic package allows explicit freeing of storage previously allocated by use of an allocator.

Ada.Wide_Text_IO (A.11)
This package is similar to Ada.Text_IO, except that the external file supports wide character representations, and the internal types are Wide_Character and Wide_String instead of Character and String. The corresponding set of nested packages and child packages are defined.

Ada.Wide_Wide_Text_IO (A.11)
This package is similar to Ada.Text_IO, except that the external file supports wide character representations, and the internal types are Wide_Character and Wide_String instead of Character and String. The corresponding set of nested packages and child packages are defined.

For packages in Interfaces and System, all the RM defined packages are available in GNAT, see the Ada 2012 RM for full details.
11 The Implementation of Standard I/O

GNAT implements all the required input-output facilities described in A.6 through A.14. These sections of the Ada Reference Manual describe the required behavior of these packages from the Ada point of view, and if you are writing a portable Ada program that does not need to know the exact manner in which Ada maps to the outside world when it comes to reading or writing external files, then you do not need to read this chapter. As long as your files are all regular files (not pipes or devices), and as long as you write and read the files only from Ada, the description in the Ada Reference Manual is sufficient.

However, if you want to do input-output to pipes or other devices, such as the keyboard or screen, or if the files you are dealing with are either generated by some other language, or to be read by some other language, then you need to know more about the details of how the GNAT implementation of these input-output facilities behaves.

In this chapter we give a detailed description of exactly how GNAT interfaces to the file system. As always, the sources of the system are available to you for answering questions at an even more detailed level, but for most purposes the information in this chapter will suffice.

Another reason that you may need to know more about how input-output is implemented arises when you have a program written in mixed languages where, for example, files are shared between the C and Ada sections of the same program. GNAT provides some additional facilities, in the form of additional child library packages, that facilitate this sharing, and these additional facilities are also described in this chapter.

11.1 Standard I/O Packages

The Standard I/O packages described in Annex A for

* Ada.Text_IO
* Ada.Text_IO.Complex_IO
* Ada.Text_IO.Text_Streams
* Ada.Wide_Text_IO
* Ada.Wide_Text_IO.Complex_IO
* Ada.Wide_Text_IO.Text_Streams
* Ada.Wide_Wide_Text_IO
* Ada.Wide_Wide_Text_IO.Complex_IO
* Ada.Wide_Wide_Text_IO.Text_Streams
* Ada.Stream_IO
* Ada.Sequential_IO
* Ada.Direct_IO

are implemented using the C library streams facility; where

* All files are opened using `fopen`.
* All input/output operations use `fread/fwrite`.
There is no internal buffering of any kind at the Ada library level. The only buffering is that provided at the system level in the implementation of the library routines that support streams. This facilitates shared use of these streams by mixed language programs. Note though that system level buffering is explicitly enabled at elaboration of the standard I/O packages and that can have an impact on mixed language programs, in particular those using I/O before calling the Ada elaboration routine (e.g., adainit). It is recommended to call the Ada elaboration routine before performing any I/O or when impractical, flush the common I/O streams and in particular Standard_Output before elaborating the Ada code.

11.2 FORM Strings

The format of a FORM string in GNAT is:

"keyword=value, keyword=value, ..., keyword=value"

where letters may be in upper or lower case, and there are no spaces between values. The order of the entries is not important. Currently the following keywords defined.

- TEXT_TRANSLATION=[YES|NO|TEXT|BINARY|U8TEXT|WTEXT|U16TEXT]
- SHARED=[YES|NO]
- WCEM=[n|h|u|s|e|8|b]
- ENCODING=[UTF8|8BITS]

The use of these parameters is described later in this section. If an unrecognized keyword appears in a form string, it is silently ignored and not considered invalid.

11.3 Direct_IO

Direct_IO can only be instantiated for definite types. This is a restriction of the Ada language, which means that the records are fixed length (the length being determined by type'Size, rounded up to the next storage unit boundary if necessary).

The records of a Direct_IO file are simply written to the file in index sequence, with the first record starting at offset zero, and subsequent records following. There is no control information of any kind. For example, if 32-bit integers are being written, each record takes 4-bytes, so the record at index K starts at offset (K-1)*4.

There is no limit on the size of Direct_IO files, they are expanded as necessary to accommodate whatever records are written to the file.

11.4 Sequential_IO

Sequential_IO may be instantiated with either a definite (constrained) or indefinite (unconstrained) type.

For the definite type case, the elements written to the file are simply the memory images of the data values with no control information of any kind. The resulting file should be read using the same type, no validity checking is performed on input.

For the indefinite type case, the elements written consist of two parts. First is the size of the data item, written as the memory image of a Interfaces.C.size_t value, followed by the memory image of the data value. The resulting file can only be read using the same (unconstrained) type. Normal assignment checks are performed on these read operations, and if these checks fail, Data_Error is raised. In particular, in the array case, the lengths
must match, and in the variant record case, if the variable for a particular read operation is constrained, the discriminants must match.

Note that it is not possible to use Sequential_IO to write variable length array items, and then read the data back into different length arrays. For example, the following will raise Data_Error:

```ada
package IO is new Sequential_IO (String);
F : IO.File_Type;
S : String (1..4);
...
IO.Create (F)
IO.Write (F, "hello!")
IO.Reset (F, Mode=>In_File);
IO.Read (F, S);
Put_Line (S);
```

On some Ada implementations, this will print hell, but the program is clearly incorrect, since there is only one element in the file, and that element is the string hello!.

In Ada 95 and Ada 2005, this kind of behavior can be legitimately achieved using Stream_IO, and this is the preferred mechanism. In particular, the above program fragment rewritten to use Stream_IO will work correctly.

### 11.5 Text_IO

Text_IO files consist of a stream of characters containing the following special control characters:

- **LF** (line feed, 16#0A#) Line Mark
- **FF** (form feed, 16#0C#) Page Mark

A canonical Text_IO file is defined as one in which the following conditions are met:

* The character LF is used only as a line mark, i.e., to mark the end of the line.
* The character FF is used only as a page mark, i.e., to mark the end of a page and consequently can appear only immediately following a LF (line mark) character.
* The file ends with either LF (line mark) or LF-FF (line mark, page mark). In the former case, the page mark is implicitly assumed to be present.

A file written using Text_IO will be in canonical form provided that no explicit LF or FF characters are written using Put or Put_Line. There will be no FF character at the end of the file unless an explicit New_Page operation was performed before closing the file.

A canonical Text_IO file that is a regular file (i.e., not a device or a pipe) can be read using any of the routines in Text_IO. The semantics in this case will be exactly as defined in the Ada Reference Manual, and all the routines in Text_IO are fully implemented.

A text file that does not meet the requirements for a canonical Text_IO file has one of the following:

* The file contains FF characters not immediately following a LF character.
* The file contains LF or FF characters written by Put or Put_Line, which are not logically considered to be line marks or page marks.
* The file ends in a character other than LF or FF, i.e., there is no explicit line mark or page mark at the end of the file.

Text IO can be used to read such non-standard text files but subprograms to do with line or page numbers do not have defined meanings. In particular, a FF character that does not follow a LF character may or may not be treated as a page mark from the point of view of page and line numbering. Every LF character is considered to end a line, and there is an implied LF character at the end of the file.

### 11.5.1 Stream Pointer Positioning

Ada.Text_IO has a definition of current position for a file that is being read. No internal buffering occurs in Text_IO, and usually the physical position in the stream used to implement the file corresponds to this logical position defined by Text_IO. There are two exceptions:

* After a call to End_Of_Page that returns True, the stream is positioned past the LF (line mark) that precedes the page mark. Text_IO maintains an internal flag so that subsequent read operations properly handle the logical position which is unchanged by the End_Of_Page call.

* After a call to End_Of_File that returns True, if the Text_IO file was positioned before the line mark at the end of file before the call, then the logical position is unchanged, but the stream is physically positioned right at the end of file (past the line mark, and past a possible page mark following the line mark. Again Text_IO maintains internal flags so that subsequent read operations properly handle the logical position.

These discrepancies have no effect on the observable behavior of Text_IO, but if a single Ada stream is shared between a C program and Ada program, or shared (using shared=yes in the form string) between two Ada files, then the difference may be observable in some situations.

### 11.5.2 Reading and Writing Non-Regular Files

A non-regular file is a device (such as a keyboard), or a pipe. Text_IO can be used for reading and writing. Writing is not affected and the sequence of characters output is identical to the normal file case, but for reading, the behavior of Text_IO is modified to avoid undesirable look-ahead as follows:

An input file that is not a regular file is considered to have no page marks. Any Ascii.FF characters (the character normally used for a page mark) appearing in the file are considered to be data characters. In particular:

* Get_Line and Skip_Line do not test for a page mark following a line mark. If a page mark appears, it will be treated as a data character.

* This avoids the need to wait for an extra character to be typed or entered from the pipe to complete one of these operations.

* End_Of_Page always returns False

* End_Of_File will return False if there is a page mark at the end of the file.

Output to non-regular files is the same as for regular files. Page marks may be written to non-regular files using New_Page, but as noted above they will not be treated as page marks on input if the output is piped to another Ada program.
Another important discrepancy when reading non-regular files is that the end of file indication is not ‘sticky’. If an end of file is entered, e.g., by pressing the EOT key, then end of file is signaled once (i.e., the test End_Of_File will yield True, or a read will raise End_Error), but then reading can resume to read data past that end of file indication, until another end of file indication is entered.

### 11.5.3 Get_Immediate

Get_Immediate returns the next character (including control characters) from the input file. In particular, Get_Immediate will return LF or FF characters used as line marks or page marks. Such operations leave the file positioned past the control character, and it is thus not treated as having its normal function. This means that page, line and column counts after this kind of Get_Immediate call are set as though the mark did not occur. In the case where a Get_Immediate leaves the file positioned between the line mark and page mark (which is not normally possible), it is undefined whether the FF character will be treated as a page mark.

### 11.5.4 Treating Text_IO Files as Streams

The package Text_IO.Streams allows a Text_IO file to be treated as a stream. Data written to a Text_IO file in this stream mode is binary data. If this binary data contains bytes 16#0A# (LF) or 16#0C# (FF), the resulting file may have non-standard format. Similarly if read operations are used to read from a Text_IO file treated as a stream, then LF and FF characters may be skipped and the effect is similar to that described above for Get_Immediate.

### 11.5.5 Text_IO Extensions

A package GNAT.IO.Aux in the GNAT library provides some useful extensions to the standard Text_IO package:

* function File_Exists (Name : String) return Boolean; Determines if a file of the given name exists.

* function Get_Line return String; Reads a string from the standard input file. The value returned is exactly the length of the line that was read.

* function Get_Line (File : Ada.Text_IO.File_Type) return String; Similar, except that the parameter File specifies the file from which the string is to be read.

### 11.5.6 Text_IO Facilities for Unbounded Strings

The package Ada.Strings.Unbounded.Text_IO in library files a-suteio.ads/adb contains some GNAT-specific subprograms useful for Text_IO operations on unbounded strings:

* function Get_Line (File : File_Type) return Unbounded_String; Reads a line from the specified file and returns the result as an unbounded string.

* procedure Put (File : File_Type; U : Unbounded_String); Writes the value of the given unbounded string to the specified file. Similar to the effect of Put (To_String (U)) except that an extra copy is avoided.

* procedure Put_Line (File : File_Type; U : Unbounded_String); Writes the value of the given unbounded string to the specified file, followed by a New_Line. Similar to the effect of Put_Line (To_String (U)) except that an extra copy is avoided.
In the above procedures, \texttt{File} is of type \texttt{Ada.Text_IO.File_Type} and is optional. If the parameter is omitted, then the standard input or output file is referenced as appropriate.

The package \texttt{Ada.Strings.Wide_Unbounded.Wide_Text_IO} in library files \texttt{a-swuwti.ads} and \texttt{a-swuwti.adb} provides similar extended \texttt{Wide_Text_IO} functionality for unbounded wide strings.

The package \texttt{Ada.Strings.Wide_Wide_Unbounded.Wide_Wide_Text_IO} in library files \texttt{a-szuzti.ads} and \texttt{a-szuzti.adb} provides similar extended \texttt{Wide_Wide_Text_IO} functionality for unbounded wide wide strings.

### 11.6 Wide_Text_IO

\texttt{Wide_Text_IO} is similar in most respects to \texttt{Text_IO}, except that both input and output files may contain special sequences that represent wide character values. The encoding scheme for a given file may be specified using a \texttt{FORM} parameter:

\begin{verbatim}
WCEM='x'
\end{verbatim}

as part of the \texttt{FORM} string (\texttt{WCEM} = wide character encoding method), where \texttt{x} is one of the following characters:

<table>
<thead>
<tr>
<th>Character</th>
<th>Encoding</th>
</tr>
</thead>
<tbody>
<tr>
<td>‘h’</td>
<td>Hex ESC encoding</td>
</tr>
<tr>
<td>‘u’</td>
<td>Upper half encoding</td>
</tr>
<tr>
<td>‘s’</td>
<td>Shift-JIS encoding</td>
</tr>
<tr>
<td>‘e’</td>
<td>EUC Encoding</td>
</tr>
<tr>
<td>‘8’</td>
<td>UTF-8 encoding</td>
</tr>
<tr>
<td>‘b’</td>
<td>Brackets encoding</td>
</tr>
</tbody>
</table>

The encoding methods match those that can be used in a source program, but there is no requirement that the encoding method used for the source program be the same as the encoding method used for files, and different files may use different encoding methods.

The default encoding method for the standard files, and for opened files for which no \texttt{WCEM} parameter is given in the \texttt{FORM} string matches the wide character encoding specified for the main program (the default being brackets encoding if no coding method was specified with \texttt{-gnatW}).

**‘Hex Coding’**

In this encoding, a wide character is represented by a five character sequence: \texttt{ESC a b c d}

where \texttt{a}, \texttt{b}, \texttt{c}, \texttt{d} are the four hexadecimal characters (using upper case letters) of the wide character code. For example, \texttt{ESC A345} is used to represent the wide character with code \texttt{16#A345#}. This scheme is compatible with use of the full \texttt{Wide_Character} set.
Chapter 11: The Implementation of Standard I/O

‘Upper Half Coding’
The wide character with encoding \texttt{16\#abcd\#}, where the upper bit is on (i.e.,
a is in the range 8-F) is represented as two bytes \texttt{16\#ab\#} and \texttt{16\#cd\#}. The
second byte may never be a format control character, but is not required to be
in the upper half. This method can be also used for shift-JIS or EUC where
the internal coding matches the external coding.

‘Shift JIS Coding’
A wide character is represented by a two character sequence \texttt{16\#ab\#} and
\texttt{16\#cd\#}, with the restrictions described for upper half encoding as described
above. The internal character code is the corresponding JIS character accord-
ing to the standard algorithm for Shift-JIS conversion. Only characters defined
in the JIS code set table can be used with this encoding method.

‘EUC Coding’
A wide character is represented by a two character sequence \texttt{16\#ab\#} and
\texttt{16\#cd\#}, with both characters being in the upper half. The internal char-
acter code is the corresponding JIS character according to the EUC encoding
algorithm. Only characters defined in the JIS code set table can be used with
this encoding method.

‘UTF-8 Coding’
A wide character is represented using UCS Transformation Format 8 (UTF-8)
as defined in Annex R of ISO 10646-1/Am.2. Depending on the character value,
the representation is a one, two, or three byte sequence:
\begin{verbatim}
16#0000#-16#007f#: 2#0xxxxxxx#
16#0080#-16#07ff#: 2#110xxxxx# 2#10xxxxxx#
16#0800#-16#ffff#: 2#1110xxxx# 2#10xxxxxx# 2#10xxxxxx#
\end{verbatim}
where the \texttt{xxx} bits correspond to the left-padded bits of the 16-bit char-
acter value. Note that all lower half ASCII characters are represented as
ASCII bytes and all upper half characters and other wide characters are
represented as sequences of upper-half (The full UTF-8 scheme allows
for encoding 31-bit characters as 6-byte sequences, but in this imple-
mentation, all UTF-8 sequences of four or more bytes length will raise a
Constraint_Error, as will all invalid UTF-8 sequences.)

‘Brackets Coding’
In this encoding, a wide character is represented by the following eight character
sequence:
\begin{verbatim}
[ " a b c d " ]
\end{verbatim}
where \texttt{a}, \texttt{b}, \texttt{c}, \texttt{d} are the four hexadecimal characters (using uppercase
letters) of the wide character code. For example, \texttt{["A345"]} is used to
represent the wide character with code \texttt{16\#A345\#}. This scheme is com-
patible with use of the full Wide_Character set. On input, brackets coding
can also be used for upper half characters, e.g., \texttt{["C1"]} for lower case \texttt{a}.
However, on output, brackets notation is only used for wide characters
with a code greater than \texttt{16\#FF\#}.

Note that brackets coding is not normally used in the context of
Wide_Text_IO or Wide_Wide_Text_IO, since it is really just designed as
Chapter 11: The Implementation of Standard I/O

a portable way of encoding source files. In the context of Wide_Text_IO or Wide_Wide_Text_IO, it can only be used if the file does not contain any instance of the left bracket character other than to encode wide character values using the brackets encoding method. In practice it is expected that some standard wide character encoding method such as UTF-8 will be used for text input output.

If brackets notation is used, then any occurrence of a left bracket in the input file which is not the start of a valid wide character sequence will cause Constraint_Error to be raised. It is possible to encode a left bracket as ["5B"] and Wide_Text_IO and Wide_Wide_Text_IO input will interpret this as a left bracket.

However, when a left bracket is output, it will be output as a left bracket and not as ["5B"]. We make this decision because for normal use of Wide_Text_IO for outputting messages, it is unpleasant to clobber left brackets. For example, if we write:

Put_Line ("Start of output [first run]");

we really do not want to have the left bracket in this message clobbered so that the output reads:

Start of output ["5B"]first run

In practice brackets encoding is reasonably useful for normal Put_Line use since we won’t get confused between left brackets and wide character sequences in the output. But for input, or when files are written out and read back in, it really makes better sense to use one of the standard encoding methods such as UTF-8.

For the coding schemes other than UTF-8, Hex, or Brackets encoding, not all wide character values can be represented. An attempt to output a character that cannot be represented using the encoding scheme for the file causes Constraint_Error to be raised. An invalid wide character sequence on input also causes Constraint_Error to be raised.

11.6.1 Stream Pointer Positioning

Ada_Wide_Text_IO is similar to Ada_Text_IO in its handling of stream pointer positioning ([Text_IO], page 239). There is one additional case:
If Ada_Wide_Text_IO.Look_Ahead reads a character outside the normal lower ASCII set, i.e. a character in the range:

Wide_Character’Val (16#0080#) .. Wide_Character’Val (16#FFFF#)

then although the logical position of the file pointer is unchanged by the Look_Ahead call, the stream is physically positioned past the wide character sequence. Again this is to avoid the need for buffering or backup, and all Wide_Text_IO routines check the internal indication that this situation has occurred so that this is not visible to a normal program using Wide_Text_IO. However, this discrepancy can be observed if the wide text file shares a stream with another file.

11.6.2 Reading and Writing Non-Regular Files

As in the case of Text_IO, when a non-regular file is read, it is assumed that the file contains no page marks (any form characters are treated as data characters), and End_Of_
Page always returns False. Similarly, the end of file indication is not sticky, so it is possible to read beyond an end of file.

11.7 Wide_Wide_Text_IO

Wide_Wide_Text_IO is similar in most respects to Text_IO, except that both input and output files may contain special sequences that represent wide wide character values. The encoding scheme for a given file may be specified using a FORM parameter:

\[ \text{WCEM='x'} \]

as part of the FORM string (WCEM = wide character encoding method), where x is one of the following characters

<table>
<thead>
<tr>
<th>Character</th>
<th>Encoding</th>
</tr>
</thead>
<tbody>
<tr>
<td>‘h’</td>
<td>Hex ESC encoding</td>
</tr>
<tr>
<td>‘u’</td>
<td>Upper half encoding</td>
</tr>
<tr>
<td>‘s’</td>
<td>Shift-JIS encoding</td>
</tr>
<tr>
<td>‘e’</td>
<td>EUC Encoding</td>
</tr>
<tr>
<td>‘8’</td>
<td>UTF-8 encoding</td>
</tr>
<tr>
<td>‘b’</td>
<td>Brackets encoding</td>
</tr>
</tbody>
</table>

The encoding methods match those that can be used in a source program, but there is no requirement that the encoding method used for the source program be the same as the encoding method used for files, and different files may use different encoding methods.

The default encoding method for the standard files, and for opened files for which no WCEM parameter is given in the FORM string matches the wide character encoding specified for the main program (the default being brackets encoding if no coding method was specified with -gnatW).

‘UTF-8 Coding’

A wide character is represented using UCS Transformation Format 8 (UTF-8) as defined in Annex R of ISO 10646-1/Am.2. Depending on the character value, the representation is a one, two, three, or four byte sequence:

\[
\begin{align*}
16\#000000\# - 16\#00007f\# & : 2\#0xxxxxxx\# \\
16\#000080\# - 16\#0007ff\# & : 2\#110xxxxx\# 2\#10xxxxxx\# \\
16\#000800\# - 16\#00ff\# & : 2\#1110xxxx\# 2\#10xxxxxx\# 2\#10xxxxxx\# \\
16\#010000\# - 16\#10ff\# & : 2\#11110xxx\# 2\#10xxxxxx\# 2\#10xxxxxx\# 2\#10xxxxxx\#
\end{align*}
\]

where the xxx bits correspond to the left-padded bits of the 21-bit character value. Note that all lower half ASCII characters are represented as ASCII bytes and all upper half characters and other wide characters are represented as sequences of upper-half characters.
Brackets Coding

In this encoding, a wide wide character is represented by the following eight character sequence if is in wide character range

[ " a b c d " ]

and by the following ten character sequence if not

[ " a b c d e f " ]

where a, b, c, d, e, and f are the four or six hexadecimal characters (using uppercase letters) of the wide wide character code. For example, ["01A345"] is used to represent the wide wide character with code 16#01A345#.

This scheme is compatible with use of the full Wide_Wide_Character set. On input, brackets coding can also be used for upper half characters, e.g., ["C1"] for lower case a. However, on output, brackets notation is only used for wide characters with a code greater than 16#FF#.

If is also possible to use the other Wide_Character encoding methods, such as Shift-JIS, but the other schemes cannot support the full range of wide wide characters. An attempt to output a character that cannot be represented using the encoding scheme for the file causes Constraint_Error to be raised. An invalid wide character sequence on input also causes Constraint_Error to be raised.

11.7.1 Stream Pointer Positioning

Ada.Wide_Wide_Text_IO is similar to Ada.Text_IO in its handling of stream pointer positioning ([Text_IO], page 239). There is one additional case:

If Ada.Wide_Wide_Text_IO.Look_Ahead reads a character outside the normal lower ASCII set, i.e. a character in the range:

Wide_Wide_Character'Val (16#0080#) .. Wide_Wide_Character'Val (16#10FFFF#)

then although the logical position of the file pointer is unchanged by the Look_Ahead call, the stream is physically positioned past the wide character sequence. Again this is to avoid the need for buffering or backup, and all Wide_Wide_Text_IO routines check the internal indication that this situation has occurred so that this is not visible to a normal program using Wide_Wide_Text_IO. However, this discrepancy can be observed if the wide text file shares a stream with another file.

11.7.2 Reading and Writing Non-Regular Files

As in the case of Text_IO, when a non-regular file is read, it is assumed that the file contains no page marks (any form characters are treated as data characters), and End_Of_Page always returns False. Similarly, the end of file indication is not sticky, so it is possible to read beyond an end of file.

11.8 Stream_IO

A stream file is a sequence of bytes, where individual elements are written to the file as described in the Ada Reference Manual. The type Stream_Element is simply a byte. There are two ways to read or write a stream file.
* The operations Read and Write directly read or write a sequence of stream elements with no control information.
* The stream attributes applied to a stream file transfer data in the manner described for stream attributes.

11.9 Text Translation

The Translation=xxx may be used as the Form parameter passed to Text.IO.Create and Text.IO.Open. Translation=xxx has no effect on Unix systems. Possible values are:
* Yes or Text is the default, which means to translate LF to/from CR/LF on Windows systems.
  No disables this translation; i.e. it uses binary mode. For output files, Translation=No may be used to create Unix-style files on Windows.
* wtext translation enabled in Unicode mode. (corresponds to O_WTEXT).
* u8text translation enabled in Unicode UTF-8 mode. (corresponds to O_U8TEXT).
* u16text translation enabled in Unicode UTF-16 mode. (corresponds to O_U16TEXT).

11.10 Shared Files

Section A.14 of the Ada Reference Manual allows implementations to provide a wide variety of behavior if an attempt is made to access the same external file with two or more internal files.

To provide a full range of functionality, while at the same time minimizing the problems of portability caused by this implementation dependence, GNAT handles file sharing as follows:
* In the absence of a shared=xxx form parameter, an attempt to open two or more files with the same full name is considered an error and is not supported. The exception Use_Error will be raised. Note that a file that is not explicitly closed by the program remains open until the program terminates.
* If the form parameter shared=no appears in the form string, the file can be opened or created with its own separate stream identifier, regardless of whether other files sharing the same external file are opened. The exact effect depends on how the C stream routines handle multiple accesses to the same external files using separate streams.
* If the form parameter shared=yes appears in the form string for each of two or more files opened using the same full name, the same stream is shared between these files, and the semantics are as described in Ada Reference Manual, Section A.14.

When a program that opens multiple files with the same name is ported from another Ada compiler to GNAT, the effect will be that Use_Error is raised.

The documentation of the original compiler and the documentation of the program should then be examined to determine if file sharing was expected, and shared=xxx parameters added to Open and Create calls as required.

When a program is ported from GNAT to some other Ada compiler, no special attention is required unless the shared=xxx form parameter is used in the program. In this case, you must examine the documentation of the new compiler to see if it supports the required file sharing semantics, and form strings modified appropriately. Of course it may be the case
that the program cannot be ported if the target compiler does not support the required functionality. The best approach in writing portable code is to avoid file sharing (and hence the use of the `shared=xxx` parameter in the form string) completely.

One common use of file sharing in Ada 83 is the use of instantiations of Sequential_IO on the same file with different types, to achieve heterogeneous input-output. Although this approach will work in GNAT if `shared=yes` is specified, it is preferable in Ada to use Stream_IO for this purpose (using the stream attributes)

### 11.11 Filenames encoding

An encoding form parameter can be used to specify the filename encoding `encoding=xxx`.

* If the form parameter `encoding=utf8` appears in the form string, the filename must be encoded in UTF-8.
* If the form parameter `encoding=8bits` appears in the form string, the filename must be a standard 8bits string.

In the absence of a `encoding=xxx` form parameter, the encoding is controlled by the `GNAT_CODE_PAGE` environment variable. And if not set `utf8` is assumed.

- **‘CP_ACP’**
  
  The current system Windows ANSI code page.

- **‘CP_UTF8’**
  
  UTF-8 encoding

This encoding form parameter is only supported on the Windows platform. On the other Operating Systems the run-time is supporting UTF-8 natively.

### 11.12 File content encoding

For text files it is possible to specify the encoding to use. This is controlled by the by the `GNAT_CCS_ENCODING` environment variable. And if not set `TEXT` is assumed.

The possible values are those supported on Windows:

- **‘TEXT’**
  
  Translated text mode

- **‘WTEXT’**
  
  Translated unicode encoding

- **‘U16TEXT’**
  
  Unicode 16-bit encoding

- **‘U8TEXT’**
  
  Unicode 8-bit encoding

This encoding is only supported on the Windows platform.
11.13 Open Modes

Open and Create calls result in a call to fopen using the mode shown in the following table:

<table>
<thead>
<tr>
<th>Open and Create Call Modes</th>
<th>‘OPEN’</th>
<th>‘CREATE’</th>
</tr>
</thead>
<tbody>
<tr>
<td>Append_File</td>
<td>“r+”</td>
<td>“w+”</td>
</tr>
<tr>
<td>In_File</td>
<td>“r”</td>
<td>“w+”</td>
</tr>
<tr>
<td>Out_File (Direct_IO)</td>
<td>“r+”</td>
<td>“w”</td>
</tr>
<tr>
<td>Out_File (all other cases)</td>
<td>“w”</td>
<td>“w”</td>
</tr>
<tr>
<td>Inout_File</td>
<td>“r+”</td>
<td>“w+”</td>
</tr>
</tbody>
</table>

If text file translation is required, then either b or t is added to the mode, depending on the setting of Text. Text file translation refers to the mapping of CR/LF sequences in an external file to LF characters internally. This mapping only occurs in DOS and DOS-like systems, and is not relevant to other systems.

A special case occurs with Stream_IO. As shown in the above table, the file is initially opened in r or w mode for the In_File and Out_File cases. If a Set_Mode operation subsequently requires switching from reading to writing or vice-versa, then the file is reopened in r+ mode to permit the required operation.

11.14 Operations on C Streams

The package Interfaces.C_Streams provides an Ada program with direct access to the C library functions for operations on C streams:

```ada
package Interfaces.C_Streams is
  -- Note: the reason we do not use the types that are in
  -- Interfaces.C is that we want to avoid dragging in the
  -- code in this unit if possible.
  subtype chars is System.Address;
  -- Pointer to null-terminated array of characters
  subtype FILEs is System.Address;
  -- Corresponds to the C type FILE*
  subtype voids is System.Address;
  -- Corresponds to the C type void*
  subtype int is Integer;
  -- Note: the above types are subtypes deliberately, and it
  -- is part of this spec that the above correspondences are
  -- guaranteed. This means that it is legitimate to, for
  -- example, use Integer instead of int. We provide these
```
-- synonyms for clarity, but in some cases it may be
-- convenient to use the underlying types (for example to
-- avoid an unnecessary dependency of a spec on the spec
-- of this unit).
type size_t is mod 2 ** Standard'Address_Size;
NULL_Stream : constant FILEs;
-- Value returned (NULL in C) to indicate an
-- fdopen/fopen/tmpfile error
----------------------------------
-- Constants Defined in stdio.h --
----------------------------------
EOF : constant int;
-- Used by a number of routines to indicate error or
-- end of file
I0FBF : constant int;
IOLBF : constant int;
IONBF : constant int;
-- Used to indicate buffering mode for setvbuf call
SEEK_CUR : constant int;
SEEK_END : constant int;
SEEK_SET : constant int;
-- Used to indicate origin for fseek call
function stdin return FILEs;
function stdout return FILEs;
function stderr return FILEs;
-- Streams associated with standard files
--------------------------
-- Standard C functions --
--------------------------
-- The functions selected below are ones that are
-- available in UNIX (but not necessarily in ANSI C).
-- These are very thin interfaces
-- which copy exactly the C headers. For more
-- documentation on these functions, see the Microsoft C
-- "Run-Time Library Reference" (Microsoft Press, 1990,
-- ISBN 1-55615-225-6), which includes useful information
-- on system compatibility.
procedure clearerr (stream : FILEs);
function fclose (stream : FILEs) return int;
function fdopen (handle : int; mode : chars) return FILEs;
function feof (stream : FILEs) return int;
function ferror (stream : FILEs) return int;
function fflush (stream : FILEs) return int;
function fgetc (stream : FILEs) return int;
function fgets (strng : chars; n : int; stream : FILEs)
           return chars;
function fileno (stream : FILEs) return int;
function fopen (filename : chars; Mode : chars)
    return FILEs;
-- Note: to maintain target independence, use
-- text_translation_required, a boolean variable defined in
-- a-sysdep.c to deal with the target dependent text
-- translation requirement. If this variable is set,
-- then b/t should be appended to the standard mode
-- argument to set the text translation mode off or on
-- as required.
function fputc (C : int; stream : FILEs) return int;
function fputs (Strng : chars; Stream : FILEs) return int;
function fread
    (buffer : voids;
    size : size_t;
    count : size_t;
    stream : FILEs)
    return size_t;
function freopen
    (filename : chars;
    mode : chars;
    stream : FILEs)
    return FILEs;
function fseek
    (stream : FILEs;
    offset : long;
    origin : int)
    return int;
function ftell (stream : FILEs) return long;
function fwrite
    (buffer : voids;
    size : size_t;
    count : size_t;
    stream : FILEs)
    return size_t;
function isatty (handle : int) return int;
procedure mktemp (template : chars);
-- The return value (which is just a pointer to template)
-- is discarded
procedure rewind (stream : FILEs);
function rmtmp return int;
function setvbuf
    (stream : FILEs;
    buffer : chars;
    mode : int;
    size : size_t)
    return int;
function tmpfile return FILEs;
function ungetc (c : int; stream : FILEs) return int;
function unlink (filename : chars) return int;

--- Extra functions ---

--- These functions supply slightly thicker bindings than
--- those above. They are derived from functions in the
--- C Run-Time Library, but may do a bit more work than
--- just directly calling one of the Library functions.
function is_regular_file (handle : int) return int;
--- Tests if given handle is for a regular file (result 1)
--- or for a non-regular file (pipe or device, result 0).

--- Control of Text/Binary Mode ---

--- If text_translation_required is true, then the following
--- functions may be used to dynamically switch a file from
--- binary to text mode or vice versa. These functions have
--- no effect if text_translation_required is false (i.e., in
--- normal UNIX mode). Use fileno to get a stream handle.
procedure set_binary_mode (handle : int);
procedure set_text_mode (handle : int);

--- Full Path Name support ---

procedure full_name (nam : chars; buffer : chars);
--- Given a NUL terminated string representing a file
--- name, returns in buffer a NUL terminated string
--- representing the full path name for the file name.
--- On systems where it is relevant the drive is also
--- part of the full path name. It is the responsibility
--- of the caller to pass an actual parameter for buffer
--- that is big enough for any full path name. Use
--- max_path_len given below as the size of buffer.
max_path_len : integer;
--- Maximum length of an allowable full path name on the
--- system, including a terminating NUL character.

end Interfaces.C_Streams;

11.15 Interfacing to C Streams

The packages in this section permit interfacing Ada files to C Stream operations.

with Interfaces.C_Streams;
package Ada.Sequential_IO.C_Streams is
    function C_Stream (F : File_Type) return Interfaces.C_Streams.FILEs;
procedure Open
   (File : in out File_Type;
    Mode : in File_Mode;
    C_Stream : in Interfaces.C_Streams.FILEs;
    Form : in String := "")
end Ada.Sequential_IO.C_Streams;

with Interfaces.C_Streams;
package Ada.Direct_IO.C_Streams is
   function C_Stream (F : File_Type)
      return Interfaces.C_Streams.FILEs;
   procedure Open
      (File : in out File_Type;
       Mode : in File_Mode;
       C_Stream : in Interfaces.C_Streams.FILEs;
       Form : in String := "")
   end Ada.Direct_IO.C_Streams;

with Interfaces.C_Streams;
package Ada.Text_IO.C_Streams is
   function C_Stream (F : File_Type)
      return Interfaces.C_Streams.FILEs;
   procedure Open
      (File : in out File_Type;
       Mode : in File_Mode;
       C_Stream : in Interfaces.C_Streams.FILEs;
       Form : in String := "")
   end Ada.Text_IO.C_Streams;

with Interfaces.C_Streams;
package Ada.Wide_Text_IO.C_Streams is
   function C_Stream (F : File_Type)
      return Interfaces.C_Streams.FILEs;
   procedure Open
      (File : in out File_Type;
       Mode : in File_Mode;
       C_Stream : in Interfaces.C_Streams.FILEs;
       Form : in String := "")
   end Ada.Wide_Text_IO.C_Streams;

with Interfaces.C_Streams;
package Ada.Wide_Wide_Text_IO.C_Streams is
   function C_Stream (F : File_Type)
      return Interfaces.C_Streams.FILEs;
   procedure Open
      (File : in out File_Type;
       Mode : in File_Mode;
       C_Stream : in Interfaces.C_Streams.FILEs;
       Form : in String := "")
   end Ada.Wide_Wide_Text_IO.C_Streams;
In each of these six packages, the \texttt{C\_Stream} function obtains the \texttt{FILE} pointer from a currently opened Ada file. It is then possible to use the \texttt{Interfaces.C\_Streams} package to operate on this stream, or the stream can be passed to a C program which can operate on it directly. Of course the program is responsible for ensuring that only appropriate sequences of operations are executed.

One particular use of relevance to an Ada program is that the \texttt{setvbuf} function can be used to control the buffering of the stream used by an Ada file. In the absence of such a call the standard default buffering is used.

The \texttt{Open} procedures in these packages open a file giving an existing C Stream instead of a file name. Typically this stream is imported from a C program, allowing an Ada file to operate on an existing C file.
12 The GNAT Library

The GNAT library contains a number of general and special purpose packages. It represents functionality that the GNAT developers have found useful, and which is made available to GNAT users. The packages described here are fully supported, and upwards compatibility will be maintained in future releases, so you can use these facilities with the confidence that the same functionality will be available in future releases.

The chapter here simply gives a brief summary of the facilities available. The full documentation is found in the spec file for the package. The full sources of these library packages, including both spec and body, are provided with all GNAT releases. For example, to find out the full specifications of the SPITBOL pattern matching capability, including a full tutorial and extensive examples, look in the `g-spipat.ads` file in the library.

For each entry here, the package name (as it would appear in a `with` clause) is given, followed by the name of the corresponding spec file in parentheses. The packages are children in four hierarchies, Ada, Interfaces, System, and GNAT, the latter being a GNAT-specific hierarchy.

Note that an application program should only use packages in one of these four hierarchies if the package is defined in the Ada Reference Manual, or is listed in this section of the GNAT Programmers Reference Manual. All other units should be considered internal implementation units and should not be directly `with`ed by application code. The use of a `with` clause that references one of these internal implementation units makes an application potentially dependent on changes in versions of GNAT, and will generate a warning message.


This child of Ada.Characters provides a set of definitions corresponding to those in the RM-defined package Ada.Characters.Latin_1 but with the few modifications required for Latin-9. The provision of such a package is specifically authorized by the Ada Reference Manual (RM A.3.3(27)).

12.2 Ada.Characters.Wide_Latin_1 (a-cwila1.ads)

This child of Ada.Characters provides a set of definitions corresponding to those in the RM-defined package Ada.Characters.Latin_1 but with the types of the constants being `Wide_Character` instead of `Character`. The provision of such a package is specifically authorized by the Ada Reference Manual (RM A.3.3(27)).


This child of Ada.Characters provides a set of definitions corresponding to those in the GNAT defined package Ada.Characters.Latin_9 but with the types of the constants being `Wide_Character` instead of `Character`. The provision of such a package is specifically authorized by the Ada Reference Manual (RM A.3.3(27)).

12.4 Ada.Characters.Wide_Wide_Latin_1 (a-chzla1.ads)

This child of Ada.Characters provides a set of definitions corresponding to those in the RM-defined package Ada.Characters.Latin_1 but with the types of the constants being
Wide_Wide_Character instead of Character. The provision of such a package is specifically authorized by the Ada Reference Manual (RM A.3.3(27)).

12.5 Ada.Characters.Wide_Wide_Latin_9 (a-chzla9.ads)
This child of Ada.Characters provides a set of definitions corresponding to those in the GNAT defined package Ada.Characters.Latin_9 but with the types of the constants being Wide_Wide_Character instead of Character. The provision of such a package is specifically authorized by the Ada Reference Manual (RM A.3.3(27)).

12.6 Ada.Containers.Bounded_Holders (a-coboho.ads)
This child of Ada.Containers defines a modified version of Indefinite_Holders that avoids heap allocation.

12.7 Ada.Command_Line.Environment (a-colien.ads)
This child of Ada.Command_Line provides a mechanism for obtaining environment values on systems where this concept makes sense.

12.8 Ada.Command_Line.Remove (a-colire.ads)
This child of Ada.Command_Line provides a mechanism for logically removing arguments from the argument list. Once removed, an argument is not visible to further calls to the subprograms in Ada.Command_Line. These calls will not see the removed argument.

This child of Ada.Command_Line provides a mechanism facilities for getting command line arguments from a text file, called a “response file”. Using a response file allow passing a set of arguments to an executable longer than the maximum allowed by the system on the command line.

12.10 Ada.Direct_IO.C_Streams (a-diocst.ads)
This package provides subprograms that allow interfacing between C streams and Direct_IO. The stream identifier can be extracted from a file opened on the Ada side, and an Ada file can be constructed from a stream opened on the C side.

12.11 Ada.Exceptions.Is_Null_Occurrence (a-einuoc.ads)
This child subprogram provides a way of testing for the null exception occurrence (Null_Occurrence) without raising an exception.

12.12 Ada.Exceptions.Last_Chance_Handler (a-elchha.ads)
This child subprogram is used for handling otherwise unhandled exceptions (hence the name last chance), and perform clean ups before terminating the program. Note that this subprogram never returns.
12.13 Ada.Exceptions.Traceback (a-exctra.ads)
This child package provides the subprogram (Tracebacks) to give a traceback array of addresses based on an exception occurrence.

12.14 Ada.Sequential_IO.C_Streams (a-siocst.ads)
This package provides subprograms that allow interfacing between C streams and Sequential_IO. The stream identifier can be extracted from a file opened on the Ada side, and an Ada file can be constructed from a stream opened on the C side.

12.15 Ada.Streams.Stream_IO.C_Streams (a-ssicst.ads)
This package provides subprograms that allow interfacing between C streams and Stream_IO. The stream identifier can be extracted from a file opened on the Ada side, and an Ada file can be constructed from a stream opened on the C side.

12.16 Ada.Strings.Unbounded.Text_IO (a-suteio.ads)
This package provides subprograms for Text_IO for unbounded strings, avoiding the necessity for an intermediate operation with ordinary strings.

12.17 Ada.Strings.Wide_Unbounded.Wide_Text_IO (a-swuwti.ads)
This package provides subprograms for Text_IO for unbounded wide strings, avoiding the necessity for an intermediate operation with ordinary wide strings.

12.18 Ada.Strings.Wide_Wide_Unbounded.Wide_Wide_Text_IO (a-szwuzti.ads)
This package provides subprograms for Text_IO for unbounded wide wide strings, avoiding the necessity for an intermediate operation with ordinary wide wide strings.

12.19 Ada.Task_Initialization (a-tasini.ads)
This package provides a way to set a global initialization handler that is automatically invoked whenever a task is activated. Handlers are parameterless procedures. Note that such a handler is only invoked for those tasks activated after the handler is set.

12.20 Ada.Text_IO.C_Streams (a-tiocst.ads)
This package provides subprograms that allow interfacing between C streams and Text_IO. The stream identifier can be extracted from a file opened on the Ada side, and an Ada file can be constructed from a stream opened on the C side.

12.21 Ada.Text_IO.Reset_Standard_Files (a-tirsfi.ads)
This procedure is used to reset the status of the standard files used by Ada.Text_IO. This is useful in a situation (such as a restart in an embedded application) where the status of the files may change during execution (for example a standard input file may be redefined to be interactive).
12.22 Ada.Wide_Characters.Unicode (a-wichun.ads)
This package provides subprograms that allow categorization of Wide_Character values according to Unicode categories.

12.23 Ada.Wide_Text_IO.C_Streams (a-wtcstr.ads)
This package provides subprograms that allow interfacing between C streams and Wide_Text_IO. The stream identifier can be extracted from a file opened on the Ada side, and an Ada file can be constructed from a stream opened on the C side.

12.24 Ada.Wide_Text_IO.Reset_Standard_Files (a-wrstfi.ads)
This procedure is used to reset the status of the standard files used by Ada.Wide_Text_IO. This is useful in a situation (such as a restart in an embedded application) where the status of the files may change during execution (for example a standard input file may be redefined to be interactive).

12.25 Ada.Wide_Wide_Characters.Unicode (a-zchuni.ads)
This package provides subprograms that allow categorization of Wide_Wide_Character values according to Unicode categories.

12.26 Ada.Wide_Wide_Text_IO.C_Streams (a-ztcstr.ads)
This package provides subprograms that allow interfacing between C streams and Wide_Wide_Text_IO. The stream identifier can be extracted from a file opened on the Ada side, and an Ada file can be constructed from a stream opened on the C side.

12.27 Ada.Wide_Wide_Text_IO.Reset_Standard_Files (a-zrstfi.ads)
This procedure is used to reset the status of the standard files used by Ada.Wide_Wide_Text_IO. This is useful in a situation (such as a restart in an embedded application) where the status of the files may change during execution (for example a standard input file may be redefined to be interactive).

12.28 GNAT.Altivec (g-altive.ads)
This is the root package of the GNAT AltiVec binding. It provides definitions of constants and types common to all the versions of the binding.

12.29 GNAT.Altivec.Conversions (g-altcon.ads)
This package provides the Vector/View conversion routines.

12.30 GNAT.Altivec.Vector_Operations (g-alveop.ads)
This package exposes the Ada interface to the AltiVec operations on vector objects. A soft emulation is included by default in the GNAT library. The hard binding is provided as a separate package. This unit is common to both bindings.
12.31 GNAT.Altivec.Vector_Types (g-alvety.ads)
This package exposes the various vector types part of the Ada binding to Altivec facilities.

12.32 GNAT.Altivec.Vector_Views (g-alvevi.ads)
This package provides public ‘View’ data types from/to which private vector representations can be converted via GNAT.Altivec.Conversions. This allows convenient access to individual vector elements and provides a simple way to initialize vector objects.

12.33 GNAT.Array_Split (g-arrspl.ads)
Useful array-manipulation routines: given a set of separators, split an array wherever the separators appear, and provide direct access to the resulting slices.

12.34 GNAT.AWK (g-awk.ads)
Provides AWK-like parsing functions, with an easy interface for parsing one or more files containing formatted data. The file is viewed as a database where each record is a line and a field is a data element in this line.

12.35 GNAT.Binary_Search (g-binsea.ads)
Allow binary search of a sorted array (or of an array-like container; the generic does not reference the array directly).

12.36 GNAT.Bind_Environment (g-binenv.ads)
Provides access to key=value associations captured at bind time. These associations can be specified using the -V binder command line switch.

12.37 GNAT.Branch_Prediction (g-brapre.ads)
Provides routines giving hints to the branch predictor of the code generator.

12.38 GNAT.Bounded_Buffers (g-boubuf.ads)
Provides a concurrent generic bounded buffer abstraction. Instances are useful directly or as parts of the implementations of other abstractions, such as mailboxes.

12.39 GNAT.Bounded_Mailboxes (g-boumai.ads)
Provides a thread-safe asynchronous intertask mailbox communication facility.

12.40 GNAT.Bubble_Sort (g-bubsor.ads)
Provides a general implementation of bubble sort usable for sorting arbitrary data items. Exchange and comparison procedures are provided by passing access-to-procedure values.
12.41 GNAT.Bubble_Sort_A (g-busora.ads)
Provides a general implementation of bubble sort usable for sorting arbitrary data items. Move and comparison procedures are provided by passing access-to-procedure values. This is an older version, retained for compatibility. Usually GNAT.Bubble_Sort will be preferable.

12.42 GNAT.Bubble_Sort_G (g-busorg.ads)
Similar to Bubble_Sort_A except that the move and sorting procedures are provided as generic parameters, this improves efficiency, especially if the procedures can be inlined, at the expense of duplicating code for multiple instantiations.

12.43 GNAT.Byte_Order_Mark (g-byorma.ads)
Provides a routine which given a string, reads the start of the string to see whether it is one of the standard byte order marks (BOM’s) which signal the encoding of the string. The routine includes detection of special XML sequences for various UCS input formats.

12.44 GNAT.Byte_Swapping (g-bytswa.ads)
General routines for swapping the bytes in 2-, 4-, and 8-byte quantities. Machine-specific implementations are available in some cases.

12.45 GNAT.Calendar (g-calend.ads)
Extends the facilities provided by Ada.Calendar to include handling of days of the week, an extended Split and Time_Of capability. Also provides conversion of Ada.Calendar.Time values to and from the C timeval format.

12.46 GNAT.Calendar.Time_IO (g-catiio.ads)

12.47 GNAT.CRC32 (g-crc32.ads)
This package implements the CRC-32 algorithm. For a full description of this algorithm see ‘Computation of Cyclic Redundancy Checks via Table Look-Up’, Communications of the ACM, Vol. 31 No. 8, pp. 1008-1013, Aug. 1988. Sarwate, D.V.

12.48 GNAT.Case_Util (g-casuti.ads)
A set of simple routines for handling upper and lower casing of strings without the overhead of the full casing tables in Ada.Characters.Handling.

12.49 GNAT.CGI (g-cgi.ads)
This is a package for interfacing a GNAT program with a Web server via the Common Gateway Interface (CGI). Basically this package parses the CGI parameters, which are a set of key/value pairs sent by the Web server. It builds a table whose index is the key and provides some services to deal with this table.
12.50 GNAT.CGI.Cookie (g-cgicoo.ads)
This is a package to interface a GNAT program with a Web server via the Common Gateway Interface (CGI). It exports services to deal with Web cookies (piece of information kept in the Web client software).

12.51 GNAT.CGI.Debug (g-cgideb.ads)
This is a package to help debugging CGI (Common Gateway Interface) programs written in Ada.

12.52 GNAT.Command_Line (g-comlin.ads)
Provides a high level interface to Ada.Command_Line facilities, including the ability to scan for named switches with optional parameters and expand file names using wildcard notations.

12.53 GNAT.Compiler_Version (g-comver.ads)
Provides a routine for obtaining the version of the compiler used to compile the program. More accurately this is the version of the binder used to bind the program (this will normally be the same as the version of the compiler if a consistent tool set is used to compile all units of a partition).

12.54 GNAT.Ctrl_C (g-ctrl_c.ads)
Provides a simple interface to handle Ctrl-C keyboard events.

12.55 GNAT.Current_Exception (g-curexc.ads)
Provides access to information on the current exception that has been raised without the need for using the Ada 95 / Ada 2005 exception choice parameter specification syntax. This is particularly useful in simulating typical facilities for obtaining information about exceptions provided by Ada 83 compilers.

12.56 GNAT.Debug_Pools (g-debpoo.ads)
Provides a debugging storage pools that helps tracking memory corruption problems. See The GNAT Debug_Pool Facility section in the GNAT User’s Guide.

12.57 GNAT.Debug_Utillities (g-debuti.ads)
Provides a few useful utilities for debugging purposes, including conversion to and from string images of address values. Supports both C and Ada formats for hexadecimal literals.

12.58 GNAT.Decode_String (g-decstr.ads)
A generic package providing routines for decoding wide character and wide wide character strings encoded as sequences of 8-bit characters using a specified encoding method. Includes validation routines, and also routines for stepping to next or previous encoded character in an encoded string. Useful in conjunction with Unicode character coding. Note there is a preinstantiation for UTF-8. See next entry.
12.59  **GNAT.Decode_UTF8_String** *(g-deutst.ads)*
A preinstantiation of GNAT.Decode_Strings for UTF-8 encoding.

12.60  **GNAT.Directory_Operations** *(g-dirope.ads)*
Provides a set of routines for manipulating directories, including changing the current directory, making new directories, and scanning the files in a directory.

12.61  **GNAT.Directory_Operations.Iteration** *(g-diopit.ads)*
A child unit of GNAT.Directory_Operations providing additional operations for iterating through directories.

12.62  **GNAT.Dynamic_HTables** *(g-dynhta.ads)*
A generic implementation of hash tables that can be used to hash arbitrary data. Provided in two forms, a simple form with built in hash functions, and a more complex form in which the hash function is supplied.
This package provides a facility similar to that of GNAT.HTable, except that this package declares a type that can be used to define dynamic instances of the hash table, while an instantiation of GNAT.HTable creates a single instance of the hash table.

12.63  **GNAT.Dynamic_Tables** *(g-dyntab.ads)*
A generic package providing a single dimension array abstraction where the length of the array can be dynamically modified.
This package provides a facility similar to that of GNAT.Table, except that this package declares a type that can be used to define dynamic instances of the table, while an instantiation of GNAT.Table creates a single instance of the table type.

12.64  **GNAT.Encode_String** *(g-encstr.ads)*
A generic package providing routines for encoding wide character and wide wide character strings as sequences of 8-bit characters using a specified encoding method. Useful in conjunction with Unicode character coding. Note there is a preinstantiation for UTF-8. See next entry.

12.65  **GNAT.Encode_UTF8_String** *(g-enutst.ads)*
A preinstantiation of GNAT.Encode_Strings for UTF-8 encoding.

12.66  **GNAT.ExceptionActions** *(g-exact.ads)*
Provides callbacks when an exception is raised. Callbacks can be registered for specific exceptions, or when any exception is raised. This can be used for instance to force a core dump to ease debugging.

12.67  **GNAT.ExceptionTraces** *(g-exctra.ads)*
Provides an interface allowing to control automatic output upon exception occurrences.
12.68 GNAT.Exceptions (g-except.ads)

Normally it is not possible to raise an exception with a message from a subprogram in a pure package, since the necessary types and subprograms are in Ada.Exceptions which is not a pure unit. GNAT.Exceptions provides a facility for getting around this limitation for a few predefined exceptions, and for example allows raising Constraint_Error with a message from a pure subprogram.

12.69 GNAT.Expect (g-expect.ads)

Provides a set of subprograms similar to what is available with the standard Tcl Expect tool. It allows you to easily spawn and communicate with an external process. You can send commands or inputs to the process, and compare the output with some expected regular expression. Currently GNAT.Expect is implemented on all native GNAT ports. It is not implemented for cross ports, and in particular is not implemented for VxWorks or LynxOS.

12.70 GNAT.Expect.TTY (g-exptty.ads)

As GNAT.Expect but using pseudo-terminal. Currently GNAT.Expect.TTY is implemented on all native GNAT ports. It is not implemented for cross ports, and in particular is not implemented for VxWorks or LynxOS.

12.71 GNAT.Float_Control (g-flocon.ads)

Provides an interface for resetting the floating-point processor into the mode required for correct semantic operation in Ada. Some third party library calls may cause this mode to be modified, and the Reset procedure in this package can be used to reestablish the required mode.

12.72 GNAT.Formatted_String (g-forstr.ads)

Provides support for C/C++ printf() formatted strings. The format is copied from the printf() routine and should therefore give identical output. Some generic routines are provided to be able to use types derived from Integer, Float or enumerations as values for the formatted string.

12.73 GNAT.Generic_Fast_Math_Functions (g-gfmafu.ads)

Provides direct access to the underlying implementation of the common mathematical functions, generally from the system mathematical library. This differs from Ada.Numerics.Generic_Elementary_Functions in that the implementation may deviate from the semantics specified for these functions in the Reference Manual, for example Numerics.Argument_Error is not raised. On selected platforms, some of these functions may also have a vector implementation that can be automatically used by the compiler when auto-vectorization is enabled.

12.74 GNAT.Heap_Sort (g-heasor.ads)

Provides a general implementation of heap sort usable for sorting arbitrary data items. Exchange and comparison procedures are provided by passing access-to-procedure values. The
algorithm used is a modified heap sort that performs approximately N*\log(N) comparisons in the worst case.

12.75 GNAT.Heap_Sort_A (g-hesora.ads)
Provides a general implementation of heap sort usable for sorting arbitrary data items. Move and comparison procedures are provided by passing access-to-procedure values. The algorithm used is a modified heap sort that performs approximately N*\log(N) comparisons in the worst case. This differs from GNAT.Heap_Sort in having a less convenient interface, but may be slightly more efficient.

12.76 GNAT.Heap_Sort_G (g-hesorg.ads)
Similar to Heap_Sort_A except that the move and sorting procedures are provided as generic parameters, this improves efficiency, especially if the procedures can be inlined, at the expense of duplicating code for multiple instantiations.

12.77 GNAT.HTable (g-htable.ads)
A generic implementation of hash tables that can be used to hash arbitrary data. Provides two approaches, one a simple static approach, and the other allowing arbitrary dynamic hash tables.

12.78 GNAT.IO (g-io.ads)
A simple preelaborable input-output package that provides a subset of simple Text_IO functions for reading characters and strings from Standard Input, and writing characters, strings and integers to either Standard Output or Standard Error.

12.79 GNAT.IO_Aux (g-io_aux.ads)
Provides some auxiliary functions for use with Text_IO, including a test for whether a file exists, and functions for reading a line of text.

12.80 GNAT.Lock_Files (g-locfil.ads)
Provides a general interface for using files as locks. Can be used for providing program level synchronization.

12.81 GNAT.MBBS_Discrete_Random (g-mbdira.ads)

12.82 GNAT.MBBS_Float_Random (g-mbflra.ads)
The original implementation of Ada.Numerics.Float_Random. Uses a modified version of the Blum-Blum-Shub generator.
12.83 GNAT.MD5 (g-md5.ads)
Implements the MD5 Message-Digest Algorithm as described in RFC 1321, and the HMAC-MD5 message authentication function as described in RFC 2104 and FIPS PUB 198.

12.84 GNAT.Memory_Dump (g-memdum.ads)
Provides a convenient routine for dumping raw memory to either the standard output or standard error files. Uses GNAT.IO for actual output.

12.85 GNAT.Most_Recent_Exception (g-moreex.ads)
Provides access to the most recently raised exception. Can be used for various logging purposes, including duplicating functionality of some Ada 83 implementation dependent extensions.

12.86 GNAT.OS_Lib (g-os_lib.ads)
Provides a range of target independent operating system interface functions, including time/date management, file operations, subprocess management, including a portable spawn procedure, and access to environment variables and error return codes.

12.87 GNAT.Perfect_Hash_Generators (g-pehage.ads)
Provides a generator of static minimal perfect hash functions. No collisions occur and each item can be retrieved from the table in one probe (perfect property). The hash table size corresponds to the exact size of the key set and no larger (minimal property). The key set has to be known in advance (static property). The hash functions are also order preserving. If w2 is inserted after w1 in the generator, their hashcode are in the same order. These hashing functions are very convenient for use with realtime applications.

12.88 GNAT.Random_Numbers (g-rannum.ads)
Provides random number capabilities which extend those available in the standard Ada library and are more convenient to use.

12.89 GNAT.Regexp (g-regexp.ads)
A simple implementation of regular expressions, using a subset of regular expression syntax copied from familiar Unix style utilities. This is the simplest of the three pattern matching packages provided, and is particularly suitable for ‘file globbing’ applications.

12.90 GNAT.Registry (g-regist.ads)
This is a high level binding to the Windows registry. It is possible to do simple things like reading a key value, creating a new key. For full registry API, but at a lower level of abstraction, refer to the Win32.Winreg package provided with the Win32Ada binding
12.91 GNAT.Regpat (g-regpat.ads)
A complete implementation of Unix-style regular expression matching, copied from the original V7 style regular expression library written in C by Henry Spencer (and binary compatible with this C library).

12.92 GNAT.Rewrite_Data (g-rewdat.ads)
A unit to rewrite on-the-fly string occurrences in a stream of data. The implementation has a very minimal memory footprint as the full content to be processed is not loaded into memory all at once. This makes this interface usable for large files or socket streams.

12.93 GNAT.Secondary_Stack_Info (g-sestin.ads)
Provides the capability to query the high water mark of the current task’s secondary stack.

12.94 GNAT.Semaphores (g-semaph.ads)
Provides classic counting and binary semaphores using protected types.

12.95 GNAT.Serial_Communications (g-sercom.ads)
Provides a simple interface to send and receive data over a serial port. This is only supported on GNU/Linux and Windows.

12.96 GNAT.SHA1 (g-sha1.ads)
Implements the SHA-1 Secure Hash Algorithm as described in FIPS PUB 180-3 and RFC 3174, and the HMAC-SHA1 message authentication function as described in RFC 2104 and FIPS PUB 198.

12.97 GNAT.SHA224 (g-sha224.ads)
Implements the SHA-224 Secure Hash Algorithm as described in FIPS PUB 180-3, and the HMAC-SHA224 message authentication function as described in RFC 2104 and FIPS PUB 198.

12.98 GNAT.SHA256 (g-sha256.ads)
Implements the SHA-256 Secure Hash Algorithm as described in FIPS PUB 180-3, and the HMAC-SHA256 message authentication function as described in RFC 2104 and FIPS PUB 198.

12.99 GNAT.SHA384 (g-sha384.ads)
Implements the SHA-384 Secure Hash Algorithm as described in FIPS PUB 180-3, and the HMAC-SHA384 message authentication function as described in RFC 2104 and FIPS PUB 198.
12.100 GNAT.SHA512 (g-sha512.ads)
Implements the SHA-512 Secure Hash Algorithm as described in FIPS PUB 180-3, and the HMAC-SHA512 message authentication function as described in RFC 2104 and FIPS PUB 198.

12.101 GNAT.Signals (g-signal.ads)
Provides the ability to manipulate the blocked status of signals on supported targets.

12.102 GNAT.Sockets (g-socket.ads)
A high level and portable interface to develop sockets based applications. This package is based on the sockets thin binding found in GNAT.Sockets.Thin. Currently GNAT.Sockets is implemented on all native GNAT ports and on VxWorks cross ports. It is not implemented for the LynxOS cross port.

12.103 GNAT.Source_Info (g-souinf.ads)
Provides subprograms that give access to source code information known at compile time, such as the current file name and line number. Also provides subprograms yielding the date and time of the current compilation (like the C macros __DATE__ and __TIME__)

12.104 GNAT.Spelling_Checker (g-speche.ads)
Provides a function for determining whether one string is a plausible near misspelling of another string.

12.105 GNAT.Spelling_Checker_Generic (g-spchge.ads)
Provides a generic function that can be instantiated with a string type for determining whether one string is a plausible near misspelling of another string.

12.106 GNAT.Spitbol.Patterns (g-spipat.ads)
A complete implementation of SNOBOL4 style pattern matching. This is the most elaborate of the pattern matching packages provided. It fully duplicates the SNOBOL4 dynamic pattern construction and matching capabilities, using the efficient algorithm developed by Robert Dewar for the SPITBOL system.

12.107 GNAT.Spitbol (g-spitbo.ads)
The top level package of the collection of SPITBOL-style functionality, this package provides basic SNOBOL4 string manipulation functions, such as Pad, Reverse, Trim, Substr capability, as well as a generic table function useful for constructing arbitrary mappings from strings in the style of the SNOBOL4 TABLE function.

12.108 GNAT.Spitbol.Table_Boolean (g-sptabo.ads)
A library level of instantiation of GNAT.Spitbol.Patterns.Table for type Standard.Boolean, giving an implementation of sets of string values.
12.109 GNAT.Spitbol.Table_Integer (g-sptain.ads)
A library level of instantiation of GNAT.Spitbol.Patterns.Table for type Standard.Integer, giving an implementation of maps from string to integer values.

12.110 GNAT.Spitbol.Table_VString (g-sptavs.ads)
A library level of instantiation of GNAT.Spitbol.Patterns.Table for a variable length string type, giving an implementation of general maps from strings to strings.

12.111 GNAT.SSE (g-sse.ads)
Root of a set of units aimed at offering Ada bindings to a subset of the Intel(r) Streaming SIMD Extensions with GNAT on the x86 family of targets. It exposes vector component types together with a general introduction to the binding contents and use.

12.112 GNAT.SSE.Vector_TYPES (g-ssvety.ads)
SSE vector types for use with SSE related intrinsics.

12.113 GNAT.String_Hash (g-strhas.ads)
Provides a generic hash function working on arrays of scalars. Both the scalar type and the hash result type are parameters.

12.114 GNAT.Strings (g-string.ads)
Common String access types and related subprograms. Basically it defines a string access and an array of string access types.

12.115 GNAT.String_Split (g-strspl.ads)
Useful string manipulation routines: given a set of separators, split a string wherever the separators appear, and provide direct access to the resulting slices. This package is instantiated from GNAT.Array_Split.

12.116 GNAT.Table (g-table.ads)
A generic package providing a single dimension array abstraction where the length of the array can be dynamically modified.
This package provides a facility similar to that of GNAT.Dynamic_Tables, except that this package declares a single instance of the table type, while an instantiation of GNAT.Dynamic_Tables creates a type that can be used to define dynamic instances of the table.

12.117 GNAT.Task_Lock (g-tasloc.ads)
A very simple facility for locking and unlocking sections of code using a single global task lock. Appropriate for use in situations where contention between tasks is very rarely expected.
12.118 GNAT.Time_Stamp (g-timsta.ads)
Provides a simple function that returns a string YYYY-MM-DD HH:MM:SS.SS that represents the current date and time in ISO 8601 format. This is a very simple routine with minimal code and there are no dependencies on any other unit.

12.119 GNAT.Threads (g-thread.ads)
Provides facilities for dealing with foreign threads which need to be known by the GNAT run-time system. Consult the documentation of this package for further details if your program has threads that are created by a non-Ada environment which then accesses Ada code.

12.120 GNAT.Traceback (g-traceb.ads)
Provides a facility for obtaining non-symbolic traceback information, useful in various debugging situations.

12.121 GNAT.Traceback.Symbolic (g-trasym.ads)

12.122 GNAT.UTF_32 (g-utf_32.ads)
This is a package intended to be used in conjunction with the Wide_Character type in Ada 95 and the Wide_Wide_Character type in Ada 2005 (available in GNAT in Ada 2005 mode). This package contains Unicode categorization routines, as well as lexical categorization routines corresponding to the Ada 2005 lexical rules for identifiers and strings, and also a lower case to upper case fold routine corresponding to the Ada 2005 rules for identifier equivalence.

12.123 GNAT.UTF_32_Spelling_Checker (g-u3spch.ads)
Provides a function for determining whether one wide wide string is a plausible near misspelling of another wide wide string, where the strings are represented using the UTF_32_String type defined in System.Wch_Cnv.

12.124 GNAT.Wide_Spelling_Checker (g-wispch.ads)
Provides a function for determining whether one wide string is a plausible near misspelling of another wide string.

12.125 GNAT.Wide_String_Split (g-wistsp.ads)
Useful wide string manipulation routines: given a set of separators, split a wide string wherever the separators appear, and provide direct access to the resulting slices. This package is instantiated from GNAT.Array_Split.

12.126 GNAT.Wide_Wide_Spelling_Checker (g-zspche.ads)
Provides a function for determining whether one wide wide string is a plausible near misspelling of another wide wide string.
12.127 GNAT.Wide_Wide_String_Split (g-zistsp.ads)
Useful wide wide string manipulation routines: given a set of separators, split a wide wide string wherever the separators appear, and provide direct access to the resulting slices. This package is instantiated from GNAT.Array_Split.

12.128 Interfaces.C.Extensions (i-cexten.ads)
This package contains additional C-related definitions, intended for use with either manually or automatically generated bindings to C libraries.

12.129 Interfaces.C.Streams (i-cstrea.ads)
This package is a binding for the most commonly used operations on C streams.

12.130 Interfaces.PackedDecimal (i-pacdec.ads)
This package provides a set of routines for conversions to and from a packed decimal format compatible with that used on IBM mainframes.

12.131 Interfaces.VxWorks (i-vxwork.ads)
This package provides a limited binding to the VxWorks API. In particular, it interfaces with the VxWorks hardware interrupt facilities.

12.132 Interfaces.VxWorks.Int_Connection (i-vxinco.ads)
This package provides a way for users to replace the use of intConnect() with a custom routine for installing interrupt handlers.

12.133 Interfaces.VxWorks.IO (i-vxwoio.ads)
This package provides a binding to the ioctl (IO/Control) function of VxWorks, defining a set of option values and function codes. A particular use of this package is to enable the use of Get_Immediate under VxWorks.

12.134 System.Address_Image (s-addima.ads)
This function provides a useful debugging function that gives an (implementation dependent) string which identifies an address.

12.135 System.Assertions (s-assert.ads)
This package provides the declaration of the exception raised by a run-time assertion failure, as well as the routine that is used internally to raise this assertion.

12.136 System.Atomic_Counters (s-atocou.ads)
This package provides the declaration of an atomic counter type, together with efficient routines (using hardware synchronization primitives) for incrementing, decrementing, and testing of these counters. This package is implemented on most targets, including all Alpha, AARCH64, ARM, ia64, PowerPC, SPARC V9, x86, and x86_64 platforms.
12.137 System.Memory (s-memory.ads)
This package provides the interface to the low level routines used by the generated code for allocation and freeing storage for the default storage pool (analogous to the C routines malloc and free). It also provides a reallocation interface analogous to the C routine realloc. The body of this unit may be modified to provide alternative allocation mechanisms for the default pool, and in addition, direct calls to this unit may be made for low level allocation uses (for example see the body of GNAT.Tables).

12.138 System.Multiprocessors (s-multip.ads)
This is an Ada 2012 unit defined in the Ada 2012 Reference Manual, but in GNAT we also make it available in Ada 95 and Ada 2005 (where it is technically an implementation-defined addition).

12.139 System.Multiprocessors.Dispatching_Domains (s-mudido.ads)
This is an Ada 2012 unit defined in the Ada 2012 Reference Manual, but in GNAT we also make it available in Ada 95 and Ada 2005 (where it is technically an implementation-defined addition).

12.140 System.Partition_Interface (s-parint.ads)
This package provides facilities for partition interfacing. It is used primarily in a distribution context when using Annex E with GLADE.

12.141 System.Pool_Global (s-pooglo.ads)
This package provides a storage pool that is equivalent to the default storage pool used for access types for which no pool is specifically declared. It uses malloc/free to allocate/free and does not attempt to do any automatic reclamation.

12.142 System.Pool_Local (s-pooloc.ads)
This package provides a storage pool that is intended for use with locally defined access types. It uses malloc/free for allocate/free, and maintains a list of allocated blocks, so that all storage allocated for the pool can be freed automatically when the pool is finalized.

12.143 System.Restrictions (s-restri.ads)
This package provides facilities for accessing at run time the status of restrictions specified at compile time for the partition. Information is available both with regard to actual restrictions specified, and with regard to compiler determined information on which restrictions are violated by one or more packages in the partition.

12.144 System.Rident (s-rident.ads)
This package provides definitions of the restrictions identifiers supported by GNAT, and also the format of the restrictions provided in package System.Restrictions. It is not nor-
nally necessary to with this generic package since the necessary instantiation is included in package System.Restrictions.

12.145 System.Strings.Stream_Ops (s-ststop.ads)
This package provides a set of stream subprograms for standard string types. It is intended primarily to support implicit use of such subprograms when stream attributes are applied to string types, but the subprograms in this package can be used directly by application programs.

12.146 System.Unsigned_Types (s-unstyp.ads)
This package contains definitions of standard unsigned types that correspond in size to the standard signed types declared in Standard, and (unlike the types in Interfaces) have corresponding names. It also contains some related definitions for other specialized types used by the compiler in connection with packed array types.

12.147 System.Wch_Cnv (s-wchcnv.ads)
This package provides routines for converting between wide and wide wide characters and a representation as a value of type Standard.String, using a specified wide character encoding method. It uses definitions in package System.Wch_Con.

12.148 System.Wch_Con (s-wchcon.ads)
This package provides definitions and descriptions of the various methods used for encoding wide characters in ordinary strings. These definitions are used by the package System.Wch_Cnv.
13 Interfacing to Other Languages

The facilities in Annex B of the Ada Reference Manual are fully implemented in GNAT, and in addition, a full interface to C++ is provided.

13.1 Interfacing to C

Interfacing to C with GNAT can use one of two approaches:
* The types in the package Interfaces.C may be used.
* Standard Ada types may be used directly. This may be less portable to other compilers, but will work on all GNAT compilers, which guarantee correspondence between the C and Ada types.

Pragma Convention C may be applied to Ada types, but mostly has no effect, since this is the default. The following table shows the correspondence between Ada scalar types and the corresponding C types.

<table>
<thead>
<tr>
<th>Ada Type</th>
<th>C Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Integer</td>
<td>int</td>
</tr>
<tr>
<td>Short_Integer</td>
<td>short</td>
</tr>
<tr>
<td>Short_Short_Integer</td>
<td>signed char</td>
</tr>
<tr>
<td>Long_Integer</td>
<td>long</td>
</tr>
<tr>
<td>Long_Long_Integer</td>
<td>long long</td>
</tr>
<tr>
<td>Short_Float</td>
<td>float</td>
</tr>
<tr>
<td>Float</td>
<td>float</td>
</tr>
<tr>
<td>Long_Float</td>
<td>double</td>
</tr>
<tr>
<td>Long_Long_Float</td>
<td></td>
</tr>
</tbody>
</table>

Additionally, there are the following general correspondences between Ada and C types:
* Ada enumeration types map to C enumeration types directly if pragma Convention C is specified, which causes them to have a length of 32 bits, except for boolean types which map to C99 bool and for which the length is 8 bits. Without pragma Convention C, Ada enumeration types map to 8, 16, or 32 bits (i.e., C types signed char, short, int, respectively) depending on the number of values passed. This is the only case in which pragma Convention C affects the representation of an Ada type.
* Ada access types map to C pointers, except for the case of pointers to unconstrained types in Ada, which have no direct C equivalent.
* Ada arrays map directly to C arrays.
* Ada records map directly to C structures.
* Packed Ada records map to C structures where all members are bit fields of the length corresponding to the type’Size value in Ada.

## 13.2 Interfacing to C++

The interface to C++ makes use of the following pragmas, which are primarily intended to be constructed automatically using a binding generator tool, although it is possible to construct them by hand.

Using these pragmas it is possible to achieve complete inter-operability between Ada tagged types and C++ class definitions. See [Implementation Defined Pragmas], page 4, for more details.

**pragma CPP_Class ([Entity =>] LOCAL_NAME)**

The argument denotes an entity in the current declarative region that is declared as a tagged or untagged record type. It indicates that the type corresponds to an externally declared C++ class type, and is to be laid out the same way that C++ would lay out the type.

Note: Pragma CPP_Class is currently obsolete. It is supported for backward compatibility but its functionality is available using pragma Import with Convention = CPP.

**pragma CPP_Constructor ([Entity =>] LOCAL_NAME)**

This pragma identifies an imported function (imported in the usual way with pragma Import) as corresponding to a C++ constructor.

A few restrictions are placed on the use of the Access attribute in conjunction with subprograms subject to convention CPP: the attribute may be used neither on primitive operations of a tagged record type with convention CPP, imported or not, nor on subprograms imported with pragma CPP_Constructor.

In addition, C++ exceptions are propagated and can be handled in an others choice of an exception handler. The corresponding Ada occurrence has no message, and the simple name of the exception identity contains Foreign_Exception. Finalization and awaiting dependent tasks works properly when such foreign exceptions are propagated.

It is also possible to import a C++ exception using the following syntax:

```ada
LOCAL_NAME : exception;
pragma Import (Cpp,
  [Entity =>] LOCAL_NAME,
  [External_Name =>] static_string_EXPRESSION);
```

The External_Name is the name of the C++ RTTI symbol. You can then cover a specific C++ exception in an exception handler.

## 13.3 Interfacing to COBOL

Interfacing to COBOL is achieved as described in section B.4 of the Ada Reference Manual.
13.4 Interfacing to Fortran

Interfacing to Fortran is achieved as described in section B.5 of the Ada Reference Manual. The pragma `Convention Fortran`, applied to a multi-dimensional array causes the array to be stored in column-major order as required for convenient interface to Fortran.

13.5 Interfacing to non-GNAT Ada code

It is possible to specify the convention `Ada` in a pragma `Import` or pragma `Export`. However, this refers to the calling conventions used by GNAT, which may or may not be similar enough to those used by some other Ada 83 / Ada 95 / Ada 2005 compiler to allow interoperation. If arguments types are kept simple, and if the foreign compiler generally follows system calling conventions, then it may be possible to integrate files compiled by other Ada compilers, provided that the elaboration issues are adequately addressed (for example by eliminating the need for any load time elaboration).

In particular, GNAT running on VMS is designed to be highly compatible with the DEC Ada 83 compiler, so this is one case in which it is possible to import foreign units of this type, provided that the data items passed are restricted to simple scalar values or simple record types without variants, or simple array types with fixed bounds.
14 Specialized Needs Annexes

Ada 95, Ada 2005, and Ada 2012 define a number of Specialized Needs Annexes, which are not required in all implementations. However, as described in this chapter, GNAT implements all of these annexes:

‘Systems Programming (Annex C)’
The Systems Programming Annex is fully implemented.

‘Real-Time Systems (Annex D)’
The Real-Time Systems Annex is fully implemented.

‘Distributed Systems (Annex E)’
Stub generation is fully implemented in the GNAT compiler. In addition, a complete compatible PCS is available as part of the GLADE system, a separate product. When the two products are used in conjunction, this annex is fully implemented.

‘Information Systems (Annex F)’
The Information Systems annex is fully implemented.

‘Numerics (Annex G)’
The Numerics Annex is fully implemented.

‘Safety and Security / High-Integrity Systems (Annex H)’
The Safety and Security Annex (termed the High-Integrity Systems Annex in Ada 2005) is fully implemented.
15 Implementation of Specific Ada Features

This chapter describes the GNAT implementation of several Ada language facilities.

15.1 Machine Code Insertions

Package Machine_Code provides machine code support as described in the Ada Reference Manual in two separate forms:

* Machine code statements, consisting of qualified expressions that fit the requirements of RM section 13.8.
* An intrinsic callable procedure, providing an alternative mechanism of including machine instructions in a subprogram.

The two features are similar, and both are closely related to the mechanism provided by the asm instruction in the GNU C compiler. Full understanding and use of the facilities in this package requires understanding the asm instruction, see the section on Extended Asm in Using the GNU Compiler Collection (GCC).

Calls to the function Asm and the procedure Asm have identical semantic restrictions and effects as described below. Both are provided so that the procedure call can be used as a statement, and the function call can be used to form a code statement.

Consider this C asm instruction:

```
asm ("fsinx %1 %0" : "=f" (result) : "f" (angle));
```

The equivalent can be written for GNAT as:

```
Asm ("fsinx %1 %0",
     My_Float'Asm_Output ("=f", result),
     My_Float'Asm_Input ("f", angle));
```

The first argument to Asm is the assembler template, and is identical to what is used in GNU C. This string must be a static expression. The second argument is the output operand list. It is either a single Asm_Output attribute reference, or a list of such references enclosed in parentheses (technically an array aggregate of such references).

The second argument of my_float'Asm_Output functions as though it were an out parameter, which is a little curious, but all names have the form of expressions, so there is no syntactic irregularity, even though normally functions would not be permitted out parameters. The third argument is the list of input operands. It is either a single Asm_Input attribute reference, or a list of such references enclosed in parentheses (technically an array aggregate of such references).
The `Asm_Input` attribute denotes a function that takes two parameters. The first is a string, the second is an expression of the type designated by the prefix. The first (string) argument is required to be a static expression, and is the constraint for the parameter, (e.g., what kind of register is required). The second argument is the value to be used as the input argument. The possible values for the constraint are the same as those used in the RTL, and are dependent on the configuration file used to built the GCC back end. No support is provided for GNU C’s symbolic names for input parameters.

If there are no input operands, this argument may either be omitted, or explicitly given as `No_Input_Operands`. The fourth argument, not present in the above example, is a list of register names, called the ‘clobber’ argument. This argument, if given, must be a static string expression, and is a space or comma separated list of names of registers that must be considered destroyed as a result of the `Asm` call. If this argument is the null string (the default value), then the code generator assumes that no additional registers are destroyed. In addition to registers, the special clobbers `memory` and `cc` as described in the GNU C docs are both supported.

The fifth argument, not present in the above example, called the ‘volatile’ argument, is by default `False`. It can be set to the literal value `True` to indicate to the code generator that all optimizations with respect to the instruction specified should be suppressed, and in particular an instruction that has outputs will still be generated, even if none of the outputs are used. See *Using the GNU Compiler Collection (GCC)* for the full description. Generally it is strongly advisable to use Volatile for any ASM statement that is missing either input or output operands or to avoid unwanted optimizations. A warning is generated if this advice is not followed.

No support is provided for GNU C’s `asm goto` feature.

The `Asm` subprograms may be used in two ways. First the procedure forms can be used anywhere a procedure call would be valid, and correspond to what the RM calls ‘intrinsic’ routines. Such calls can be used to intersperse machine instructions with other Ada statements. Second, the function forms, which return a dummy value of the limited private type `Asm_Insn`, can be used in code statements, and indeed this is the only context where such calls are allowed. Code statements appear as aggregates of the form:

- `Asm_Insn'(Asm (...));`
- `Asm_Insn'(Asm_Volatile (...));`

In accordance with RM rules, such code statements are allowed only within subprograms whose entire body consists of such statements. It is not permissible to intermix such statements with other Ada statements.

Typically the form using intrinsic procedure calls is more convenient and more flexible. The code statement form is provided to meet the RM suggestion that such a facility should be made available. The following is the exact syntax of the call to `Asm`. As usual, if named notation is used, the arguments may be given in arbitrary order, following the normal rules for use of positional and named arguments:

- `ASM_CALL ::= Asm (  
  [Template =>] static_string_EXPRESSION 
  [, [Outputs =>] OUTPUT_OPERAND_LIST ]  
  [, [Inputs =>] INPUT_OPERAND_LIST ]  
  [, [Clobber =>] static_string_EXPRESSION ]  
)`
Chapter 15: Implementation of Specific Ada Features

Chapter 15: Implementation of Specific Ada Features

[,[Volatile =>] static_boolean_EXPRESSION] )

OUTPUT_OPERAND_LIST ::= 
  [PREFIX.]No_Output_Operands 
  | OUTPUT_OPERAND_ATTRIBUTE 
  | (OUTPUT_OPERAND_ATTRIBUTE {,OUTPUT_OPERAND_ATTRIBUTE})

OUTPUT_OPERAND_ATTRIBUTE ::= 
  SUBTYPE_MARK’Asm_Output (static_string_EXPRESSION, NAME)

INPUT_OPERAND_LIST ::= 
  [PREFIX.]No_Input_Operands 
  | INPUT_OPERAND_ATTRIBUTE 
  | (INPUT_OPERAND_ATTRIBUTE {,INPUT_OPERAND_ATTRIBUTE})

INPUT_OPERAND_ATTRIBUTE ::= 
  SUBTYPE_MARK’Asm_Input (static_string_EXPRESSION, EXPRESSION)

The identifiers No_Input_Operands and No_Output_Operands are declared in the package Machine_Code and must be referenced according to normal visibility rules. In particular if there is no use clause for this package, then appropriate package name qualification is required.

15.2 GNAT Implementation of Tasking

This chapter outlines the basic GNAT approach to tasking (in particular, a multi-layered library for portability) and discusses issues related to compliance with the Real-Time Systems Annex.

15.2.1 Mapping Ada Tasks onto the Underlying Kernel Threads

GNAT’s run-time support comprises two layers:
  * GNARL (GNAT Run-time Layer)
  * GNULL (GNAT Low-level Library)

In GNAT, Ada’s tasking services rely on a platform and OS independent layer known as GNARL. This code is responsible for implementing the correct semantics of Ada’s task creation, rendezvous, protected operations etc.

GNARL decomposes Ada’s tasking semantics into simpler lower level operations such as create a thread, set the priority of a thread, yield, create a lock, lock/unlock, etc. The spec for these low-level operations constitutes GNULLI, the GNULL Interface. This interface is directly inspired from the POSIX real-time API.

If the underlying executive or OS implements the POSIX standard faithfully, the GNULL Interface maps as is to the services offered by the underlying kernel. Otherwise, some target dependent glue code maps the services offered by the underlying kernel to the semantics expected by GNARL.

Whatever the underlying OS (VxWorks, UNIX, Windows, etc.) the key point is that each Ada task is mapped on a thread in the underlying kernel. For example, in the case of VxWorks, one Ada task = one VxWorks task.
In addition Ada task priorities map onto the underlying thread priorities. Mapping Ada
tasks onto the underlying kernel threads has several advantages:

* The underlying scheduler is used to schedule the Ada tasks. This makes Ada tasks as
efficient as kernel threads from a scheduling standpoint.
* Interaction with code written in C containing threads is eased since at the lowest level
Ada tasks and C threads map onto the same underlying kernel concept.
* When an Ada task is blocked during I/O the remaining Ada tasks are able to proceed.
* On multiprocessor systems Ada tasks can execute in parallel.

Some threads libraries offer a mechanism to fork a new process, with the child process
duplicating the threads from the parent. GNAT does not support this functionality when
the parent contains more than one task.

15.2.2 Ensuring Compliance with the Real-Time Annex

Although mapping Ada tasks onto the underlying threads has significant advantages, it does
create some complications when it comes to respecting the scheduling semantics specified
in the real-time annex (Annex D).

For instance the Annex D requirement for the FIFO_Within_Priorities scheduling policy
states:

‘When the active priority of a ready task that is not running changes,
or the setting of its base priority takes effect, the task is removed from
the ready queue for its old active priority and is added at the tail of the
ready queue for its new active priority, except in the case where the active
priority is lowered due to the loss of inherited priority, in which case the
task is added at the head of the ready queue for its new active priority.’

While most kernels do put tasks at the end of the priority queue when a task changes its
priority, (which respects the main FIFO_Within_Priorities requirement), almost none keep
a thread at the beginning of its priority queue when its priority drops from the loss of
inherited priority.

As a result most vendors have provided incomplete Annex D implementations.

The GNAT run-time, has a nice cooperative solution to this problem which ensures that
accurate FIFO_Within_Priorities semantics are respected.

The principle is as follows. When an Ada task T is about to start running, it checks whether
some other Ada task R with the same priority as T has been suspended due to the loss
of priority inheritance. If this is the case, T yields and is placed at the end of its priority
queue. When R arrives at the front of the queue it executes.

Note that this simple scheme preserves the relative order of the tasks that were ready to
execute in the priority queue where R has been placed at the end.

15.2.3 Support for Locking Policies

This section specifies which policies specified by pragma Locking_Policy are supported on
which platforms.

GNAT supports the standard Ceiling_Locking policy, and the implementation defined
Inheritance_Locking and Concurrent_Readers_Locking policies.
Ceiling_Locking is supported on all platforms if the operating system supports it. In particular, Ceiling_Locking is not supported on VxWorks. Inheritance_Locking is supported on Linux, Darwin (Mac OS X), LynxOS 178, and VxWorks. Concurrent_Readers_Locking is supported on Linux.

Notes about Ceiling_Locking on Linux: If the process is running as ‘root’, ceiling locking is used. If the capabilities facility is installed (“sudo apt-get –assume-yes install libcap-dev” on Ubuntu, for example), and the program is linked against that library (“-largs -lcap”), and the executable file has the cap_sys_nice capability (“sudo /sbin/setcap cap_sys_nice=ep executable_file_name”), then ceiling locking is used. Otherwise, the Ceiling_Locking policy is ignored.

15.3 GNAT Implementation of Shared Passive Packages

GNAT fully implements the pragma Shared_Passive for the purpose of designating shared passive packages. This allows the use of passive partitions in the context described in the Ada Reference Manual; i.e., for communication between separate partitions of a distributed application using the features in Annex E.

However, the implementation approach used by GNAT provides for more extensive usage as follows:

‘Communication between separate programs’

This allows separate programs to access the data in passive partitions, using protected objects for synchronization where needed. The only requirement is that the two programs have a common shared file system. It is even possible for programs running on different machines with different architectures (e.g., different endianness) to communicate via the data in a passive partition.

‘Persistence between program runs’

The data in a passive package can persist from one run of a program to another, so that a later program sees the final values stored by a previous run of the same program.

The implementation approach used is to store the data in files. A separate stream file is created for each object in the package, and an access to an object causes the corresponding file to be read or written.

The environment variable SHARED_MEMORY_DIRECTORY should be set to the directory to be used for these files. The files in this directory have names that correspond to their fully qualified names. For example, if we have the package

```
package X is
  pragma Shared_Passive (X);
  Y : Integer;
  Z : Float;
end X;
```

and the environment variable is set to /stemp/, then the files created will have the names:
```
/stemp/x.y
/stemp/x.z
```

These files are created when a value is initially written to the object, and the files are retained until manually deleted. This provides the persistence semantics. If no file exists,
Chapter 15: Implementation of Specific Ada Features

it means that no partition has assigned a value to the variable; in this case the initial value declared in the package will be used. This model ensures that there are no issues in synchronizing the elaboration process, since elaboration of passive packages elaborates the initial values, but does not create the files.

The files are written using normal Stream IO access. If you want to be able to communicate between programs or partitions running on different architectures, then you should use the XDR versions of the stream attribute routines, since these are architecture independent.

If active synchronization is required for access to the variables in the shared passive package, then as described in the Ada Reference Manual, the package may contain protected objects used for this purpose. In this case a lock file (whose name is ___lock, with three underscores) is created in the shared memory directory.

This is used to provide the required locking semantics for proper protected object synchronization.

15.4 Code Generation for Array Aggregates

Aggregates have a rich syntax and allow the user to specify the values of complex data structures by means of a single construct. As a result, the code generated for aggregates can be quite complex and involve loops, case statements and multiple assignments. In the simplest cases, however, the compiler will recognize aggregates whose components and constraints are fully static, and in those cases the compiler will generate little or no executable code. The following is an outline of the code that GNAT generates for various aggregate constructs. For further details, you will find it useful to examine the output produced by the -gnatG flag to see the expanded source that is input to the code generator. You may also want to examine the assembly code generated at various levels of optimization.

The code generated for aggregates depends on the context, the component values, and the type. In the context of an object declaration the code generated is generally simpler than in the case of an assignment. As a general rule, static component values and static subtypes also lead to simpler code.

15.4.1 Static constant aggregates with static bounds

For the declarations:

```ada
type One_Dim is array (1..10) of integer;
ar0 : constant One_Dim := (1, 2, 3, 4, 5, 6, 7, 8, 9, 0);
```

GNAT generates no executable code: the constant ar0 is placed in static memory. The same is true for constant aggregates with named associations:

```ada
Cr1 : constant One_Dim := (4 => 16, 2 => 4, 3 => 9, 1 => 1, 5 .. 10 => 0);
Cr3 : constant One_Dim := (others => 7777);
```

The same is true for multidimensional constant arrays such as:

```ada
type two_dim is array (1..3, 1..3) of integer;
Unit : constant two_dim := ( (1,0,0), (0,1,0), (0,0,1));
```

The same is true for arrays of one-dimensional arrays: the following are static:

```ada
type ar1b is array (1..3) of boolean;
type ar_ar is array (1..3) of ar1b;
None : constant ar1b := (others => false); -- fully static
```
None2 : constant ar_ar := (1..3 => None); -- fully static

However, for multidimensional aggregates with named associations, GNAT will generate assignments and loops, even if all associations are static. The following two declarations generate a loop for the first dimension, and individual component assignments for the second dimension:

Zero1: constant two_dim := (1..3 => (1..3 => 0));
Zero2: constant two_dim := (others => (others => 0));

15.4.2 Constant aggregates with unconstrained nominal types

In such cases the aggregate itself establishes the subtype, so that associations with others cannot be used. GNAT determines the bounds for the actual subtype of the aggregate, and allocates the aggregate statically as well. No code is generated for the following:

type One_Unc is array (natural range <>) of integer;
Cr_Unc : constant One_Unc := (12, 24, 36);

15.4.3 Aggregates with static bounds

In all previous examples the aggregate was the initial (and immutable) value of a constant. If the aggregate initializes a variable, then code is generated for it as a combination of individual assignments and loops over the target object. The declarations

Cr_Var1 : One_Dim := (2, 5, 7, 11, 0, 0, 0, 0, 0, 0);
Cr_Var2 : One_Dim := (others > -1);

generate the equivalent of

Cr_Var1 (1) := 2;
Cr_Var1 (2) := 3;
Cr_Var1 (3) := 5;
Cr_Var1 (4) := 11;

for I in Cr_Var2’range loop
    Cr_Var2 (I) := -1;
end loop;

15.4.4 Aggregates with nonstatic bounds

If the bounds of the aggregate are not statically compatible with the bounds of the nominal subtype of the target, then constraint checks have to be generated on the bounds. For a multidimensional array, constraint checks may have to be applied to sub-arrays individually, if they do not have statically compatible subtypes.

15.4.5 Aggregates in assignment statements

In general, aggregate assignment requires the construction of a temporary, and a copy from the temporary to the target of the assignment. This is because it is not always possible to convert the assignment into a series of individual component assignments. For example, consider the simple case:

A := (A(2), A(1));

This cannot be converted into:
A(1) := A(2);
A(2) := A(1);

So the aggregate has to be built first in a separate location, and then copied into the target. GNAT recognizes simple cases where this intermediate step is not required, and the assignments can be performed in place, directly into the target. The following sufficient criteria are applied:

* The bounds of the aggregate are static, and the associations are static.
* The components of the aggregate are static constants, names of simple variables that are not renamings, or expressions not involving indexed components whose operands obey these rules.

If any of these conditions are violated, the aggregate will be built in a temporary (created either by the front-end or the code generator) and then that temporary will be copied onto the target.

15.5 The Size of Discriminated Records with Default Discriminants

If a discriminated type \( T \) has discriminants with default values, it is possible to declare an object of this type without providing an explicit constraint:

```ada
type Size is range 1..100;

type Rec (D : Size := 15) is record
  Name : String (1..D);
end T;

Word : Rec;
```

Such an object is said to be ‘unconstrained’. The discriminant of the object can be modified by a full assignment to the object, as long as it preserves the relation between the value of the discriminant, and the value of the components that depend on it:

```ada
Word := (3, "yes");

Word := (5, "maybe");

Word := (5, "no"); -- raises Constraint_Error
```

In order to support this behavior efficiently, an unconstrained object is given the maximum size that any value of the type requires. In the case above, \( \text{Word} \) has storage for the discriminant and for a \text{String} of length 100. It is important to note that unconstrained objects do not require dynamic allocation. It would be an improper implementation to place on the heap those components whose size depends on discriminants. (This improper implementation was used by some Ada83 compilers, where the \text{Name} component above would have been stored as a pointer to a dynamic string). Following the principle that dynamic storage management should never be introduced implicitly, an Ada compiler should reserve the full size for an unconstrained declared object, and place it on the stack.

This maximum size approach has been a source of surprise to some users, who expect the default values of the discriminants to determine the size reserved for an unconstrained
object: “If the default is 15, why should the object occupy a larger size?” The answer, of course, is that the discriminant may be later modified, and its full range of values must be taken into account. This is why the declaration:

```ada
   type Rec (D : Positive := 15) is record
      Name : String (1..D);
   end record;
   Too_Large : Rec;
```

is flagged by the compiler with a warning: an attempt to create `Too_Large` will raise `Storage_Error`, because the required size includes `Positive'Last` bytes. As the first example indicates, the proper approach is to declare an index type of ‘reasonable’ range so that unconstrained objects are not too large.

One final wrinkle: if the object is declared to be aliased, or if it is created in the heap by means of an allocator, then it is ‘not’ unconstrained: it is constrained by the default values of the discriminants, and those values cannot be modified by full assignment. This is because in the presence of aliasing all views of the object (which may be manipulated by different tasks, say) must be consistent, so it is imperative that the object, once created, remain invariant.

### 15.6 Image Values For Nonscalar Types

Ada 2022 defines the `Image`, `Wide_Image`, and `Wide_Wide` image attributes for nonscalar types; earlier Ada versions defined these attributes only for scalar types. Ada RM 4.10 provides some general guidance regarding the default implementation of these attributes and the GNAT compiler follows that guidance. However, beyond that the precise details of the image text generated in these cases are deliberately not documented and are subject to change. In particular, users should not rely on formatting details (such as spaces or line breaking), record field order, image values for access types, image values for types that have ancestor or subcomponent types declared in non-Ada2022 code, image values for predefined types, or the compiler’s choices regarding the implementation permissions described in Ada RM 4.10. This list is not intended to be exhaustive. If more precise control of image text is required for some type `T`, then `T'Put_Image` should be explicitly specified.

### 15.7 Strict Conformance to the Ada Reference Manual

The dynamic semantics defined by the Ada Reference Manual impose a set of run-time checks to be generated. By default, the GNAT compiler will insert many run-time checks into the compiled code, including most of those required by the Ada Reference Manual. However, there are two checks that are not enabled in the default mode for efficiency reasons: checks for access before elaboration on subprogram calls, and stack overflow checking (most operating systems do not perform this check by default).

Strict conformance to the Ada Reference Manual can be achieved by adding two compiler options for dynamic checks for access-before-elaboration on subprogram calls and generic instantiations (`-gnatE`), and stack overflow checking (`-fstack-check`).

Note that the result of a floating point arithmetic operation in overflow and invalid situations, when the `Machine_Overflows` attribute of the result type is `False`, is to generate IEEE NaN and infinite values. This is the case for machines compliant with the IEEE
floating-point standard, but on machines that are not fully compliant with this standard, such as Alpha, the ‘-mieee’ compiler flag must be used for achieving IEEE confirming behavior (although at the cost of a significant performance penalty), so infinite and NaN values are properly generated.
16 Implementation of Ada 2012 Features

This chapter contains a complete list of Ada 2012 features that have been implemented. Generally, these features are only available if the ‘-gnat12’ (Ada 2012 features enabled) option is set, which is the default behavior, or if the configuration pragma Ada_2012 is used.

However, new pragmas, attributes, and restrictions are unconditionally available, since the Ada 95 standard allows the addition of new pragmas, attributes, and restrictions (there are exceptions, which are documented in the individual descriptions), and also certain packages were made available in earlier versions of Ada.

An ISO date (YYYY-MM-DD) appears in parentheses on the description line. This date shows the implementation date of the feature. Any wavefront subsequent to this date will contain the indicated feature, as will any subsequent releases. A date of 0000-00-00 means that GNAT has always implemented the feature, or implemented it as soon as it appeared as a binding interpretation.

Each feature corresponds to an Ada Issue (‘AI’) approved by the Ada standardization group (ISO/IEC JTC1/SC22/WG9) for inclusion in Ada 2012. The features are ordered based on the relevant sections of the Ada Reference Manual (“RM”). When a given AI relates to multiple points in the RM, the earliest is used.

A complete description of the AIs may be found in http://www.ada-auth.org/ai05-summary.html.

* ‘AI-0176 Quantified expressions (2010-09-29)’
  Both universally and existentially quantified expressions are implemented. They use the new syntax for iterators proposed in AI05-139-2, as well as the standard Ada loop syntax.
  RM References: 1.01.04 (12) 2.09 (2/2) 4.04 (7) 4.05.09 (0)

* ‘AI-0079 Allow other_format characters in source (2010-07-10)’
  Wide characters in the unicode category ‘other_format’ are now allowed in source programs between tokens, but not within a token such as an identifier.
  RM References: 2.01 (4/2) 2.02 (7)

* ‘AI-0091 Do not allow other_format in identifiers (0000-00-00)’
  Wide characters in the unicode category ‘other_format’ are not permitted within an identifier, since this can be a security problem. The error message for this case has been improved to be more specific, but GNAT has never allowed such characters to appear in identifiers.
  RM References: 2.03 (3.1/2) 2.03 (4/2) 2.03 (5/2) 2.03 (5.1/2) 2.03 (5.2/2) 2.03 (5.3/2) 2.09 (2/2)

* ‘AI-0100 Placement of pragmas (2010-07-01)’
  This AI is an earlier version of AI-163. It simplifies the rules for legal placement of pragmas. In the case of lists that allow pragmas, if the list may have no elements, then the list may consist solely of pragmas.
  RM References: 2.08 (7)
* ‘AI-0163 Pragmas in place of null (2010-07-01)’

A statement sequence may be composed entirely of pragmas. It is no longer necessary to add a dummy `null` statement to make the sequence legal.

RM References: 2.08 (7) 2.08 (16)

* ‘AI-0080 ‘View of’ not needed if clear from context (0000-00-00)’

This is an editorial change only, described as non-testable in the AI.

RM References: 3.01 (7)

* ‘AI-0183 Aspect specifications (2010-08-16)’

Aspect specifications have been fully implemented except for pre and post-conditions, and type invariants, which have their own separate AI’s. All forms of declarations listed in the AI are supported. The following is a list of the aspects supported (with GNAT implementation aspects marked)

<table>
<thead>
<tr>
<th>Supported Aspect</th>
<th>Source</th>
</tr>
</thead>
<tbody>
<tr>
<td>Ada_2005</td>
<td>– GNAT</td>
</tr>
<tr>
<td>Ada_2012</td>
<td>– GNAT</td>
</tr>
<tr>
<td>Address</td>
<td></td>
</tr>
<tr>
<td>Alignment</td>
<td></td>
</tr>
<tr>
<td>Atomic</td>
<td></td>
</tr>
<tr>
<td>Atomic.Components</td>
<td></td>
</tr>
<tr>
<td>Bit.Order</td>
<td></td>
</tr>
<tr>
<td>Component_Size</td>
<td></td>
</tr>
<tr>
<td>Contract.Cases</td>
<td>– GNAT</td>
</tr>
<tr>
<td>Discard.Names</td>
<td></td>
</tr>
<tr>
<td>External.Tag</td>
<td></td>
</tr>
<tr>
<td>Favor_Top_Level</td>
<td>– GNAT</td>
</tr>
<tr>
<td>Inline</td>
<td></td>
</tr>
<tr>
<td>Inline_Always</td>
<td>– GNAT</td>
</tr>
<tr>
<td>Invariant</td>
<td>– GNAT</td>
</tr>
</tbody>
</table>
Machine_Radix

No_Return

Object_Size – GNAT

Pack

Persistent_BSS – GNAT

Post

Pre

Predicate

Preelaborable_Initialization

Pure_Function – GNAT

Remote_Access_Type – GNAT

Shared – GNAT

Size

Storage_Pool

Storage_Size

Stream_Size

Suppress

Suppress_Debug_Info – GNAT

Test_Case – GNAT

Thread_Local_Storage – GNAT

Type_Invariant

Unchecked_Union

Universal_Aliasing – GNAT
Note that for aspects with an expression, e.g. Size, the expression is treated like a default expression (visibility is analyzed at the point of occurrence of the aspect, but evaluation of the expression occurs at the freeze point of the entity involved).

RM References: 3.02.01 (3) 3.02.02 (2) 3.03.01 (2/2) 3.08 (6) 3.09.03 (1.1/2) 6.01 (2/2) 6.07 (2/2) 9.05.02 (2/2) 7.01 (3) 7.03 (2) 7.03 (3) 9.01 (2/2) 9.01 (3/2) 9.04 (2/2) 9.04 (3/2) 9.05.02 (2/2) 11.01 (2) 12.01 (3) 12.03 (2/2) 12.04 (2/2) 12.05 (2) 12.06 (2.1/2) 12.06 (2.2/2) 12.07 (2) 13.01 (0.1/2) 13.03 (5/1) 13.03.01 (0)

* ‘AI-0128 Inequality is a primitive operation (0000-00-00)’
If an equality operator ("=") is declared for a type, then the implicitly declared inequality operator ("/=") is a primitive operation of the type. This is the only reasonable interpretation, and is the one always implemented by GNAT, but the RM was not entirely clear in making this point.
RM References: 3.02.03 (6) 6.06 (6)

* ‘AI-0003 Qualified expressions as names (2010-07-11)’
In Ada 2012, a qualified expression is considered to be syntactically a name, meaning that constructs such as A'(F(X)).B are now legal. This is useful in disambiguating some cases of overloading.
RM References: 3.03 (11) 3.03 (21) 4.01 (2) 4.04 (7) 4.07 (3) 5.04 (7)

* ‘AI-0120 Constant instance of protected object (0000-00-00)’
This is an RM editorial change only. The section that lists objects that are constant failed to include the current instance of a protected object within a protected function. This has always been treated as a constant in GNAT.
RM References: 3.03 (21)

* ‘AI-0008 General access to constrained objects (0000-00-00)’
The wording in the RM implied that if you have a general access to a constrained object, it could be used to modify the discriminants. This was obviously not intended. Constraint_Error should be raised, and GNAT has always done so in this situation.
Chapter 16: Implementation of Ada 2012 Features

RM References: 3.03 (23) 3.10.02 (26/2) 4.01 (9) 6.04.01 (17) 8.05.01 (5/2)

* ‘AI-0093 Additional rules use immutably limited (0000-00-00)’
This is an editorial change only, to make more widespread use of the Ada 2012 ‘immutably limited’.
RM References: 3.03 (23.4/3)

* ‘AI-0096 Deriving from formal private types (2010-07-20)’
In general it is illegal for a type derived from a formal limited type to be nonlimited. This AI makes an exception to this rule: derivation is legal if it appears in the private part of the generic, and the formal type is not tagged. If the type is tagged, the legality check must be applied to the private part of the package.
RM References: 3.04 (5.1/2) 6.02 (7)

* ‘AI-0181 Soft hyphen is a non-graphic character (2010-07-23)’
From Ada 2005 on, soft hyphen is considered a non-graphic character, which means that it has a special name (SOFT_HYPHEN) in conjunction with the Image and Value attributes for the character types. Strictly speaking this is an inconsistency with Ada 95, but in practice the use of these attributes is so obscure that it will not cause problems.
RM References: 3.05.02 (2/2) A.01 (35/2) A.03.03 (21)

* ‘AI-0182 Additional forms for’ Character’Value ‘(0000-00-00)’
This AI allows Character’Value to accept the string ’?’ where ? is any character including non-graphic control characters. GNAT has always accepted such strings. It also allows strings such as HEX_00000041 to be accepted, but GNAT does not take advantage of this permission and raises Constraint_Error, as is certainly still permitted.
RM References: 3.05 (56/2)

* ‘AI-0214 Defaulted discriminants for limited tagged (2010-10-01)’
Ada 2012 relaxes the restriction that forbids discriminants of tagged types to have default expressions by allowing them when the type is limited. It is often useful to define a default value for a discriminant even though it can’t be changed by assignment.
RM References: 3.07 (9.1/2) 3.07.02 (3)

* ‘AI-0102 Some implicit conversions are illegal (0000-00-00)’
It is illegal to assign an anonymous access constant to an anonymous access variable. The RM did not have a clear rule to prevent this, but GNAT has always generated an error for this usage.
RM References: 3.07 (16) 3.07.01 (9) 6.04.01 (6) 8.06 (27/2)

* ‘AI-0158 Generalizing membership tests (2010-09-16)’
This AI extends the syntax of membership tests to simplify complex conditions that can be expressed as membership in a subset of values of any type. It introduces syntax for a list of expressions that may be used in loop contexts as well.
RM References: 3.08.01 (5) 4.04 (3) 4.05.02 (3) 4.05.02 (5) 4.05.02 (27)

* ‘AI-0173 Testing if tags represent abstract types (2010-07-03)’
The function Ada.Tags.Type_Is_Abstract returns True if invoked with the tag of an abstract type, and False otherwise.
RM References: 3.09 (7.4/2) 3.09 (12.4/2)
* ‘AI-0076 function with controlling result (0000-00-00)’
This is an editorial change only. The RM defines calls with controlling results, but uses
the term ‘function with controlling result’ without an explicit definition.
RM References: 3.09.02 (2/2)
* ‘AI-0126 Dispatching with no declared operation (0000-00-00)’
This AI clarifies dispatching rules, and simply confirms that dispatching executes the
operation of the parent type when there is no explicitly or implicitly declared operation
for the descendant type. This has always been the case in all versions of GNAT.
RM References: 3.09.02 (20/2) 3.09.02 (20.1/2) 3.09.02 (20.2/2)
* ‘AI-0097 Treatment of abstract null extension (2010-07-19)’
The RM as written implied that in some cases it was possible to create an object of
an abstract type, by having an abstract extension inherit a non-abstract constructor
from its parent type. This mistake has been corrected in GNAT and in the RM, and
this construct is now illegal.
RM References: 3.09.03 (4/2)
* ‘AI-0203 Extended return cannot be abstract (0000-00-00)’
A return_subtype_indication cannot denote an abstract subtype. GNAT has never
permitted such usage.
RM References: 3.09.03 (8/3)
* ‘AI-0198 Inheriting abstract operators (0000-00-00)’
This AI resolves a conflict between two rules involving inherited abstract operations
and predefined operators. If a derived numeric type inherits an abstract operator, it
overrides the predefined one. This interpretation was always the one implemented in
GNAT.
RM References: 3.09.03 (4/3)
* ‘AI-0073 Functions returning abstract types (2010-07-10)’
This AI covers a number of issues regarding returning abstract types. In particular
generic functions cannot have abstract result types or access result types designated
an abstract type. There are some other cases which are detailed in the AI. Note that
this binding interpretation has not been retrofitted to operate before Ada 2012 mode,
since it caused a significant number of regressions.
RM References: 3.09.03 (8) 3.09.03 (10) 6.05 (8/2)
* ‘AI-0070 Elaboration of interface types (0000-00-00)’
This is an editorial change only, there are no testable consequences short of checking
for the absence of generated code for an interface declaration.
RM References: 3.09.04 (18/2)
* ‘AI-0208 Characteristics of incomplete views (0000-00-00)’
The wording in the Ada 2005 RM concerning characteristics of incomplete views was
incorrect and implied that some programs intended to be legal were now illegal. GNAT
had never considered such programs illegal, so it has always implemented the intent of
this AI.
Incomplete types are made more useful by allowing them to be completed by private types and private extensions.

An unintentional omission in the RM implied some inconsistent restrictions on the use of anonymous access to subprogram values. These restrictions were not intentional, and have never been enforced by GNAT.

A choice list in a record aggregate can include several components of (distinct) anonymous access types as long as they have matching designated subtypes.

This AI addresses a wording problem in the RM that appears to permit some complex cases of aggregates with nonstatic discriminants. GNAT has always implemented the intended semantics.

This AI confirms that an association of the form `Indx => <>` in an array aggregate must raise `Constraint_Error` if `Indx` is out of range. The RM specified a range check on other associations, but not when the value of the association was defaulted. GNAT has always inserted a constraint check on the index value.

Equality of untagged record composes, so that the predefined equality for a composite type that includes a component of some untagged record type `R` uses the equality operation of `R` (which may be user-defined or predefined). This makes the behavior of untagged records identical to that of tagged types in this respect.

This change is an incompatibility with previous versions of Ada, but it corrects a non-uniformity that was often a source of confusion. Analysis of a large number of industrial programs indicates that in those rare cases where a composite type had an untagged
record component with a user-defined equality, either there was no use of the composite
equality, or else the code expected the same composability as for tagged types, and thus
had a bug that would be fixed by this change.
RM References: 4.05.02 (9.7/2) 4.05.02 (14) 4.05.02 (15) 4.05.02 (24) 8.05.04 (8)

* ‘AI-0088 The value of exponentiation (0000-00-00)’
This AI clarifies the equivalence rule given for the dynamic semantics of exponentiation:
the value of the operation can be obtained by repeated multiplication, but the operation
can be implemented otherwise (for example using the familiar divide-by-two-and-square
algorithm, even if this is less accurate), and does not imply repeated reads of a volatile
base.
RM References: 4.05.06 (11)

* ‘AI-0188 Case expressions (2010-01-09)’
Case expressions are permitted. This allows use of constructs such as:

\[ X := (\text{case } Y \text{ is when } 1 => 2, \text{ when } 2 => 3, \text{ when others } => 31) \]

RM References: 4.05.07 (0) 4.05.08 (0) 4.09 (12) 4.09 (33)

* ‘AI-0104 Null exclusion and uninitialized allocator (2010-07-15)’
The assignment \( \text{Ptr} := \text{new not null Some_Ptr} \); will raise \text{Constraint_Error} because
the default value of the allocated object is ‘null’. This useless construct is illegal in
Ada 2012.
RM References: 4.08 (2)

* ‘AI-0157 Allocation/Deallocation from empty pool (2010-07-11)’
Allocation and Deallocation from an empty storage pool (i.e. allocation or deallocation
of a pointer for which a static storage size clause of zero has been given) is now illegal
and is detected as such. GNAT previously gave a warning but not an error.
RM References: 4.08 (5.3/2) 13.11.02 (4) 13.11.02 (17)

* ‘AI-0179 Statement not required after label (2010-04-10)’
It is not necessary to have a statement following a label, so a label can appear at the end
of a statement sequence without the need for putting a null statement afterwards, but
it is not allowable to have only labels and no real statements in a statement sequence.
RM References: 5.01 (2)

* ‘AI-0139-2 Syntactic sugar for iterators (2010-09-29)’
The new syntax for iterating over arrays and containers is now implemented. Iteration
over containers is for now limited to read-only iterators. Only default iterators are
supported, with the syntax: \text{for Elem of C}.
RM References: 5.05

* ‘AI-0134 Profiles must match for full conformance (0000-00-00)’
For full conformance, the profiles of anonymous-access-to-subprogram parameters must
match. GNAT has always enforced this rule.
RM References: 6.03.01 (18)

* ‘AI-0207 Mode conformance and access constant (0000-00-00)’
This AI confirms that access_to_constant indication must match for mode conformance.
This was implemented in GNAT when the qualifier was originally introduced in Ada
2005.
For full conformance, in the case of access parameters, the null exclusion must match (either both or neither must have not null).

This AI clarifies the rules for named associations in subprogram calls and generic instantiations. The rules have been in place since Ada 83.

Null exclusion checks are not made for out parameters when evaluating the actual parameters. GNAT has never generated these checks.

The return object declared in an `extended_return_statement` may be declared constant. This was always intended, and GNAT has always allowed it.

If a function returns a class-wide type, the object of an extended return statement can be declared with a specific type that is covered by the class-wide type. This has been implemented in GNAT since the introduction of extended returns. Note AI-0103 complements this AI by imposing matching rules for constrained return types.

If the return subtype of a function is an elementary type or a constrained type, the subtype indication in an extended return statement must match statically this return subtype.

The RM had some incorrect wording implying wrong treatment of abnormal completion in an extended return. GNAT has always implemented the intended correct semantics as described by this AI.

The implementation permissions for raising `Constraint_Error` early on a function call when it was clear an exception would be raised were over-permissive and allowed mishandling of discriminants in some cases. GNAT did not take advantage of these incorrect permissions in any case.

In Ada 2012, the declaration of a primitive operation of a type extension or private extension can also override an inherited primitive that is not visible at the point of this declaration.
RM References: 7.03.01 (6) 8.03 (23) 8.03.01 (5/2) 8.03.01 (6/2)

* ‘AI-0062 Null exclusions and deferred constants (0000-00-00)’
  A full constant may have a null exclusion even if its associated deferred constant does not. GNAT has always allowed this.
RM References: 7.04 (6/2) 7.04 (7.1/2)

* ‘AI-0178 Incomplete views are limited (0000-00-00)’
  This AI clarifies the role of incomplete views and plugs an omission in the RM. GNAT always correctly restricted the use of incomplete views and types.
RM References: 7.05 (3/2) 7.05 (6/2)

* ‘AI-0087 Actual for formal nonlimited derived type (2010-07-15)’
  The actual for a formal nonlimited derived type cannot be limited. In particular, a formal derived type that extends a limited interface but which is not explicitly limited cannot be instantiated with a limited type.
RM References: 7.05 (5/2) 12.05.01 (5.1/2)

* ‘AI-0099 Tag determines whether finalization needed (0000-00-00)’
  This AI clarifies that ‘needs finalization’ is part of dynamic semantics, and therefore depends on the run-time characteristics of an object (i.e. its tag) and not on its nominal type. As the AI indicates: “we do not expect this to affect any implementation”.
RM References: 7.06.01 (6) 7.06.01 (7) 7.06.01 (8) 7.06.01 (9/2)

* ‘AI-0064 Redundant finalization rule (0000-00-00)’
  This is an editorial change only. The intended behavior is already checked by an existing ACATS test, which GNAT has always executed correctly.
RM References: 7.06.01 (17.1/1)

* ‘AI-0026 Missing rules for Unchecked_Union (2010-07-07)’
  Record representation clauses concerning Unchecked_Union types cannot mention the discriminant of the type. The type of a component declared in the variant part of an Unchecked_Union cannot be controlled, have controlled components, nor have protected or task parts. If an Unchecked_Union type is declared within the body of a generic unit or its descendants, then the type of a component declared in the variant part cannot be a formal private type or a formal private extension declared within the same generic unit.
RM References: 7.06 (9.4/2) B.03.03 (9/2) B.03.03 (10/2)

* ‘AI-0205 Extended return declares visible name (0000-00-00)’
  This AI corrects a simple omission in the RM. Return objects have always been visible within an extended return statement.
RM References: 8.03 (17)

* ‘AI-0042 Overriding versus implemented-by (0000-00-00)’
  This AI fixes a wording gap in the RM. An operation of a synchronized interface can be implemented by a protected or task entry, but the abstract operation is not being overridden in the usual sense, and it must be stated separately that this implementation is legal. This has always been the case in GNAT.
RM References: 9.01 (9.2/2) 9.04 (11.1/2)
* ‘AI-0030 Requeue on synchronized interfaces (2010-07-19)’
Requeue is permitted to a protected, synchronized or task interface primitive providing it is known that the overriding operation is an entry. Otherwise the requeue statement has the same effect as a procedure call. Use of pragma Implemented provides a way to impose a static requirement on the overriding operation by adhering to one of the implementation kinds: entry, protected procedure or any of the above.
RM References: 9.05 (9) 9.05.04 (2) 9.05.04 (3) 9.05.04 (5) 9.05.04 (6) 9.05.04 (7) 9.05.04 (12)

* ‘AI-0201 Independence of atomic object components (2010-07-22)’
If an Atomic object has a pragma Pack or a Component_Size attribute, then individual components may not be addressable by independent tasks. However, if the representation clause has no effect (is confirming), then independence is not compromised. Furthermore, in GNAT, specification of other appropriately addressable component sizes (e.g. 16 for 8-bit characters) also preserves independence. GNAT now gives very clear warnings both for the declaration of such a type, and for any assignment to its components.
RM References: 9.10 (1/3) C.06 (22/2) C.06 (23/2)

* ‘AI-0009 Pragma Independent[.Components] (2010-07-23)’
This AI introduces the new pragmas Independent and Independent_Components, which control guaranteeing independence of access to objects and components. The AI also requires independence not unaffected by confirming rep clauses.
RM References: 9.10 (1) 13.01 (15/1) 13.02 (9) 13.03 (13) C.06 (2) C.06 (4) C.06 (6) C.06 (9) C.06 (13) C.06 (14)

* ‘AI-0072 Task signalling using ‘Terminated (0000-00-00)’
This AI clarifies that task signalling for reading ‘Terminated only occurs if the result is True. GNAT semantics has always been consistent with this notion of task signalling.
RM References: 9.10 (6.1/1)

* ‘AI-0108 Limited incomplete view and discriminants (0000-00-00)’
This AI confirms that an incomplete type from a limited view does not have discriminants. This has always been the case in GNAT.
RM References: 10.01.01 (12.3/2)

* ‘AI-0129 Limited views and incomplete types (0000-00-00)’
This AI clarifies the description of limited views: a limited view of a package includes only one view of a type that has an incomplete declaration and a full declaration (there is no possible ambiguity in a client package). This AI also fixes an omission: a nested package in the private part has no limited view. GNAT always implemented this correctly.
RM References: 10.01.01 (12.2/2) 10.01.01 (12.3/2)

* ‘AI-0077 Limited withs and scope of declarations (0000-00-00)’
This AI clarifies that a declaration does not include a context clause, and confirms that it is illegal to have a context in which both a limited and a nonlimited view of a package are accessible. Such double visibility was always rejected by GNAT.
RM References: 10.01.02 (12/2) 10.01.02 (21/2) 10.01.02 (22/2)
* ‘AI-0122 Private with and children of generics (0000-00-00)’
This AI clarifies the visibility of private children of generic units within instantiations of a parent. GNAT has always handled this correctly.
RM References: 10.01.02 (12/2)

* ‘AI-0040 Limited with clauses on descendant (0000-00-00)’
This AI confirms that a limited with clause in a child unit cannot name an ancestor of the unit. This has always been checked in GNAT.
RM References: 10.01.02 (20/2)

* ‘AI-0132 Placement of library unit pragmas (0000-00-00)’
This AI fills a gap in the description of library unit pragmas. The pragma clearly must apply to a library unit, even if it does not carry the name of the enclosing unit. GNAT has always enforced the required check.
RM References: 10.01.05 (7)

* ‘AI-0034 Categorization of limited views (0000-00-00)’
The RM makes certain limited with clauses illegal because of categorization considerations, when the corresponding normal with would be legal. This is not intended, and GNAT has always implemented the recommended behavior.
RM References: 10.02.01 (11/1) 10.02.01 (17/2)

* ‘AI-0035 Inconsistencies with Pure units (0000-00-00)’
This AI remedies some inconsistencies in the legality rules for Pure units. Derived access types are legal in a pure unit (on the assumption that the rule for a zero storage pool size has been enforced on the ancestor type). The rules are enforced in generic instances and in subunits. GNAT has always implemented the recommended behavior.
RM References: 10.02.01 (15.1/2) 10.02.01 (15.4/2) 10.02.01 (15.5/2) 10.02.01 (17/2)

* ‘AI-0219 Pure permissions and limited parameters (2010-05-25)’
This AI refines the rules for the cases with limited parameters which do not allow the implementations to omit ‘redundant’. GNAT now properly conforms to the requirements of this binding interpretation.
RM References: 10.02.01 (18/2)

* ‘AI-0043 Rules about raising exceptions (0000-00-00)’
This AI covers various omissions in the RM regarding the raising of exceptions. GNAT has always implemented the intended semantics.
RM References: 11.04.01 (10.1/2) 11 (2)

* ‘AI-0200 Mismatches in formal package declarations (0000-00-00)’
This AI plugs a gap in the RM which appeared to allow some obviously intended illegal instantiations. GNAT has never allowed these instantiations.
RM References: 12.07 (16)

* ‘AI-0112 Detection of duplicate pragmas (2010-07-24)’
This AI concerns giving names to various representation aspects, but the practical effect is simply to make the use of duplicate Atomic[.Components], Volatile[.Components], and Independent[.Components] pragmas illegal, and GNAT now performs this required check.
The RM appeared to allow representation pragmas on generic formal parameters, but this was not intended, and GNAT has never permitted this usage.

It is now illegal to give an inappropriate component size or a pragma Pack that attempts to change the component size in the case of atomic or aliased components. Previously GNAT ignored such an attempt with a warning.

The RM permitted the use of dynamic expressions (such as ptr.all) for stream attributes, but these were never useful and are now illegal. GNAT has always regarded such expressions as illegal.

This AI requires that the alignment of a class-wide object be no greater than the alignment of any type in the class. GNAT has always followed this recommendation.

Type invariants may be specified for private types using the aspect notation. Aspect Type_Invariant may be specified for any private type, Type_Invariant'Class can only be specified for tagged types, and is inherited by any descendent of the tagged types. The invariant is a boolean expression that is tested for being true in the following situations: conversions to the private type, object declarations for the private type that are default initialized, and ['in'] 'out' parameters and returned result on return from any primitive operation for the type that is visible to a client. GNAT defines the synonyms Invariant for Type_Invariant and Invariant'Class for Type_Invariant'Class.

In Ada 2012, compilers are required to support unchecked conversion where the target alignment is a multiple of the source alignment. GNAT always supported this case (and indeed all cases of differing alignments, doing copies where required if the alignment was reduced).

The handling of invalid values is now designated to be implementation defined. This is a documentation change only, requiring Annex M in the GNAT Reference Manual.
to document this handling. In GNAT, checks for invalid values are made only when necessary to avoid erroneous behavior. Operations like assignments which cannot cause erroneous behavior ignore the possibility of invalid values and do not do a check. The date given above applies only to the documentation change, this behavior has always been implemented by GNAT.
RM References: 13.09.01 (10)
* ‘AI-0193 Alignment of allocators (2010-09-16)’
This AI introduces a new attribute \texttt{Max\_Alignment\_For\_Allocation}, analogous to \texttt{Max\_Size\_In\_Storage\_Elements}, but for alignment instead of size.
RM References: 13.11 (16) 13.11 (21) 13.11.01 (0) 13.11.01 (1) 13.11.01 (2) 13.11.01 (3)
* ‘AI-0177 Parameterized expressions (2010-07-10)’
The new Ada 2012 notion of parameterized expressions is implemented. The form is:

\begin{verbatim}
function-specification is (expression)
\end{verbatim}

This is exactly equivalent to the corresponding function body that returns the expression, but it can appear in a package spec. Note that the expression must be parenthesized.
RM References: 13.11.01 (3/2)
* ‘AI-0033 Attach/Interrupt \_ Handler in generic (2010-07-24)’
Neither of these two pragmas may appear within a generic template, because the generic might be instantiated at other than the library level.
RM References: 13.11.02 (16) C.03.01 (7/2) C.03.01 (8/2)
* ‘AI-0161 Restriction No\_Default\_Stream\_Attributes (2010-09-11)’
A new restriction \texttt{No\_Default\_Stream\_Attributes} prevents the use of any of the default stream attributes for elementary types. If this restriction is in force, then it is necessary to provide explicit subprograms for any stream attributes used.
RM References: 13.12.01 (4/2) 13.13.02 (40/2) 13.13.02 (52/2)
* ‘AI-0194 Value of Stream\_Size attribute (0000-00-00)’
The \texttt{Stream\_Size} attribute returns the default number of bits in the stream representation of the given type. This value is not affected by the presence of stream subprogram attributes for the type. GNAT has always implemented this interpretation.
RM References: 13.13.02 (1.2/2)
* ‘AI-0109 Redundant check in S\_Class\_Input (0000-00-00)’
This AI is an editorial change only. It removes the need for a tag check that can never fail.
RM References: 13.13.02 (34/2)
* ‘AI-0007 Stream read and private scalar types (0000-00-00)’
The RM as written appeared to limit the possibilities of declaring read attribute procedures for private scalar types. This limitation was not intended, and has never been enforced by GNAT.
RM References: 13.13.02 (50/2) 13.13.02 (51/2)
* ‘AI-0065 Remote access types and external streaming (0000-00-00)’
This AI clarifies the fact that all remote access types support external streaming. This fixes an obvious oversight in the definition of the language, and GNAT always implemented the intended correct rules.
RM References: 13.13.02 (52/2)

* ‘AI-0019 Freezing of primitives for tagged types (0000-00-00)’
The RM suggests that primitive subprograms of a specific tagged type are frozen when the tagged type is frozen. This would be an incompatible change and is not intended. GNAT has never attempted this kind of freezing and its behavior is consistent with the recommendation of this AI.

* ‘AI-0017 Freezing and incomplete types (0000-00-00)’
So-called ‘Taft-amendment types’ (i.e., types that are completed in package bodies) are not frozen by the occurrence of bodies in the enclosing declarative part. GNAT always implemented this properly.
RM References: 13.14 (3/1)

* ‘AI-0060 Extended definition of remote access types (0000-00-00)’
This AI extends the definition of remote access types to include access to limited, synchronized, protected or task class-wide interface types. GNAT already implemented this extension.
RM References: A (4) E.02.02 (9/1) E.02.02 (9.2/1) E.02.02 (14/2) E.02.02 (18)

* ‘AI-0114 Classification of letters (0000-00-00)’
The code points 170 (FEMININE ORDINAL INDICATOR), 181 (MICRO SIGN), and 186 (MASCULINE ORDINAL INDICATOR) are technically considered lower case letters by Unicode. However, they are not allowed in identifiers, and they return False to Ada.Characters.Handling.Is_Letter/Is_Lower. This behavior is consistent with that defined in Ada 95.
RM References: A.03.02 (59) A.04.06 (7)

* ‘AI-0185 Ada.Wide_[Wide_]Characters.Handling (2010-07-06)’
Two new packages Ada.Wide_[Wide_]Characters.Handling provide classification functions for Wide_Character and Wide_Wide_Character, as well as providing case folding routines for Wide_[Wide_]Character and Wide_[Wide_]String.
RM References: A.03.05 (0) A.03.06 (0)

* ‘AI-0031 Add From parameter to Find_Token (2010-07-25)’
A new version of Find_Token is added to all relevant string packages, with an extra parameter From. Instead of starting at the first character of the string, the search for a matching Token starts at the character indexed by the value of From. These procedures are available in all versions of Ada but if used in versions earlier than Ada 2012 they will generate a warning that an Ada 2012 subprogram is being used.
RM References: A.04.03 (16) A.04.03 (67) A.04.03 (68/1) A.04.04 (51) A.04.05 (46)

* ‘AI-0056 Index on null string returns zero (0000-00-00)’
The wording in the Ada 2005 RM implied an incompatible handling of the Index functions, resulting in raising an exception instead of returning zero in some situations.
This was not intended and has been corrected. GNAT always returned zero, and is thus consistent with this AI.

RM References: A.04.03 (56.2/2) A.04.03 (58.5/2)

* ‘AI-0137 String encoding package (2010-03-25)’
The packages Ada.Strings.UTF_Encoding, together with its child packages, Conversions, Strings, Wide.Strings, and Wide_Wide.Strings have been implemented. These packages (whose documentation can be found in the spec files a-stuten.ads, a-suenco.ads, a-suenst.ads, a-suewst.ads, a-suezst.ads) allow encoding and decoding of String, Wide_String, and Wide_Wide_String values using UTF coding schemes (including UTF-8, UTF-16LE, UTF-16BE, and UTF-16), as well as conversions between the different UTF encodings. With the exception of Wide_Wide.Strings, these packages are available in Ada 95 and Ada 2005 mode as well as Ada 2012 mode. The Wide_Wide.Strings package is available in Ada 2005 mode as well as Ada 2012 mode (but not in Ada 95 mode since it uses Wide_Wide.Character).

RM References: A.04.11

* ‘AI-0038 Minor errors in Text_IO (0000-00-00)’
These are minor errors in the description on three points. The intent on all these points has always been clear, and GNAT has always implemented the correct intended semantics.

RM References: A.10.05 (37) A.10.07 (8/1) A.10.07 (10) A.10.07 (12) A.10.08 (10) A.10.08 (24)

* ‘AI-0044 Restrictions on container instantiations (0000-00-00)’
This AI places restrictions on allowed instantiations of generic containers. These restrictions are not checked by the compiler, so there is nothing to change in the implementation. This affects only the RM documentation.

RM References: A.18 (4/2) A.18.02 (231/2) A.18.03 (145/2) A.18.06 (56/2) A.18.08 (66/2) A.18.09 (79/2) A.18.26 (5/2) A.18.26 (9/2)

* ‘AI-0127 Adding Locale Capabilities (2010-09-29)’
This package provides an interface for identifying the current locale.


* ‘AI-0002 Export C with unconstrained arrays (0000-00-00)’
The compiler is not required to support exporting an Ada subprogram with convention C if there are parameters or a return type of an unconstrained array type (such as String). GNAT allows such declarations but generates warnings. It is possible, but complicated, to write the corresponding C code and certainly such code would be specific to GNAT and non-portable.

RM References: B.01 (17) B.03 (62) B.03 (71.1/2)

* ‘AI-0216 No_Task_Hierarchy forbids local tasks (0000-00-00)’
It is clearly the intention that No_Task_Hierarchy is intended to forbid tasks declared locally within subprograms, or functions returning task objects, and that is the implementation that GNAT has always provided. However the language in the RM was not sufficiently clear on this point. Thus this is a documentation change in the RM only.
RM References: D.07 (3/3)
* ‘AI-0211 No_Relative_Delays forbids Set_Handler use (2010-07-09)’
The restriction No_Relative_Delays forbids any calls to the subprogram Ada.Real_Time.Timing_Events.Set_Handler.
RM References: D.07 (5) D.07 (10/2) D.07 (10.4/2) D.07 (10.7/2)
* ‘AI-0190 pragma Default_Storage_Pool (2010-09-15)’
This AI introduces a new pragma Default_Storage_Pool, which can be used to control storage pools globally. In particular, you can force every access type that is used for allocation (‘new’) to have an explicit storage pool, or you can declare a pool globally to be used for all access types that lack an explicit one.
RM References: D.07 (8)
* ‘AI-0189 No_Allocators_After_Elaboration (2010-01-23)’
This AI introduces a new restriction No_Allocators_After_Elaboration, which says that no dynamic allocation will occur once elaboration is completed. In general this requires a run-time check, which is not required, and which GNAT does not attempt. But the static cases of allocators in a task body or in the body of the main program are detected and flagged at compile or bind time.
RM References: D.07 (19.1/2) H.04 (23.3/2)
* ‘AI-0171 Pragma CPU and Ravenscar Profile (2010-09-24)’
A new package System.Multiprocessors is added, together with the definition of pragma CPU for controlling task affinity. A new no dependence restriction, on System.Multiprocessors.Dispatching_Domains, is added to the Ravenscar profile.
RM References: D.13.01 (4/2) D.16
* ‘AI-0210 Correct Timing_Events metric (0000-00-00)’
This is a documentation only issue regarding wording of metric requirements, that does not affect the implementation of the compiler.
RM References: D.15 (24/2)
* ‘AI-0206 Remote types packages and preelaborate (2010-07-24)’
Remote types packages are now allowed to depend on preelaborated packages. This was formerly considered illegal.
RM References: E.02.02 (6)
* ‘AI-0152 Restriction No_Anonymous_Allocators (2010-09-08)’
Restriction No_Anonymous_Allocators prevents the use of allocators where the type of the returned value is an anonymous access type.
RM References: H.04 (8/1)
17 Security Hardening Features

This chapter describes Ada extensions aimed at security hardening that are provided by GNAT.

The features in this chapter are currently experimental and subject to change.

17.1 Register Scrubbing

GNAT can generate code to zero-out hardware registers before returning from a subprogram. It can be enabled with the `-fzero-call-used-regs=choice` command-line option, to affect all subprograms in a compilation, and with a `Machine_Attribute` pragma, to affect only specific subprograms.

```ada
procedure Foo;
pragma Machine_Attribute (Foo, "zero_call_used_regs", "used");
-- Before returning, Foo scrubs only call-clobbered registers
-- that it uses itself.
```

```ada
function Bar return Integer;
pragma Machine_Attribute (Bar, "zero_call_used_regs", "all");
-- Before returning, Bar scrubs all call-clobbered registers.
```

For usage and more details on the command-line option, on the `zero_call_used_regs` attribute, and on their use with other programming languages, see Using the GNU Compiler Collection (GCC).

17.2 Stack Scrubbing

GNAT can generate code to zero-out stack frames used by subprograms. It can be activated with the `Machine_Attribute` pragma, on specific subprograms and variables, or their types. (This attribute always applies to a type, even when it is associated with a subprogram or a variable.)

```ada
function Foo returns Integer;
pragma Machine_Attribute (Foo, "strub");
-- Foo and its callers are modified so as to scrub the stack
-- space used by Foo after it returns. Shorthand for:
-- pragma Machine_Attribute (Foo, "strub", "at-calls");
```

```ada
procedure Bar;
pragma Machine_Attribute (Bar, "strub", "internal");
-- Bar is turned into a wrapper for its original body,
-- and they scrub the stack used by the original body.
```

```ada
Var : Integer;
pragma Machine_Attribute (Var, "strub");
-- Reading from Var in a subprogram enables stack scrubbing
-- of the stack space used by the subprogram. Furthermore, if
-- Var is declared within a subprogram, this also enables
```
-- scrubbing of the stack space used by that subprogram.

Given these declarations, Foo has its type and body modified as follows:

```ada
function Foo (<WaterMark> : in out System.Address) returns Integer is
  -- ...
begin
  <__strub_update> (<WaterMark>); -- Updates the stack WaterMark.
  -- ...
end;
```

whereas its callers are modified from:

```ada
X := Foo;
```


to:

```ada
declare
  <WaterMark> : System.Address;
begin
  <__strub_enter> (<WaterMark>); -- Initialize <WaterMark>.
  X := Foo (<WaterMark>);
  <__strub_leave> (<WaterMark>); -- Scrubs stack up to <WaterMark>.
end;
```

As for Bar, because it is strubbed in internal mode, its callers are not modified. Its definition is modified roughly as follows:

```ada
procedure Bar is
  <WaterMark> : System.Address;
  procedure Strubbed_Bar (<WaterMark> : in out System.Address) is
    begin
      <__strub_update> (<WaterMark>); -- Updates the stack WaterMark.
      -- original Bar body.
    end Strubbed_Bar;
  begin
    <__strub_enter> (<WaterMark>); -- Initialize <WaterMark>.
    Strubbed_Bar (<WaterMark>);
    <__strub_leave> (<WaterMark>); -- Scrubs stack up to <WaterMark>.
end Bar;
```

There are also `-fstrub='choice'` command-line options to control default settings. For usage and more details on the command-line options, on the `strub` attribute, and their use with other programming languages, see Using the GNU Compiler Collection (GCC).

Note that Ada secondary stacks are not scrubbed. The restriction `No_Secondary_Stack` avoids their use, and thus their accidental preservation of data that should be scrubbed.

Attributes `Access` and `Unconstrained_Access` of variables and constants with `strub` enabled require types with `strub` enabled; there is no way to express an access-to-strub type otherwise. `Unchecked_Access` bypasses this constraint, but the resulting access type designates a non-strub type.

```ada
VI : aliased Integer;
pragma Machine_Attribute (VI, "strub");
```
Every access-to-subprogram type, renaming, and overriding and overridden dispatching operations that may refer to a subprogram with an attribute-modified interface must be annotated with the same interface-modifying attribute. Access-to-subprogram types can be explicitly converted to different strub modes, as long as they are interface-compatible (i.e., adding or removing at-calls is not allowed). For example, a strub-disabled subprogram can be turned callable through such an explicit conversion:

```
XsVI : access Integer := VI'Access; -- Error.
UXsVI : access Integer := VI'Unchecked_Access; -- OK,
-- UXsVI does *not* enable strub in subprograms that
dereference it to obtain the UXsVI.all value.

type Strub_Int is new Integer;
pragma Machine_Attribute (Strub_Int, "strub");
VSI : aliased Strub_Int;
XsVSI : access Strub_Int := VSI'Access; -- OK,
-- VSI and XsVSI.all both enable strub in subprograms that
-- read their values.
```

17.3 Hardened Conditionals

GNAT can harden conditionals to protect against control-flow attacks. This is accomplished by two complementary transformations, each activated by a separate command-line option.

The option `-fharden-compares` enables hardening of compares that compute results stored in variables, adding verification that the reversed compare yields the opposite result, turning:

```
B := X = Y;
```

into:

```
B := X = Y;
```
declare
  NotB : Boolean := X /= Y; -- Computed independently of B.
begin
  if B = NotB then
    //__builtin_trap;
  end if;
end;

The option -fharden-conditional-branches enables hardening of compares that guard conditional branches, adding verification of the reversed compare to both execution paths, turning:

```plaintext
if X = Y then
  X := Z + 1;
else
  Y := Z - 1;
end if;
```

into:

```plaintext
if X = Y then
  if X /= Y then -- Computed independently of X = Y.
    //__builtin_trap;
  end if;
  X := Z + 1;
else
  if X /= Y then -- Computed independently of X = Y.
    null;
  else
    //__builtin_trap;
  end if;
  Y := Z - 1;
end if;
```

These transformations are introduced late in the compilation pipeline, long after boolean expressions are decomposed into separate compares, each one turned into either a conditional branch or a compare whose result is stored in a boolean variable or temporary. Compiler optimizations, if enabled, may also turn conditional branches into stored compares, and vice-versa, or into operations with implied conditionals (e.g. MIN and MAX). Conditionals may also be optimized out entirely, if their value can be determined at compile time, and occasionally multiple compares can be combined into one.

It is thus difficult to predict which of these two options will affect a specific compare operation expressed in source code. Using both options ensures that every compare that is neither optimized out nor optimized into implied conditionals will be hardened.

The addition of reversed compares can be observed by enabling the dump files of the corresponding passes, through command-line options -fdump-tree-hardcmp and -fdump-tree-hardcbr, respectively.

They are separate options, however, because of the significantly different performance impact of the hardening transformations.
For usage and more details on the command-line options, see *Using the GNU Compiler Collection (GCC)*. These options can be used with other programming languages supported by GCC.

### 17.4 Hardened Booleans

Ada has built-in support for introducing boolean types with alternative representations, using representation clauses:

```ada
type HBool is new Boolean;
for HBool use (16#5a#, 16#a5#);
for HBool'Size use 8;
```

When validity checking is enabled, the compiler will check that variables of such types hold values corresponding to the selected representations.

There are multiple strategies for where to introduce validity checking (see `-gnatV` options). Their goal is to guard against various kinds of programming errors, and GNAT strives to omit checks when program logic rules out an invalid value, and optimizers may further remove checks found to be redundant.

For additional hardening, the `hardbool Machine_Attribute` pragma can be used to annotate boolean types with representation clauses, so that expressions of such types used as conditions are checked even when compiling with `-gnatVT`:

```ada
pragma Machine_Attribute (HBool, "hardbool");
```

```ada
function To_Boolean (X : HBool) returns Boolean is (Boolean (X));
```

is compiled roughly like:

```ada
function To_Boolean (X : HBool) returns Boolean is
begin
  if X not in True | False then
    raise Constraint_Error;
  elsif X in True then
    return True;
  else
    return False;
  end if;
end To_Boolean;
```

Note that `-gnatVn` will disable even `hardbool` testing.

Analogous behavior is available as a GCC extension to the C and Objective C programming languages, through the `hardbool` attribute, with the difference that, instead of raising a `Constraint_Error` exception, when a hardened boolean variable is found to hold a value that stands for neither True nor False, the program traps. For usage and more details on that attribute, see *Using the GNU Compiler Collection (GCC)*.

### 17.5 Control Flow Redundancy

GNAT can guard against unexpected execution flows, such as branching into the middle of subprograms, as in Return Oriented Programming exploits.
In units compiled with \texttt{-fharden-control-flow-redundancy}, subprograms are instrumented so that, every time they are called, basic blocks take note as control flows through them, and, before returning, subprograms verify that the taken notes are consistent with the control-flow graph.

Functions with too many basic blocks, or with multiple return points, call a run-time function to perform the verification. Other functions perform the verification inline before returning.

Optimizing the inlined verification can be quite time consuming, so the default upper limit for the inline mode is set at 16 blocks. Command-line option \texttt{--param hardcfr-max-inline-blocks=} can override it.

Even though typically sparse control-flow graphs exhibit run-time verification time nearly proportional to the block count of a subprogram, it may become very significant for generated subprograms with thousands of blocks. Command-line option \texttt{--param hardcfr-max-blocks=} can set an upper limit for instrumentation.

For each block that is marked as visited, the mechanism checks that at least one of its predecessors, and at least one of its successors, are also marked as visited.

Verification is performed just before a subprogram returns. The following fragment:

\begin{verbatim}
if X then
  Y := F (Z);
  return;
end if;
\end{verbatim}

gets turned into:

\begin{verbatim}
type Visited_Bitmap is array (1..N) of Boolean with Pack;
Visited : aliased Visited_Bitmap := (others => False);
-- Bitmap of visited blocks. N is the basic block count.
[...
-- Basic block #I
Visited(I) := True;
if X then
  -- Basic block #J
  Visited(J) := True;
  Y := F (Z);
  CFR.Check (N, Visited'Access, CFG'Access);
  -- CFR is a hypothetical package whose Check procedure calls
  -- libgcc's \_hardcfr\_check, that traps if the Visited bitmap
  -- does not hold a valid path in CFG, the run-time
  -- representation of the control flow graph in the enclosing
  -- subprogram.
  return;
end if;
-- Basic block #K
Visited(K) := True;
\end{verbatim}

Verification would also be performed before tail calls, if any front-ends marked them as mandatory or desirable, but none do. Regular calls are optimized into tail calls too late for this transformation to act on it.
In order to avoid adding verification after potential tail calls, which would prevent tail-call optimization, we recognize returning calls, i.e., calls whose result, if any, is returned by the calling subprogram to its caller immediately after the call returns. Verification is performed before such calls, whether or not they are ultimately optimized to tail calls. This behavior is enabled by default whenever sibcall optimization is enabled (see `-foptimize-sibling-calls`); it may be disabled with `-fno-hardcfr-check-returning-calls`, or enabled with `-fhardcfr-check-returning-calls`, regardless of the optimization, but the lack of other optimizations may prevent calls from being recognized as returning calls:

```
-- CFR.Check here, with -fhardcfr-check-returning-calls.
P (X);
-- CFR.Check here, with -fno-hardcfr-check-returning-calls.
return;

or:

-- CFR.Check here, with -fhardcfr-check-returning-calls.
R := F (X);
-- CFR.Check here, with -fno-hardcfr-check-returning-calls.
return R;
```

Any subprogram from which an exception may escape, i.e., that may raise or propagate an exception that isn’t handled internally, is conceptually enclosed by a cleanup handler that performs verification, unless this is disabled with `-fno-hardcfr-check-exceptions`. With this feature enabled, a subprogram body containing:

```
-- ...
Y := F (X); -- May raise exceptions.
-- ...
raise E; -- Not handled internally.
-- ...
```

gets modified as follows:

```
begin
-- ...
Y := F (X); -- May raise exceptions.
-- ...
raise E; -- Not handled internally.
-- ...
exception
when others =>
  CFR.Check (N, Visited’Access, CFG’Access);
  raise;
end;
```

Verification may also be performed before No_Return calls, whether only nothrow ones, with `-fhardcfr-check-noreturn-calls=nothrow`, or all of them, with `-fhardcfr-check-noreturn-calls=always`. The default is `-fhardcfr-check-noreturn-calls=never` for this feature, that disables checking before No_Return calls.

When a No_Return call returns control to its caller through an exception, verification may have already been performed before the call, if `-fhardcfr-check-noreturn-calls=always` is in effect. The compiler arranges for already-checked No_Return calls without a preexisting
handler to bypass the implicitly-added cleanup handler and thus the redundant check, but a local exception or cleanup handler, if present, will modify the set of visited blocks, and checking will take place again when the caller reaches the next verification point, whether it is a return or reraise statement after the exception is otherwise handled, or even another No_Return call.

The instrumentation for hardening with control flow redundancy can be observed in dump files generated by the command-line option -fdump-tree-hardcfr.

For more details on the control flow redundancy command-line options, see Using the GNU Compiler Collection (GCC). These options can be used with other programming languages supported by GCC.
Chapter 18: Obsolescent Features

18 Obsolescent Features

This chapter describes features that are provided by GNAT, but are considered obsolescent since there are preferred ways of achieving the same effect. These features are provided solely for historical compatibility purposes.

18.1 pragma No_Run_Time

The pragma `No_Run_Time` is used to achieve an affect similar to the use of the “Zero Foot Print” configurable run time, but without requiring a specially configured run time. The result of using this pragma, which must be used for all units in a partition, is to restrict the use of any language features requiring run-time support code. The preferred usage is to use an appropriately configured run-time that includes just those features that are to be made accessible.

18.2 pragma Ravenscar

The pragma `Ravenscar` has exactly the same effect as pragma `Profile (Ravenscar)`. The latter usage is preferred since it is part of the new Ada 2005 standard.

18.3 pragma Restricted_Run_Time

The pragma `Restricted_Run_Time` has exactly the same effect as pragma `Profile (Restricted)`. The latter usage is preferred since the Ada 2005 pragma `Profile` is intended for this kind of implementation dependent addition.

18.4 pragma Task_Info

The functionality provided by pragma `Task_Info` is now part of the Ada language. The CPU aspect and the package `System.Multiprocessors` offer a less system-dependent way to specify task affinity or to query the number of processors.

Syntax

```
pragma Task_Info (EXPRESSION);
```

This pragma appears within a task definition (like pragma `Priority`) and applies to the task in which it appears. The argument must be of type `System.Task_Info.Task_Info_Type`. The `Task_Info` pragma provides system dependent control over aspects of tasking implementation, for example, the ability to map tasks to specific processors. For details on the facilities available for the version of GNAT that you are using, see the documentation in the spec of package `System.Task_Info` in the runtime library.

18.5 package System.Task_Info (s-tasinf.ads)

This package provides target dependent functionality that is used to support the `Task_Info` pragma. The predefined Ada package `System.Multiprocessors` and the CPU aspect now provide a standard replacement for GNAT’s `Task_Info` functionality.
19 Compatibility and Porting Guide

This chapter presents some guidelines for developing portable Ada code, describes the compatibility issues that may arise between GNAT and other Ada compilation systems (including those for Ada 83), and shows how GNAT can expedite porting applications developed in other Ada environments.

19.1 Writing Portable Fixed-Point Declarations

The Ada Reference Manual gives an implementation freedom to choose bounds that are narrower by Small from the given bounds. For example, if we write

\[ \text{type } F1 \text{ is delta 1.0 range } -128.0 \ldots +128.0; \]

then the implementation is allowed to choose -128.0 .. +127.0 if it likes, but is not required to do so.

This leads to possible portability problems, so let’s have a closer look at this, and figure out how to avoid these problems.

First, why does this freedom exist, and why would an implementation take advantage of it? To answer this, take a closer look at the type declaration for \( F1 \) above. If the compiler uses the given bounds, it would need 9 bits to hold the largest positive value (and typically that means 16 bits on all machines). But if the implementation chooses the +127.0 bound then it can fit values of the type in 8 bits.

Why not make the user write +127.0 if that’s what is wanted? The rationale is that if you are thinking of fixed point as a kind of ‘poor man’s floating-point’, then you don’t want to be thinking about the scaled integers that are used in its representation. Let’s take another example:

\[ \text{type } F2 \text{ is delta } 2.0^{(-15)} \text{ range } -1.0 \ldots +1.0; \]

Looking at this declaration, it seems casually as though it should fit in 16 bits, but again that extra positive value +1.0 has the scaled integer equivalent of \( 2^{15} \) which is one too big for signed 16 bits. The implementation can treat this as:

\[ \text{type } F2 \text{ is delta } 2.0^{(-15)} \text{ range } -1.0 \ldots +1.0-(2.0^{(-15)}); \]

and the Ada language design team felt that this was too annoying to require. We don’t need to debate this decision at this point, since it is well established (the rule about narrowing the ranges dates to Ada 83).

But the important point is that an implementation is not required to do this narrowing, so we have a potential portability problem. We could imagine three types of implementation:

a. those that narrow the range automatically if they can figure out that the narrower range will allow storage in a smaller machine unit,

b. those that will narrow only if forced to by a ‘Size clause, and

c. those that will never narrow.

Now if we are language theoreticians, we can imagine a fourth approach: to narrow all the time, e.g. to treat

\[ \text{type } F3 \text{ is delta 1.0 range } -10.0 \ldots +23.0; \]

as though it had been written:
type F3 is delta 1.0 range -9.0 .. +22.0;

But although technically allowed, such a behavior would be hostile and silly, and no real
compiler would do this. All real compilers will fall into one of the categories (a), (b) or (c)
above.

So, how do you get the compiler to do what you want? The answer is give the actual bounds
you want, and then use a 'Small clause and a 'Size clause to absolutely pin down what
the compiler does. E.g., for F2 above, we will write:

\[
\begin{align*}
\text{My}_{\text{Small}} : \text{constant} & := 2.0^{(-15)}; \\
\text{My}_{\text{First}} : \text{constant} & := -1.0; \\
\text{My}_{\text{Last}} : \text{constant} & := +1.0 - \text{My}_{\text{Small}};
\end{align*}
\]

type F2 is delta \text{My}_{\text{Small}} range \text{My}_{\text{First}} .. \text{My}_{\text{Last}};

and then add

\[
\begin{align*}
\text{for F2'Small use} & \text{my}_{\text{Small}}; \\
\text{for F2'Size use} & 16;
\end{align*}
\]

In practice all compilers will do the same thing here and will give you what you want, so
the above declarations are fully portable. If you really want to play language lawyer and
guard against ludicrous behavior by the compiler you could add

\[
\begin{align*}
\text{Test1} : \text{constant} & := 1 / \text{Boolean'Pos (F2'First = My}_{\text{First}}; \\
\text{Test2} : \text{constant} & := 1 / \text{Boolean'Pos (F2'Last = My}_{\text{Last}};
\end{align*}
\]

One or other or both are allowed to be illegal if the compiler is behaving in a silly manner,
but at least the silly compiler will not get away with silently messing with your (very clear)
intentions.

If you follow this scheme you will be guaranteed that your fixed-point types will be portable.

19.2 Compatibility with Ada 83

Ada 95 and the subsequent revisions Ada 2005 and Ada 2012 are highly upwards compatible
with Ada 83. In particular, the design intention was that the difficulties associated with
moving from Ada 83 to later versions of the standard should be no greater than those that
occur when moving from one Ada 83 system to another.

However, there are a number of points at which there are minor incompatibilities. The Ada
95 Annotated Reference Manual contains full details of these issues as they relate to Ada
95, and should be consulted for a complete treatment. In practice the following subsections
treat the most likely issues to be encountered.

19.2.1 Legal Ada 83 programs that are illegal in Ada 95

Some legal Ada 83 programs are illegal (i.e., they will fail to compile) in Ada 95 and later
versions of the standard:

* ‘Character literals’

Some uses of character literals are ambiguous. Since Ada 95 has introduced Wide_
Character as a new predefined character type, some uses of character literals that
were legal in Ada 83 are illegal in Ada 95. For example:
for Char in 'A' .. 'Z' loop ... end loop;

The problem is that 'A' and 'Z' could be from either Character or Wide_Character. The simplest correction is to make the type explicit; e.g.:

for Char in Character range 'A' .. 'Z' loop ... end loop;

* 'New reserved words'
The identifiers abstract, aliased, protected, requeue, tagged, and until are reserved in Ada 95. Existing Ada 83 code using any of these identifiers must be edited to use some alternative name.

* 'Freezing rules'
The rules in Ada 95 are slightly different with regard to the point at which entities are frozen, and representation pragmas and clauses are not permitted past the freeze point. This shows up most typically in the form of an error message complaining that a representation item appears too late, and the appropriate corrective action is to move the item nearer to the declaration of the entity to which it refers.

A particular case is that representation pragmas cannot be applied to a subprogram body. If necessary, a separate subprogram declaration must be introduced to which the pragma can be applied.

* 'Optional bodies for library packages'
In Ada 83, a package that did not require a package body was nevertheless allowed to have one. This lead to certain surprises in compiling large systems (situations in which the body could be unexpectedly ignored by the binder). In Ada 95, if a package does not require a body then it is not permitted to have a body. To fix this problem, simply remove a redundant body if it is empty, or, if it is non-empty, introduce a dummy declaration into the spec that makes the body required. One approach is to add a private part to the package declaration (if necessary), and define a parameterless procedure called Requires_Body, which must then be given a dummy procedure body in the package body, which then becomes required. Another approach (assuming that this does not introduce elaboration circularities) is to add an Elaborate_Body pragma to the package spec, since one effect of this pragma is to require the presence of a package body.

* 'Numeric_Error is the same exception as Constraint_Error'
In Ada 95, the exception Numeric_Error is a renaming of Constraint_Error. This means that it is illegal to have separate exception handlers for the two exceptions. The fix is simply to remove the handler for the Numeric_Error case (since even in Ada 83, a compiler was free to raise Constraint_Error in place of Numeric_Error in all cases).

* 'Indefinite subtypes in generics'
In Ada 83, it was permissible to pass an indefinite type (e.g, String) as the actual for a generic formal private type, but then the instantiation would be illegal if there were any instances of declarations of variables of this type in the generic body. In Ada 95, to avoid this clear violation of the methodological principle known as the ‘contract model’, the generic declaration explicitly indicates whether or not such instantiations are permitted. If a generic formal parameter has explicit unknown discriminants, indicated by using (<>) after the subtype name, then it can be instantiated with indefinite types, but no stand-alone variables can be declared of this type. Any attempt to declare such
a variable will result in an illegality at the time the generic is declared. If the (<>)
notation is not used, then it is illegal to instantiate the generic with an indefinite type.
This is the potential incompatibility issue when porting Ada 83 code to Ada 95. It will
show up as a compile time error, and the fix is usually simply to add the (<>)
to the generic declaration.

19.2.2 More deterministic semantics

* ‘Conversions’
Conversions from real types to integer types round away from 0. In Ada 83 the conver-
sion Integer(2.5) could deliver either 2 or 3 as its value. This implementation freedom
was intended to support unbiased rounding in statistical applications, but in practice
it interfered with portability. In Ada 95 the conversion semantics are unambiguous,
and rounding away from 0 is required. Numeric code may be affected by this change
in semantics. Note, though, that this issue is no worse than already existed in Ada 83
when porting code from one vendor to another.

* ‘Tasking’
The Real-Time Annex introduces a set of policies that define the behavior of features
that were implementation dependent in Ada 83, such as the order in which open select
branches are executed.

19.2.3 Changed semantics
The worst kind of incompatibility is one where a program that is legal in Ada 83 is also legal
in Ada 95 but can have an effect in Ada 95 that was not possible in Ada 83. Fortunately
this is extremely rare, but the one situation that you should be alert to is the change in the
predefined type Character from 7-bit ASCII to 8-bit Latin-1.

* ‘Range of type ‘Character’’
The range of Standard.Character is now the full 256 characters of Latin-1, whereas
in most Ada 83 implementations it was restricted to 128 characters. Although some
of the effects of this change will be manifest in compile-time rejection of legal Ada
83 programs it is possible for a working Ada 83 program to have a different effect
in Ada 95, one that was not permitted in Ada 83. As an example, the expression
Character’Pos(Character’Last) returned 127 in Ada 83 and now delivers 255 as
its value. In general, you should look at the logic of any character-processing Ada
83 program and see whether it needs to be adapted to work correctly with Latin-1.
Note that the predefined Ada 95 API has a character handling package that may be
relevant if code needs to be adapted to account for the additional Latin-1 elements.
The desirable fix is to modify the program to accommodate the full character set, but
in some cases it may be convenient to define a subtype or derived type of Character
that covers only the restricted range.

19.2.4 Other language compatibility issues

* ‘-gnat83’ switch
All implementations of GNAT provide a switch that causes GNAT to operate in Ada
83 mode. In this mode, some but not all compatibility problems of the type described
above are handled automatically. For example, the new reserved words introduced
in Ada 95 and Ada 2005 are treated simply as identifiers as in Ada 83. However, in practice, it is usually advisable to make the necessary modifications to the program to remove the need for using this switch. See the Compiling Different Versions of Ada section in the GNAT User’s Guide.

* Support for removed Ada 83 pragmas and attributes

A number of pragmas and attributes from Ada 83 were removed from Ada 95, generally because they were replaced by other mechanisms. Ada 95 and Ada 2005 compilers are allowed, but not required, to implement these missing elements. In contrast with some other compilers, GNAT implements all such pragmas and attributes, eliminating this compatibility concern. These include pragma Interface and the floating point type attributes (Emax, Mantissa, etc.), among other items.

19.3 Compatibility between Ada 95 and Ada 2005

Although Ada 2005 was designed to be upwards compatible with Ada 95, there are a number of incompatibilities. Several are enumerated below; for a complete description please see the Annotated Ada 2005 Reference Manual, or section 9.1.1 in Rationale for Ada 2005.

* ‘New reserved words.’

The words interface, overriding and synchronized are reserved in Ada 2005. A pre-Ada 2005 program that uses any of these as an identifier will be illegal.

* ‘New declarations in predefined packages.’


* ‘Access parameters.’

A nondispatching subprogram with an access parameter cannot be renamed as a dispatching operation. This was permitted in Ada 95.

* ‘Access types, discriminants, and constraints.’

Rule changes in this area have led to some incompatibilities; for example, constrained subtypes of some access types are not permitted in Ada 2005.

* ‘Aggregates for limited types.’

The allowance of aggregates for limited types in Ada 2005 raises the possibility of ambiguities in legal Ada 95 programs, since additional types now need to be considered in expression resolution.

* ‘Fixed-point multiplication and division.’

Certain expressions involving ‘*’ or ‘/’ for a fixed-point type, which were legal in Ada 95 and invoked the predefined versions of these operations, are now ambiguous. The ambiguity may be resolved either by applying a type conversion to the expression, or by explicitly invoking the operation from package Standard.

* ‘Return-by-reference types.’

The Ada 95 return-by-reference mechanism has been removed. Instead, the user can declare a function returning a value from an anonymous access type.
19.4 Implementation-dependent characteristics

Although the Ada language defines the semantics of each construct as precisely as practical, in some situations (for example for reasons of efficiency, or where the effect is heavily dependent on the host or target platform) the implementation is allowed some freedom. In porting Ada 83 code to GNAT, you need to be aware of whether / how the existing code exercised such implementation dependencies. Such characteristics fall into several categories, and GNAT offers specific support in assisting the transition from certain Ada 83 compilers.

19.4.1 Implementation-defined pragmas

Ada compilers are allowed to supplement the language-defined pragmas, and these are a potential source of non-portability. All GNAT-defined pragmas are described in [Implementation Defined Pragmas], page 4, and these include several that are specifically intended to correspond to other vendors’ Ada 83 pragmas. For migrating from VADS, the pragma Use_VADS_Size may be useful. For compatibility with HP Ada 83, GNAT supplies the pragmas Extend_System, Ident, Inline_Generic, Interface_Name, Passive, Suppress_All, and Volatile. Other relevant pragmas include External and Link_With. Some vendor-specific Ada 83 pragmas (Share_Generic, Subtitle, and Title) are recognized, thus avoiding compiler rejection of units that contain such pragmas; they are not relevant in a GNAT context and hence are not otherwise implemented.

19.4.2 Implementation-defined attributes

Analogous to pragmas, the set of attributes may be extended by an implementation. All GNAT-defined attributes are described in [Implementation Defined Attributes], page 115, and these include several that are specifically intended to correspond to other vendors’ Ada 83 attributes. For migrating from VADS, the attribute VADS_Size may be useful. For compatibility with HP Ada 83, GNAT supplies the attributes Bit, Machine_Size and Type_Class.

19.4.3 Libraries

Vendors may supply libraries to supplement the standard Ada API. If Ada 83 code uses vendor-specific libraries then there are several ways to manage this in Ada 95 and later versions of the standard:

* If the source code for the libraries (specs and bodies) are available, then the libraries can be migrated in the same way as the application.
* If the source code for the specs but not the bodies are available, then you can reimplement the bodies.
* Some features introduced by Ada 95 obviate the need for library support. For example most Ada 83 vendors supplied a package for unsigned integers. The Ada 95 modular type feature is the preferred way to handle this need, so instead of migrating or reimplementing the unsigned integer package it may be preferable to retrofit the application using modular types.

19.4.4 Elaboration order

The implementation can choose any elaboration order consistent with the unit dependency relationship. This freedom means that some orders can result in Program_Error being
raised due to an ‘Access Before Elaboration’: an attempt to invoke a subprogram before its body has been elaborated, or to instantiate a generic before the generic body has been elaborated. By default GNAT attempts to choose a safe order (one that will not encounter access before elaboration problems) by implicitly inserting Elaborate or Elaborate_All pragmas where needed. However, this can lead to the creation of elaboration circularities and a resulting rejection of the program by gnatbind. This issue is thoroughly described in the ‘Elaboration Order Handling in GNAT’ appendix in the GNAT User’s Guide. In brief, there are several ways to deal with this situation:

* Modify the program to eliminate the circularities, e.g., by moving elaboration-time code into explicitly-invoked procedures

* Constrain the elaboration order by including explicit Elaborate_Body or Elaborate pragmas, and then inhibit the generation of implicit Elaborate_All pragmas either globally (as an effect of the ‘-gnatE’ switch) or locally (by selectively suppressing elaboration checks via pragma Suppress(Elaboration_Check) when it is safe to do so).

19.4.5 Target-specific aspects

Low-level applications need to deal with machine addresses, data representations, interfacing with assembler code, and similar issues. If such an Ada 83 application is being ported to different target hardware (for example where the byte endianness has changed) then you will need to carefully examine the program logic; the porting effort will heavily depend on the robustness of the original design. Moreover, Ada 95 (and thus Ada 2005 and Ada 2012) are sometimes incompatible with typical Ada 83 compiler practices regarding implicit packing, the meaning of the Size attribute, and the size of access values. GNAT’s approach to these issues is described in [Representation Clauses], page 320.

19.5 Compatibility with Other Ada Systems

If programs avoid the use of implementation dependent and implementation defined features, as documented in the Ada Reference Manual, there should be a high degree of portability between GNAT and other Ada systems. The following are specific items which have proved troublesome in moving Ada 95 programs from GNAT to other Ada 95 compilers, but do not affect porting code to GNAT. (As of January 2007, GNAT is the only compiler available for Ada 2005; the following issues may or may not arise for Ada 2005 programs when other compilers appear.)

* ‘Ada 83 Pragmas and Attributes’

Ada 95 compilers are allowed, but not required, to implement the missing Ada 83 pragmas and attributes that are no longer defined in Ada 95. GNAT implements all such pragmas and attributes, eliminating this as a compatibility concern, but some other Ada 95 compilers reject these pragmas and attributes.

* ‘Specialized Needs Annexes’

GNAT implements the full set of special needs annexes. At the current time, it is the only Ada 95 compiler to do so. This means that programs making use of these features may not be portable to other Ada 95 compilation systems.

* ‘Representation Clauses’
Some other Ada 95 compilers implement only the minimal set of representation clauses required by the Ada 95 reference manual. GNAT goes far beyond this minimal set, as described in the next section.

19.6 Representation Clauses

The Ada 83 reference manual was quite vague in describing both the minimal required implementation of representation clauses, and also their precise effects. Ada 95 (and thus also Ada 2005) are much more explicit, but the minimal set of capabilities required is still quite limited.

GNAT implements the full required set of capabilities in Ada 95 and Ada 2005, but also goes much further, and in particular an effort has been made to be compatible with existing Ada 83 usage to the greatest extent possible.

A few cases exist in which Ada 83 compiler behavior is incompatible with the requirements in Ada 95 (and thus also Ada 2005). These are instances of intentional or accidental dependence on specific implementation dependent characteristics of these Ada 83 compilers. The following is a list of the cases most likely to arise in existing Ada 83 code.

* ‘Implicit Packing’

Some Ada 83 compilers allowed a Size specification to cause implicit packing of an array or record. This could cause expensive implicit conversions for change of representation in the presence of derived types, and the Ada design intends to avoid this possibility. Subsequent AI’s were issued to make it clear that such implicit change of representation in response to a Size clause is inadvisable, and this recommendation is represented explicitly in the Ada 95 (and Ada 2005) Reference Manuals as implementation advice that is followed by GNAT. The problem will show up as an error message rejecting the size clause. The fix is simply to provide the explicit pragma Pack, or for more fine tuned control, provide a Component_Size clause.

* ‘Meaning of Size Attribute’

The Size attribute in Ada 95 (and Ada 2005) for discrete types is defined as the minimal number of bits required to hold values of the type. For example, on a 32-bit machine, the size of Natural will typically be 31 and not 32 (since no sign bit is required). Some Ada 83 compilers gave 31, and some 32 in this situation. This problem will usually show up as a compile time error, but not always. It is a good idea to check all uses of the ‘Size attribute when porting Ada 83 code. The GNAT specific attribute Object_Size can provide a useful way of duplicating the behavior of some Ada 83 compiler systems.

* ‘Size of Access Types’

A common assumption in Ada 83 code is that an access type is in fact a pointer, and that therefore it will be the same size as a System.Address value. This assumption is true for GNAT in most cases with one exception. For the case of a pointer to an unconstrained array type (where the bounds may vary from one value of the access type to another), the default is to use a ‘fat pointer’, which is represented as two separate pointers, one to the bounds, and one to the array. This representation has a number of advantages, including improved efficiency. However, it may cause some difficulties in porting existing Ada 83 code which makes the assumption that, for example, pointers fit in 32 bits on a machine with 32-bit addressing.
To get around this problem, GNAT also permits the use of ‘thin pointers’ for access types in this case (where the designated type is an unconstrained array type). These thin pointers are indeed the same size as a System.Address value. To specify a thin pointer, use a size clause for the type, for example:

```ada
type X is access all String;
for X'Size use Standard'Address_Size;
```

which will cause the type X to be represented using a single pointer. When using this representation, the bounds are right behind the array. This representation is slightly less efficient, and does not allow quite such flexibility in the use of foreign pointers or in using the Unrestricted_Access attribute to create pointers to non-aliased objects. But for any standard portable use of the access type it will work in a functionally correct manner and allow porting of existing code. Note that another way of forcing a thin pointer representation is to use a component size clause for the element size in an array, or a record representation clause for an access field in a record.

See the documentation of Unrestricted_Access in the GNAT RM for a full discussion of possible problems using this attribute in conjunction with thin pointers.

### 19.7 Compatibility with HP Ada 83

All the HP Ada 83 pragmas and attributes are recognized, although only a subset of them can sensibly be implemented. The description of pragmas in [Implementation Defined Pragmas], page 4 indicates whether or not they are applicable to GNAT.

* ‘Default floating-point representation’

  In GNAT, the default floating-point format is IEEE, whereas in HP Ada 83, it is VMS format.

* ‘System’

  the package System in GNAT exactly corresponds to the definition in the Ada 95 reference manual, which means that it excludes many of the HP Ada 83 extensions. However, a separate package Aux_DEC is provided that contains the additional definitions, and a special pragma, Extend_System allows this package to be treated transparently as an extension of package System.
Chapter 20: GNU Free Documentation License

20 GNU Free Documentation License

Version 1.3, 3 November 2008
Everyone is permitted to copy and distribute verbatim copies of this license document, but changing it is not allowed.

‘Preamble’
The purpose of this License is to make a manual, textbook, or other functional and useful document “free” in the sense of freedom: to assure everyone the effective freedom to copy and redistribute it, with or without modifying it, either commercially or noncommercially. Secondarily, this License preserves for the author and publisher a way to get credit for their work, while not being considered responsible for modifications made by others.

This License is a kind of “copyleft”, which means that derivative works of the document must themselves be free in the same sense. It complements the GNU General Public License, which is a copyleft license designed for free software.

We have designed this License in order to use it for manuals for free software, because free software needs free documentation: a free program should come with manuals providing the same freedoms that the software does. But this License is not limited to software manuals; it can be used for any textual work, regardless of subject matter or whether it is published as a printed book. We recommend this License principally for works whose purpose is instruction or reference.

‘1. APPLICABILITY AND DEFINITIONS’
This License applies to any manual or other work, in any medium, that contains a notice placed by the copyright holder saying it can be distributed under the terms of this License. Such a notice grants a world-wide, royalty-free license, unlimited in duration, to use that work under the conditions stated herein. The ‘Document’, below, refers to any such manual or work. Any member of the public is a licensee, and is addressed as “you”. You accept the license if you copy, modify or distribute the work in a way requiring permission under copyright law.

A “Modified Version” of the Document means any work containing the Document or a portion of it, either copied verbatim, or with modifications and/or translated into another language.

A “Secondary Section” is a named appendix or a front-matter section of the Document that deals exclusively with the relationship of the publishers or authors of the Document to the Document’s overall subject (or to related matters) and contains nothing that could fall directly within that overall subject. (Thus, if the Document is in part a textbook of mathematics, a Secondary Section may not explain any mathematics.) The relationship could be a matter of historical connection with the subject or with related matters, or of legal, commercial, philosophical, ethical or political position regarding them.

The “Invariant Sections” are certain Secondary Sections whose titles are designated, as being those of Invariant Sections, in the notice that says that the Document is released under this License. If a section does not fit the above definition of Secondary then it is not allowed to be designated as Invariant. The Document may contain zero Invariant Sections. If the Document does not identify any Invariant Sections then there are none.
The “Cover Texts” are certain short passages of text that are listed, as Front-Cover Texts or Back-Cover Texts, in the notice that says that the Document is released under this License. A Front-Cover Text may be at most 5 words, and a Back-Cover Text may be at most 25 words.

A “Transparent” copy of the Document means a machine-readable copy, represented in a format whose specification is available to the general public, that is suitable for revising the document straightforwardly with generic text editors or (for images composed of pixels) generic paint programs or (for drawings) some widely available drawing editor, and that is suitable for input to text formatters or for automatic translation to a variety of formats suitable for input to text formatters. A copy made in an otherwise Transparent file format whose markup, or absence of markup, has been arranged to thwart or discourage subsequent modification by readers is not Transparent. An image format is not Transparent if used for any substantial amount of text. A copy that is not “Transparent” is called ‘Opaque’.

Examples of suitable formats for Transparent copies include plain ASCII without markup, TeXinfo input format, LaTeX input format, SGML or XML using a publicly available DTD, and standard-conforming simple HTML, PostScript or PDF designed for human modification. Examples of transparent image formats include PNG, XCF and JPG. Opaque formats include proprietary formats that can be read and edited only by proprietary word processors, SGML or XML for which the DTD and/or processing tools are not generally available, and the machine-generated HTML, PostScript or PDF produced by some word processors for output purposes only.

The “Title Page” means, for a printed book, the title page itself, plus such following pages as are needed to hold, legibly, the material this License requires to appear in the title page. For works in formats which do not have any title page as such, “Title Page” means the text near the most prominent appearance of the work’s title, preceding the beginning of the body of the text.

The “publisher” means any person or entity that distributes copies of the Document to the public.

A section “Entitled XYZ” means a named subunit of the Document whose title either is precisely XYZ or contains XYZ in parentheses following text that translates XYZ in another language. (Here XYZ stands for a specific section name mentioned below, such as “Acknowledgements”, “Dedications”, “Endorsements”, or “History”.) To “Preserve the Title” of such a section when you modify the Document means that it remains a section “Entitled XYZ” according to this definition.

The Document may include Warranty Disclaimers next to the notice which states that this License applies to the Document. These Warranty Disclaimers are considered to be included by reference in this License, but only as regards disclaiming warranties: any other implication that these Warranty Disclaimers may have is void and has no effect on the meaning of this License.

‘2. VERBATIM COPYING’

You may copy and distribute the Document in any medium, either commercially or noncommercially, provided that this License, the copyright notices, and the license notice saying this License applies to the Document are reproduced in all copies, and that you add no other conditions whatsoever to those of this License. You may not use technical measures to obstruct or control the reading or further copying of the copies you make or distribute.
However, you may accept compensation in exchange for copies. If you distribute a large enough number of copies you must also follow the conditions in section 3.

You may also lend copies, under the same conditions stated above, and you may publicly display copies.

‘3. COPYING IN QUANTITY’

If you publish printed copies (or copies in media that commonly have printed covers) of the Document, numbering more than 100, and the Document’s license notice requires Cover Texts, you must enclose the copies in covers that carry, clearly and legibly, all these Cover Texts: Front-Cover Texts on the front cover, and Back-Cover Texts on the back cover. Both covers must also clearly and legibly identify you as the publisher of these copies. The front cover must present the full title with all words of the title equally prominent and visible. You may add other material on the covers in addition. Copying with changes limited to the covers, as long as they preserve the title of the Document and satisfy these conditions, can be treated as verbatim copying in other respects.

If the required texts for either cover are too voluminous to fit legibly, you should put the first ones listed (as many as fit reasonably) on the actual cover, and continue the rest onto adjacent pages.

If you publish or distribute Opaque copies of the Document numbering more than 100, you must either include a machine-readable Transparent copy along with each Opaque copy, or state in or with each Opaque copy a computer-network location from which the general network-using public has access to download using public-standard network protocols a complete Transparent copy of the Document, free of added material. If you use the latter option, you must take reasonably prudent steps, when you begin distribution of Opaque copies in quantity, to ensure that this Transparent copy will remain thus accessible at the stated location until at least one year after the last time you distribute an Opaque copy (directly or through your agents or retailers) of that edition to the public.

It is requested, but not required, that you contact the authors of the Document well before redistributing any large number of copies, to give them a chance to provide you with an updated version of the Document.

‘4. MODIFICATIONS’

You may copy and distribute a Modified Version of the Document under the conditions of sections 2 and 3 above, provided that you release the Modified Version under precisely this License, with the Modified Version filling the role of the Document, thus licensing distribution and modification of the Modified Version to whoever possesses a copy of it. In addition, you must do these things in the Modified Version:

A. Use in the Title Page (and on the covers, if any) a title distinct from that of the Document, and from those of previous versions (which should, if there were any, be listed in the History section of the Document). You may use the same title as a previous version if the original publisher of that version gives permission.

B. List on the Title Page, as authors, one or more persons or entities responsible for authorship of the modifications in the Modified Version, together with at least five of the principal authors of the Document (all of its principal authors, if it has fewer than five), unless they release you from this requirement.

C. State on the Title page the name of the publisher of the Modified Version, as the publisher.
D. Preserve all the copyright notices of the Document.
E. Add an appropriate copyright notice for your modifications adjacent to the other copyright notices.
F. Include, immediately after the copyright notices, a license notice giving the public permission to use the Modified Version under the terms of this License, in the form shown in the Addendum below.
G. Preserve in that license notice the full lists of Invariant Sections and required Cover Texts given in the Document’s license notice.
H. Include an unaltered copy of this License.
I. Preserve the section Entitled “History”, Preserve its Title, and add to it an item stating at least the title, year, new authors, and publisher of the Modified Version as given on the Title Page. If there is no section Entitled “History” in the Document, create one stating the title, year, authors, and publisher of the Document as given on its Title Page, then add an item describing the Modified Version as stated in the previous sentence.
J. Preserve the network location, if any, given in the Document for public access to a Transparent copy of the Document, and likewise the network locations given in the Document for previous versions it was based on. These may be placed in the “History” section. You may omit a network location for a work that was published at least four years before the Document itself, or if the original publisher of the version it refers to gives permission.
K. For any section Entitled “Acknowledgements” or “Dedications”, Preserve the Title of the section, and preserve in the section all the substance and tone of each of the contributor acknowledgements and/or dedications given therein.
L. Preserve all the Invariant Sections of the Document, unaltered in their text and in their titles. Section numbers or the equivalent are not considered part of the section titles.
M. Delete any section Entitled “Endorsements”. Such a section may not be included in the Modified Version.
N. Do not retitle any existing section to be Entitled “Endorsements” or to conflict in title with any Invariant Section.
O. Preserve any Warranty Disclaimers.

If the Modified Version includes new front-matter sections or appendices that qualify as Secondary Sections and contain no material copied from the Document, you may at your option designate some or all of these sections as invariant. To do this, add their titles to the list of Invariant Sections in the Modified Version’s license notice. These titles must be distinct from any other section titles.

You may add a section Entitled “Endorsements”, provided it contains nothing but endorsements of your Modified Version by various parties—for example, statements of peer review or that the text has been approved by an organization as the authoritative definition of a standard.

You may add a passage of up to five words as a Front-Cover Text, and a passage of up to 25 words as a Back-Cover Text, to the end of the list of Cover Texts in the Modified Version. Only one passage of Front-Cover Text and one of Back-Cover Text may be added by (or through arrangements made by) any one entity. If the Document already includes
a cover text for the same cover, previously added by you or by arrangement made by the
same entity you are acting on behalf of, you may not add another; but you may replace the
old one, on explicit permission from the previous publisher that added the old one.
The author(s) and publisher(s) of the Document do not by this License give permission to
use their names for publicity for or to assert or imply endorsement of any Modified Version.

5. COMBINING DOCUMENTS
You may combine the Document with other documents released under this License, under
the terms defined in section 4 above for modified versions, provided that you include in the
combination all of the Invariant Sections of all of the original documents, unmodified, and
list them all as Invariant Sections of your combined work in its license notice, and that you
preserve all their Warranty Disclaimers.

The combined work need only contain one copy of this License, and multiple identical
Invariant Sections may be replaced with a single copy. If there are multiple Invariant
Sections with the same name but different contents, make the title of each such section
unique by adding at the end of it, in parentheses, the name of the original author or
publisher of that section if known, or else a unique number. Make the same adjustment to
the section titles in the list of Invariant Sections in the license notice of the combined work.

In the combination, you must combine any sections Entitled “History” in the various original
documents, forming one section Entitled “History”; likewise combine any sections Entitled
“Acknowledgements”, and any sections Entitled “Dedications”. You must delete all sections
Entitled “Endorsements”.

6. COLLECTIONS OF DOCUMENTS
You may make a collection consisting of the Document and other documents released under
this License, and replace the individual copies of this License in the various documents with
a single copy that is included in the collection, provided that you follow the rules of this
License for verbatim copying of each of the documents in all other respects.

You may extract a single document from such a collection, and distribute it individually
under this License, provided you insert a copy of this License into the extracted document,
and follow this License in all other respects regarding verbatim copying of that document.

7. AGGREGATION WITH INDEPENDENT WORKS
A compilation of the Document or its derivatives with other separate and independent
documents or works, in or on a volume of a storage or distribution medium, is called an
“aggregate” if the copyright resulting from the compilation is not used to limit the legal
rights of the compilation’s users beyond what the individual works permit. When the
Document is included in an aggregate, this License does not apply to the other works in
the aggregate which are not themselves derivative works of the Document.

If the Cover Text requirement of section 3 is applicable to these copies of the Document,
then if the Document is less than one half of the entire aggregate, the Document’s Cover
Texts may be placed on covers that bracket the Document within the aggregate, or the
electronic equivalent of covers if the Document is in electronic form. Otherwise they must
appear on printed covers that bracket the whole aggregate.

8. TRANSLATION
Translation is considered a kind of modification, so you may distribute translations of the
Document under the terms of section 4. Replacing Invariant Sections with translations
requires special permission from their copyright holders, but you may include translations of some or all Invariant Sections in addition to the original versions of these Invariant Sections. You may include a translation of this License, and all the license notices in the Document, and any Warranty Disclaimers, provided that you also include the original English version of this License and the original versions of those notices and disclaimers. In case of a disagreement between the translation and the original version of this License or a notice or disclaimer, the original version will prevail.

If a section in the Document is Entitled “Acknowledgements”, “Dedications”, or “History”, the requirement (section 4) to Preserve its Title (section 1) will typically require changing the actual title.

‘9. TERMINATION’

You may not copy, modify, sublicense, or distribute the Document except as expressly provided under this License. Any attempt otherwise to copy, modify, sublicense, or distribute it is void, and will automatically terminate your rights under this License.

However, if you cease all violation of this License, then your license from a particular copyright holder is reinstated (a) provisionally, unless and until the copyright holder explicitly and finally terminates your license, and (b) permanently, if the copyright holder fails to notify you of the violation by some reasonable means prior to 60 days after the cessation. Moreover, your license from a particular copyright holder is reinstated permanently if the copyright holder notifies you of the violation by some reasonable means, this is the first time you have received notice of violation of this License (for any work) from that copyright holder, and you cure the violation prior to 30 days after your receipt of the notice.

Termination of your rights under this section does not terminate the licenses of parties who have received copies or rights from you under this License. If your rights have been terminated and not permanently reinstated, receipt of a copy of some or all of the same material does not give you any rights to use it.

‘10. FUTURE REVISIONS OF THIS LICENSE’

The Free Software Foundation may publish new, revised versions of the GNU Free Documentation License from time to time. Such new versions will be similar in spirit to the present version, but may differ in detail to address new problems or concerns. See https://www.gnu.org/copyleft/.

Each version of the License is given a distinguishing version number. If the Document specifies that a particular numbered version of this License “or any later version” applies to it, you have the option of following the terms and conditions either of that specified version or of any later version that has been published (not as a draft) by the Free Software Foundation. If the Document does not specify a version number of this License, you may choose any version ever published (not as a draft) by the Free Software Foundation. If the Document specifies that a proxy can decide which future versions of this License can be used, that proxy’s public statement of acceptance of a version permanently authorizes you to choose that version for the Document.

‘11. RELICENSING’

“Massive Multiauthor Collaboration Site” (or “MMC Site”) means any World Wide Web server that publishes copyrightable works and also provides prominent facilities for anybody to edit those works. A public wiki that anybody can edit is an example of such a server. A
“Massive Multiauthor Collaboration” (or “MMC”) contained in the site means any set of copyrightable works thus published on the MMC site.

“CC-BY-SA” means the Creative Commons Attribution-Share Alike 3.0 license published by Creative Commons Corporation, a not-for-profit corporation with a principal place of business in San Francisco, California, as well as future copyleft versions of that license published by that same organization.

“Incorporate” means to publish or republish a Document, in whole or in part, as part of another Document.

An MMC is “eligible for relicensing” if it is licensed under this License, and if all works that were first published under this License somewhere other than this MMC, and subsequently incorporated in whole or in part into the MMC, (1) had no cover texts or invariant sections, and (2) were thus incorporated prior to November 1, 2008.

The operator of an MMC Site may republish an MMC contained in the site under CC-BY-SA on the same site at any time before August 1, 2009, provided the MMC is eligible for relicensing.

‘ADDENDUM: How to use this License for your documents’

To use this License in a document you have written, include a copy of the License in the document and put the following copyright and license notices just after the title page:

Copyright © YEAR YOUR NAME. Permission is granted to copy, distribute and/or modify this document under the terms of the GNU Free Documentation License, Version 1.3 or any later version published by the Free Software Foundation; with no Invariant Sections, no Front-Cover Texts, and no Back-Cover Texts. A copy of the license is included in the section entitled “GNU Free Documentation License”.

If you have Invariant Sections, Front-Cover Texts and Back-Cover Texts, replace the “with . . . Texts.” line with this:

with the Invariant Sections being LIST THEIR TITLES, with the Front-Cover Texts being LIST, and with the Back-Cover Texts being LIST.

If you have Invariant Sections without Cover Texts, or some other combination of the three, merge those two alternatives to suit the situation.

If your document contains nontrivial examples of program code, we recommend releasing these examples in parallel under your choice of free software license, such as the GNU General Public License, to permit their use in free software.
Index

-  
  -gnat12 option (gcc) .................................. 287
  -gnatR (gcc) ........................................ 223
  
  __lock file (for shared passive packages) ........ 282

A
  
  Abort_Signal ........................................ 116
  Abstract_State ...................................... 108
  Access .............................................. 133
  Access values ...................................... 121
  Accuracy ........................................... 172
  Accuracy requirements .............................. 172
  Ada 2012 implementation status .................. 287
  Ada 83 attributes ................................ 119, 120, 123, 124, 127, 131
  Ada 95 Language Reference Manual .............. 3
  Ada Extensions ..................................... 34
  Ada.Characters.Handling ........................... 163
  Ada.Characters.Wide_Latin_1 (a-cwila1.ads) .... 255
  Ada.Characters.Wide_Wide_Latin_1 (a-chzla1.ads)
     .................................................. 255
     .................................................. 255
  Ada.Command_Line.Environment (a-colien.ads)
     .................................................. 256
  Ada.Command_Line.Remove (a-colire.ads) ........ 256
     .................................................. 256
  Ada.Containers.Bounded_Holders (a-coboho.ads)
     .................................................. 256
  Ada.Direct_IO.C_Streams (a-diocest.ads) ....... 256
  Ada.Exceptions.Is_Null_Occurrence (a-einoc.ads)
     .................................................. 256
  Ada.Exceptions.Last_Chance_Handler
     (a-elchha.ads) .................................. 256
  Ada.Exceptions.Traceback (a-extrax.ads) ......... 257
  Ada.Sequential_IO.C_Streams (a-siocest.ads) ... 257
  Ada.Streams.Stream_IO.C_Streams (a-ssicst.ads)
     .................................................. 257
  Ada.Strings.Unbounded.Text_IO (a-suteio.ads)
     .................................................. 257
  Ada.Strings.Wide_Unbounded.Wide_Text_IO
     (a-swuuti.ads) .................................. 257
  Ada.Strings.Wide_Wide_Unbounded.Wide_Wide_Text
     .................................................. 257
  Ada.Text_IO.C_Streams (a-tiocst.ads) ............ 257

  Ada.Text_IO.Reset_Standard_Files (a-tirfsi.ads)
     .................................................. 257
  Ada.Wide_Characters.Unicode (a-wichuni.ads)
     .................................................. 257
  Ada.Wide_Text_IO.C_Streams (a-wtcstr.ads)
     .................................................. 257
  Ada.Wide_Text_IO.Reset_Standard_Files
     (a-wrstfi.ads) .................................. 257
  Ada.Wide_Wide_Characters.Unicode
     (a-zchuni.ads) .................................. 257
  Ada.Wide_Wide_Text_IO.C_Streams (a-ztcstr.ads)
     .................................................. 257
  Ada.Wide_Wide_Text_IO.Reset_Standard_Files
     (a-zrstfi.ads) .................................. 257
  Ada_2012 configuration pragma .................... 287
  Address ............................................ 161
  Address Clause .................................... 215
  Address clauses ................................... 158
  Address image ..................................... 270
  Address of subprogram code ....................... 177
  Address_Size ...................................... 116
  AI-0002 (Ada 2012 feature) ....................... 302
  AI-0003 (Ada 2012 feature) ....................... 290
  AI-0007 (Ada 2012 feature) ....................... 300
  AI-0008 (Ada 2012 feature) ....................... 290
  AI-0009 (Ada 2012 feature) ....................... 297
  AI-0012 (Ada 2012 feature) ....................... 299
  AI-0015 (Ada 2012 feature) ....................... 295
  AI-0017 (Ada 2012 feature) ....................... 301
  AI-0019 (Ada 2012 feature) ....................... 301
  AI-0026 (Ada 2012 feature) ....................... 296
  AI-0030 (Ada 2012 feature) ....................... 297
  AI-0031 (Ada 2012 feature) ....................... 301
  AI-0032 (Ada 2012 feature) ....................... 295
  AI-0033 (Ada 2012 feature) ....................... 300
  AI-0034 (Ada 2012 feature) ....................... 298
  AI-0035 (Ada 2012 feature) ....................... 298
  AI-0037 (Ada 2012 feature) ....................... 293
  AI-0038 (Ada 2012 feature) ....................... 302
  AI-0039 (Ada 2012 feature) ....................... 299
  AI-0040 (Ada 2012 feature) ....................... 298
  AI-0042 (Ada 2012 feature) ....................... 296
  AI-0043 (Ada 2012 feature) ....................... 298
  AI-0044 (Ada 2012 feature) ....................... 302
  AI-0046 (Ada 2012 feature) ....................... 295
  AI-0050 (Ada 2012 feature) ....................... 295
  AI-0056 (Ada 2012 feature) ....................... 301
  AI-0058 (Ada 2012 feature) ....................... 295
  AI-0060 (Ada 2012 feature) ....................... 301
  AI-0062 (Ada 2012 feature) ....................... 296
  AI-0064 (Ada 2012 feature) ....................... 296
  AI-0070 (Ada 2012 feature) ....................... 292
  AI-0072 (Ada 2012 feature) ....................... 297
  AI-0073 (Ada 2012 feature) ....................... 292
<table>
<thead>
<tr>
<th>Index</th>
<th>330</th>
</tr>
</thead>
</table>

| AI-0076 (Ada 2012 feature) | 292 |
| AI-0077 (Ada 2012 feature) | 297 |
| AI-0078 (Ada 2012 feature) | 299 |
| AI-0079 (Ada 2012 feature) | 287 |
| AI-0080 (Ada 2012 feature) | 288 |
| AI-0087 (Ada 2012 feature) | 296 |
| AI-0088 (Ada 2012 feature) | 294 |
| AI-0091 (Ada 2012 feature) | 287 |
| AI-0093 (Ada 2012 feature) | 291 |
| AI-0095 (Ada 2012 feature) | 299 |
| AI-0096 (Ada 2012 feature) | 291 |
| AI-0097 (Ada 2012 feature) | 292 |
| AI-0098 (Ada 2012 feature) | 293 |
| AI-0099 (Ada 2012 feature) | 296 |
| AI-0100 (Ada 2012 feature) | 287 |
| AI-0102 (Ada 2012 feature) | 291 |
| AI-0103 (Ada 2012 feature) | 295 |
| AI-0104 (Ada 2012 feature) | 294 |
| AI-0106 (Ada 2012 feature) | 299 |
| AI-0108 (Ada 2012 feature) | 297 |
| AI-0109 (Ada 2012 feature) | 300 |
| AI-0112 (Ada 2012 feature) | 298 |
| AI-0114 (Ada 2012 feature) | 301 |
| AI-0116 (Ada 2012 feature) | 299 |
| AI-0118 (Ada 2012 feature) | 295 |
| AI-0120 (Ada 2012 feature) | 290 |
| AI-0122 (Ada 2012 feature) | 298 |
| AI-0123 (Ada 2012 feature) | 293 |
| AI-0125 (Ada 2012 feature) | 295 |
| AI-0126 (Ada 2012 feature) | 292 |
| AI-0127 (Ada 2012 feature) | 302 |
| AI-0128 (Ada 2012 feature) | 290 |
| AI-0129 (Ada 2012 feature) | 297 |
| AI-0132 (Ada 2012 feature) | 298 |
| AI-0133 (Ada 2012 feature) | 294 |
| AI-0137 (Ada 2012 feature) | 302 |
| AI-0139-2 (Ada 2012 feature) | 294 |
| AI-0146 (Ada 2012 feature) | 299 |
| AI-0147 (Ada 2012 feature) | 293 |
| AI-0152 (Ada 2012 feature) | 303 |
| AI-0157 (Ada 2012 feature) | 294 |
| AI-0158 (Ada 2012 feature) | 291 |
| AI-0161 (Ada 2012 feature) | 300 |
| AI-0162 (Ada 2012 feature) | 293 |
| AI-0163 (Ada 2012 feature) | 288 |
| AI-0171 (Ada 2012 feature) | 303 |
| AI-0173 (Ada 2012 feature) | 291 |
| AI-0176 (Ada 2012 feature) | 287 |
| AI-0177 (Ada 2012 feature) | 300 |
| AI-0178 (Ada 2012 feature) | 296 |
| AI-0179 (Ada 2012 feature) | 294 |
| AI-0181 (Ada 2012 feature) | 291 |
| AI-0182 (Ada 2012 feature) | 291 |
| AI-0183 (Ada 2012 feature) | 288 |
| AI-0185 (Ada 2012 feature) | 301 |
| AI-0188 (Ada 2012 feature) | 294 |
| AI-0189 (Ada 2012 feature) | 303 |
| AI-0190 (Ada 2012 feature) | 303 |

| AI-0193 (Ada 2012 feature) | 300 |
| AI-0194 (Ada 2012 feature) | 300 |
| AI-0195 (Ada 2012 feature) | 299 |
| AI-0196 (Ada 2012 feature) | 295 |
| AI-0198 (Ada 2012 feature) | 292 |
| AI-0199 (Ada 2012 feature) | 293 |
| AI-0200 (Ada 2012 feature) | 298 |
| AI-0201 (Ada 2012 feature) | 297 |
| AI-0203 (Ada 2012 feature) | 292 |
| AI-0205 (Ada 2012 feature) | 296 |
| AI-0206 (Ada 2012 feature) | 303 |
| AI-0207 (Ada 2012 feature) | 294 |
| AI-0208 (Ada 2012 feature) | 292 |
| AI-0210 (Ada 2012 feature) | 303 |
| AI-0211 (Ada 2012 feature) | 303 |
| AI-0214 (Ada 2012 feature) | 291 |
| AI-0219 (Ada 2012 feature) | 298 |
| AI-0220 (Ada 2012 feature) | 293 |
| AI05-0216 (Ada 2012 feature) | 302 |
| Alignment | 64, 124, 131, 196 |
| Alignment Clause | 195 |
| Alignment clauses | 158 |
| Alignments of components | 19 |
| allocator | 131 |
| Alternative Character Sets | 154 |
| AltiVec | 258, 259 |
| Annex E | 281 |
| Annotate | 108 |
| Anonymous access types | 221 |
| Argument passing mechanisms | 30 |
| argument removal | 256 |
| Array packing | 43 |
| Array splitter | 259 |
| Arrays | 155, 262, 268 |
| as private type | 161 |
| Asm_Input | 116 |
| Asm_Output | 116 |
| Assert_Failure | 270 |
| Assertions | 14, 16, 270 |
| Async_Readers | 109 |
| Async_Writers | 109 |
| Atomic Synchronization | 26, 29 |
| Atomic_Always_Lock_Free | 117 |
| Attribute | 216 |
| Attribute Loop_Entry | 96 |
| Attribute Old | 96 |
| AWK | 259 |

B

Biased representation | 200
Big endian | 118
Binary search | 259
Bind environment | 259
Bit | 117
Bit ordering | 161
bit ordering | 204
Bit_Order Clause | 204
Index

GNAT.Heap_Sort_G (g-hesorg.ads) ............. 264
GNAT.HTable (g-htable.ads) ................. 264
GNAT.IO (g-io.ads) .......................... 264
GNAT.IO_Aux (g-io_aux.ads) .................. 264
GNAT.Lock_Files (g-locfil.ads) ............... 267
GNAT.MBBS_Discrete_Random (g-mbdira.ads) ... 267
GNAT.MBBS_Float_Random (g-mbdira.ads) ... 267
GNAT.IO (g-io.ads) .......................... 267
GNAT.MD5 (g-md5.ads) ........................ 267
GNAT.Memory_Dump (g-memdump.ads) .......... 267
GNAT.Most_Recent_Exception (g-moreex.ads) 267
GNAT.OS_Lib (g-os_lib.ads) ................. 267
GNAT.Perfect_Hash_Generators (g-phehga.ads) 267
GNAT.Random_Numbers (g-rannum.ads) ....... 267
GNAT.Regexp (g-regexp.ads) ................. 267
GNAT.Regist (g-regist.ads) ................... 267
GNAT.RegsMap (g-regpsmt.ads) ............... 267
GNAT.Rewrite_Data (g-rewdat.ads) .......... 267
GNAT.Secondary_Stack_Info (g-sestin.ads) 267
GNAT.Semaphores (g-semaph.ads) ............. 268
GNAT.Serial_Communications (g-sercom.ads) 268
GNAT.SHA1 (g-sha1.ads) ........................ 268
GNAT.SHA24 (g-sha24.ads) .................... 268
GNAT.SHA256 (g-sha256.ads) .................. 268
GNAT.SHA384 (g-sha384.ads) .................. 268
GNAT.SHA512 (g-sha512.ads) .................. 268
GNAT.Signals (g-signal.ads) ................. 268
GNAT.Sockets (g-socket.ads) ................. 268
GNAT.Source_Info (g-sourinf.ads) ........... 268
GNAT.Spelling_Checker (g-speche.ads) ...... 268
GNAT.Spelling_Checker_Generic (g-spehge.ads) 268
GNAT.Spitbol (g-sptbols.ads) ............... 268
GNAT.Spitbol.Patterns (g-spitpat.ads) ..... 268
GNAT.Spitbol.Table_Boolean (g-sptabo.ads) 268
GNAT.Spitbol.Table_Integer (g-sptain.ads) 268
GNAT.Spitbol.Table_VString (g-sptavs.ads) 268
GNAT.SSE (g-sse.ads) ......................... 268
GNAT.SSE.Vector_Types (g-svetyt.ads) ..... 269
GNAT.String_Hash (g-strhsh.ads) ........... 269
GNAT.String_Split (g-strspl.ads) ............. 269
GNAT.Strings (g-string.ads) .................. 269
GNAT.Table (g-table.ads) ..................... 269
GNAT.Task_Lock (g-tasloc.ads) .............. 269
GNAT.Threads (g-thread.ads) ................. 269
GNAT.Time_Stamp (g-timsta.ads) .............. 269
GNAT.Traceback (g-traceb.ads) .............. 269
GNAT.Traceback_Symbolic (g-trasym.ads) .... 269
GNAT.UTF_32 (g-utf_32.ads) .................. 269
GNAT.UTF_32_Spelling_Checker (g-u3spch.ads) 269
GNAT.Wide_Spelling_Checker (g-wispch.ads) 269
GNAT.Wide_String_Split (g-wistps.ads) ..... 269

GNAT.Wide_Wide_Spelling_Checker
(g-wispche.ads) ............................... 269
GNAT.Wide_Wide_String_Split (g-wistsp.ads) 270

H

handling long command lines .................. 256
Handling of Records with Holes ............... 213
Has_Access_Values .. .......................... 121
Has_Discriminators ................................ 121
Has_Tagged_Values .................................. 122
Hash functions .................................... 265, 268
Hash tables ....................................... 262, 264
Heap usage ........................................ 162

I

I/O interfacing ................................... 270
IBM Packed Format .................................. 270
Image .................................................. 270
Img ...................................................... 122
Immediate_Reclamation .......................... 139
Implementation-dependent features ............. 2
implicit .............................................. 162
Import ............................................... 217
Initial_Condition .................................. 111
Initialization ...................................... 92
Initialized ......................................... 122
Initializes .......................................... 111
Inline_Aways ....................................... 111
Input/Output facilities .......................... 264
Integer maps ........................................ 268
Integer types ....................................... 155
Integer_Value ...................................... 122
Interfaces ......................................... 164
Interfaces.C.Extensions (i-cexten.ads) .... 270
Interfaces.C.Streams (i-cstrea.ads) ......... 270
Interfaces.Packed_Decimal (i-pacedec.ads) 270
Interfaces.VxWorks (i-vxwork.ads) .......... 270
Interfaces.VxWorks.Int_Connection (i-vxinco.ads) 270

Interfaces.VxWorks.IO (i-vxwoio.ads) ....... 270
interfacing ......................................... 270
Interfacing to C++ ................................. 23, 77
Interfacing to VxWorks .......................... 270
Interfacing to VxWorks’ I/O .................... 270
interfacing with .................................. 165, 166
Interfacing with “Text_IO” .................. 257
Interfacing with “Wide_Text_IO” ............. 258
Interfacing with “Wide_Wide_Text_IO” ...... 258
Interfacing with C++ ............................. 22, 23
Interfacing with Direct_IO ..................... 256
Interfacing with Sequential_IO ............... 257
Interfacing with Stream_IO ..................... 257
Interrupt ............................................. 261
Interrupt support .................................. 167
Interrupts .......................................... 168
Intrinsic operator ........................................ 192
Intrinsic Subprograms ................................. 192
Invalid representations ............................... 13
Invalid values ........................................ 13
Invalid_Value ........................................... 122
Invariant ................................................ 111
Invariant’Class ........................................... 111
IO support .............................................. 257
Iterable ............................................... 112, 123

L
Large ..................................................... 123
Latin-1 .................................................... 316
Latin_1 constants for Wide_Character ............. 255
Latin_1 constants for WideWide_Character ....... 255
Latin_9 constants for Character .................... 255
Latin_9 constants for Wide_Character ............. 255
Latin_9 constants for WideWide_Character ....... 256
Library_Level ........................................... 123
License checking ....................................... 51
Line ....................................................... 193
Linker_Section ......................................... 112
Little endian ............................................ 118
local ...................................................... 271
Local storage pool .................................... 271
Lock_Free ............................................... 112
Locking .................................................. 268
Locking_Policies ....................................... 169
Locking using files .................................... 264
Loop_Entry ............................................. 123

M
Machine Code insertions .............................. 277
Machine operations ................................... 166
Machine_Size .......................................... 123
Mailboxes .............................................. 259
Mantissa ............................................... 124
Maps ...................................................... 268
Mathematical functions ................................ 263
Max_Asyncous_Select_Nesting ....................... 139
Max_Entry_Queue_Depth ................................ 139
Max_Entry_Queue_Length ................................ 139
Max_Integer_Size ...................................... 124
Max_Protected_Entries ................................ 139
Max_Queue_Length ...................................... 113
Max_Select_Alanatives ................................ 139
Max_Storage_At_Blocking .............................. 140
Max_Task_Entries ...................................... 140
Max_Tasks .............................................. 140
maximum ............................................... 124
Maximum_Alignment ................................... 124
Maximum_Alignment attribute ....................... 195
Mechanism_Code ....................................... 124
Memory allocation ..................................... 271
Memory corruption debugging ....................... 261
Memory-mapped I/O .................................... 220
Message Digest MD5 .................................... 265
monotonic .............................................. 169
multidimensional ...................................... 155
Multidimensional arrays ............................. 155
Multiprocessor interface ............................. 271

N
Named assertions ....................................... 14, 16
Named numbers ........................................ 133
NoAbort_Statements .................................. 140
NoAccess_Parameter_Allocators ..................... 140
NoAccess_Subprograms ................................ 140
NoAllocators ......................................... 140
NoAnonymous_Allocators .............................. 140
NoAsynchronous_Control .............................. 140
NoCaching ............................................. 113
NoCalendar ............................................ 140
NoCoextensions ....................................... 140
NoDefault_Initialization ............................. 141
NoDelay ................................................. 141
NoDependence ........................................ 141
NoDirect_Boolean_Operators ......................... 141
NoDispatch ............................................. 141
NoDispatching_Calls .................................. 141
NoDynamic_Accessibility_Checks .................... 150
NoDynamic_Attachment ................................ 142
NoDynamicInterruptedException ...................... 142
NoDynamic_Priorities ................................ 142
NoDynamic_Sized_Objects ............................. 151
NoElimilation_Code ................................... 149
NoElimilation_Code_All ............................... 113
NoEntry_Calls_In_Elaboration_Code .................. 142
NoEntry_Queue ......................................... 151
NoEnumeration_Maps .................................. 143
NoException_Handlers ................................ 143
NoException_Propagation ............................. 143
NoException_Registration ............................ 143
NoExceptions ........................................... 143
NoFinalization ........................................ 143
NoFixed_Point ......................................... 144
NoFloating_Point ...................................... 144
NoImplementation_Aspect_Specifications .......... 151
NoImplementation_Attributes ....................... 151
NoImplementation_Identifiers ....................... 151
NoImplementation_Pragmas ........................... 151
NoImplementation_Restrictions ...................... 151
NoImplementation_Units ................................ 152
NoImplicit_Aliasing ................................... 152
NoImplicit_Conditionals ............................ 144
NoImplicit_Dynamic_Code ............................. 144
NoImplicit_Heap_Allocations ........................ 144
NoImplicit_Loops ....................................... 152
NoImplicit_Protected_Object_Allocations .......... 145
NoImplicit_Task_Allocations ........................ 145
NoInitializeScalars ................................... 145
NoInline ............................................... 113
NoIO ................................................... 145
<table>
<thead>
<tr>
<th>Real-Time Systems Annex compliance</th>
<th>Shift_Right</th>
<th>Shift_Right_Arithmetic</th>
</tr>
</thead>
<tbody>
<tr>
<td>Record Representation Clause</td>
<td>Signals</td>
<td>85, 130</td>
</tr>
<tr>
<td>Record representation clauses</td>
<td>simple</td>
<td>264</td>
</tr>
<tr>
<td>records</td>
<td>Simple I/O</td>
<td>85, 130</td>
</tr>
<tr>
<td>Refined_Deps</td>
<td>Simple storage pool</td>
<td>149</td>
</tr>
<tr>
<td>Refined_Global</td>
<td>Simple_Barriers</td>
<td>114, 130</td>
</tr>
<tr>
<td>Refined_Initialization</td>
<td>Simple_Storage_Pool</td>
<td>114</td>
</tr>
<tr>
<td>Refined_Post</td>
<td>Simple_Storage_Pool_Type</td>
<td>114</td>
</tr>
<tr>
<td>Refined_State</td>
<td>Size ..........</td>
<td>100, 124, 138, 198, 200</td>
</tr>
<tr>
<td>Regular expressions</td>
<td>size ..........</td>
<td>198</td>
</tr>
<tr>
<td>Remote_Access_Type</td>
<td>Size_Clause</td>
<td>196</td>
</tr>
<tr>
<td>Removing command line arguments</td>
<td>Size clauses</td>
<td>159</td>
</tr>
<tr>
<td>Representation</td>
<td>Size for biased representation</td>
<td>200</td>
</tr>
<tr>
<td>representation</td>
<td>Size of &quot;Address&quot;</td>
<td>116</td>
</tr>
<tr>
<td>Representation Clause</td>
<td>Small ..........</td>
<td>131</td>
</tr>
<tr>
<td>Representation clauses</td>
<td>Small_Denominator</td>
<td>131</td>
</tr>
<tr>
<td>representation of</td>
<td>Small_Numerator</td>
<td>131</td>
</tr>
<tr>
<td>Representation of enums</td>
<td>Sockets</td>
<td>267</td>
</tr>
<tr>
<td>Representation of wide characters</td>
<td>Sorting ..........</td>
<td>259, 260, 263, 264</td>
</tr>
<tr>
<td>RepresentationPragma</td>
<td>Source Information</td>
<td>267</td>
</tr>
<tr>
<td>response file</td>
<td>Source_Location</td>
<td>194</td>
</tr>
<tr>
<td>Response file for command line</td>
<td>SPARK_05</td>
<td>152</td>
</tr>
<tr>
<td>Restriction_Set</td>
<td>SPARK_Mode</td>
<td>114</td>
</tr>
<tr>
<td>Restrictions</td>
<td>Spawn capability</td>
<td>265</td>
</tr>
<tr>
<td>Restrictions definitions</td>
<td>Spell checking</td>
<td>267, 269</td>
</tr>
<tr>
<td>Result</td>
<td>SPITBOL interface</td>
<td>267</td>
</tr>
<tr>
<td>Return values</td>
<td>SPITBOL pattern matching</td>
<td>267</td>
</tr>
<tr>
<td>Rewrite data</td>
<td>SPITBOL Tables</td>
<td>267, 268</td>
</tr>
<tr>
<td>Rotate_Left</td>
<td>Static_Dispatch_Tables</td>
<td>152</td>
</tr>
<tr>
<td>Rotate_Right</td>
<td>Static_Priorities</td>
<td>149</td>
</tr>
<tr>
<td>Run-time restrictions access</td>
<td>Static_Storage_Size</td>
<td>149</td>
</tr>
<tr>
<td>Storage place attributes</td>
<td>Storage pool</td>
<td>85, 130, 271</td>
</tr>
<tr>
<td>Storage_Size Clause</td>
<td>Storage_Unit</td>
<td>131</td>
</tr>
<tr>
<td>Storage_Unit (in pragma Component_Alignment)</td>
<td>Storage_Unit (in pragma Component_Alignment)</td>
<td>19</td>
</tr>
<tr>
<td>Scalar storage order</td>
<td>String decoding</td>
<td>261</td>
</tr>
<tr>
<td>Scalar_Storage_Order</td>
<td>String encoding</td>
<td>262</td>
</tr>
<tr>
<td>Secondary Stack Info</td>
<td>String maps</td>
<td>268</td>
</tr>
<tr>
<td>Secondary_STACK_Size</td>
<td>String splitter</td>
<td>268</td>
</tr>
<tr>
<td>Secure Hash Algorithm SHA-1</td>
<td>String stream operations</td>
<td>272</td>
</tr>
<tr>
<td>Secure Hash Algorithm SHA-224</td>
<td>Stub_Type</td>
<td>131</td>
</tr>
<tr>
<td>Secure Hash Algorithm SHA-256</td>
<td>Subprogram address</td>
<td>117</td>
</tr>
<tr>
<td>Secure Hash Algorithm SHA-384</td>
<td>subtypes</td>
<td>196</td>
</tr>
<tr>
<td>Secure Hash Algorithm SHA-512</td>
<td>Suppress_Debug_Info</td>
<td>115</td>
</tr>
<tr>
<td>Semaphores</td>
<td>Suppress_Initialization</td>
<td>115</td>
</tr>
<tr>
<td>Sequential elaboration policy</td>
<td>Suppressing external name</td>
<td>31, 32, 33</td>
</tr>
<tr>
<td>Serial_Communications</td>
<td>Suppressing initialization</td>
<td>92</td>
</tr>
<tr>
<td>Sets of strings</td>
<td>suppression of</td>
<td>92, 156</td>
</tr>
<tr>
<td>setting for not-first subtype</td>
<td>Suppression of checks</td>
<td>156</td>
</tr>
<tr>
<td>Shared</td>
<td>synonyms</td>
<td>21, 81</td>
</tr>
<tr>
<td>Shared passive packages</td>
<td>System</td>
<td>33</td>
</tr>
<tr>
<td>SHARED_MEMORY_DIRECTORY environment</td>
<td>System.Address_Image (s-addima.ads)</td>
<td>270</td>
</tr>
<tr>
<td>Shift operators</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Shift_Left</td>
<td></td>
<td></td>
</tr>
<tr>
<td>X</td>
<td>Z</td>
<td></td>
</tr>
<tr>
<td>---</td>
<td>---</td>
<td></td>
</tr>
<tr>
<td>XDR representation</td>
<td>Zero address</td>
<td></td>
</tr>
<tr>
<td>163</td>
<td>124</td>
<td></td>
</tr>
</tbody>
</table>